

Converging outer approximations of classical network correlations

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June 29th, 2022 - QIFQT Workshop, ENS Lyon

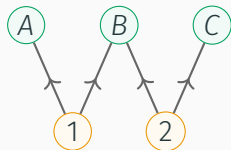
Quantum Information Theory group, Institute for Theoretical Physics, ETH Zürich
arXiv:2202.04103

Introduction

Causal compatibility

- **Causal structure** = observed and unobserved nodes + theory for unobserved nodes + connectivity

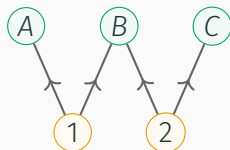
e.g., classical bilocal network



Causal compatibility

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- **Causal compatibility** = compatibility of a distribution with a given causal structure

- Bell's theorem

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- non-locality without inputs

Renou, Bäumer, Boreiri, Brunner, Gisin, Beigi, arXiv:1905.04902

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- quantifying freedom of choice in Bell's theorem

Chaves et al, arXiv:2105.05721

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- quantifying freedom of choice in Bell's theorem

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- multipartite quantum or post-quantum entanglement

Coiteux-Roy, Wolfe, Renou, arXiv:2105.09381

Causal compatibility: computational aspects

- Bell-like scenarios: **one** unobserved node \rightarrow convex problem \rightarrow easy 🍰

Causal compatibility: computational aspects

- Bell-like scenarios: **one** unobserved node → convex problem → easy 🍰
- More general networks: **several independent** unobserved nodes → nonconvex → hard 🤖

Causal compatibility: approximations

- Inner approximations: based on human or computational exploration of the search space



set of distributions compatible with network

(lower bounds for maximization problems,
certify feasibility)

Kriváchy et al, arXiv:1907.10552

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- Outer approximations:
 - nontight: entropic or algebraic inequalities
 - powerful and asymptotically tight method: **inflation**



(upper bounds for maximization problems,
certify infeasibility)

Objectives of the talk





- Define causal compatibility formally
- Construct outer approximations that are apparently converging

Causal compatibility in classical networks, a.k.a. network locality







Tensor notation

	Standard	Tensor	Components
Outcome distribution	$p(\cdot, \cdot, \cdot)$	\boxed{p}	$p(a, b, c) = \boxed{p}$









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Outcome distribution	$p(\cdot, \cdot, \cdot)$		$p(a, b, c) =$ 
Strategy with two inputs	$p_A(\cdot \cdot, \cdot)$		$p_A(a \alpha, \beta) =$ 

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$\dots\{1, \dots, n\}$			 $= \frac{1}{n}$

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Uniform distribution on...			
... $\{1, \dots, n\}$			 $= \frac{1}{n}$
... $[0, 1]$			 $= 1$

Special case: **deterministic** strategies \boxed{A} such that

$$\boxed{A} = \delta(a - f(\alpha, \beta))$$

for some function $f(\cdot, \cdot)$

Causal compatibility in classical networks

Define

$$\boxed{p} \in \mathcal{L} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ \textcircled{1} \quad \textcircled{2} \end{array} \right)$$

iff there exist \boxed{A} , \boxed{B} , \boxed{C} such that

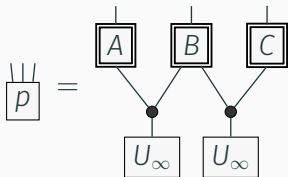
$$\boxed{p} = \begin{array}{c} \boxed{A} \quad \boxed{B} \quad \boxed{C} \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \boxed{U_\infty} \quad \boxed{U_\infty} \end{array}$$

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iff there exist \boxed{A} , \boxed{B} , \boxed{C} such that



i.e., for all a, b, c :

$$\boxed{p} \stackrel{a \ b \ c}{=} \int d\alpha d\beta \begin{array}{c} a \\ \boxed{A} \\ \alpha \end{array} \begin{array}{c} b \\ \boxed{B} \\ \alpha \ \beta \end{array} \begin{array}{c} c \\ \boxed{C} \\ \beta \end{array}$$

Causal compatibility in classical networks

Define

$$\boxed{p} \in \mathcal{L} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ \textcircled{1} \quad \textcircled{2} \end{array} \right)$$

iff there exist \boxed{A} , \boxed{B} , \boxed{C} such that

$$\boxed{p} = \begin{array}{c} \boxed{A} \quad \boxed{B} \quad \boxed{C} \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ \bullet \quad \bullet \\ \boxed{U_\infty} \quad \boxed{U_\infty} \end{array} \quad (\text{w.l.o.g.})$$

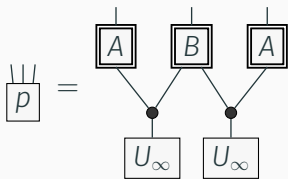
$$\text{i.e., for all } a, b, c : \begin{array}{c} a \ b \ c \\ \swarrow \quad \searrow \\ \boxed{p} \end{array} = \int d\alpha d\beta \begin{array}{c} a \\ \boxed{A} \\ \alpha \end{array} \begin{array}{c} b \\ \boxed{B} \\ \alpha \ \beta \end{array} \begin{array}{c} c \\ \boxed{C} \\ \beta \end{array}$$

Causal compatibility: identical strategies

Define

$$\boxed{p} \in \mathcal{L} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{A} \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ \textcircled{1} \quad \textcircled{2} \end{array} \right)$$

iff there exist \boxed{A} , \boxed{B} such that



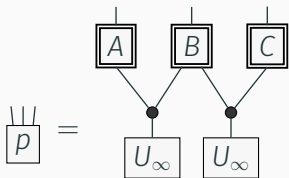
i.e., for all a, b, a' :

$$\boxed{p} \begin{array}{c} a \quad b \quad a' \\ \swarrow \quad \downarrow \quad \swarrow \\ \end{array} = \int d\alpha d\beta \begin{array}{c} a \\ \downarrow \\ \boxed{A} \\ \alpha \end{array} \begin{array}{c} b \\ \downarrow \\ \boxed{B} \\ \alpha \quad \beta \end{array} \begin{array}{c} a' \\ \downarrow \\ \boxed{A} \\ \beta \end{array}$$

Constructing converging outer approximations

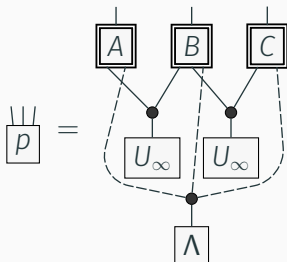
Step 1: Convexification

$\boxed{p} \in \mathcal{L} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right)$ iff there exist \boxed{A} , \boxed{B} , \boxed{C}
such that



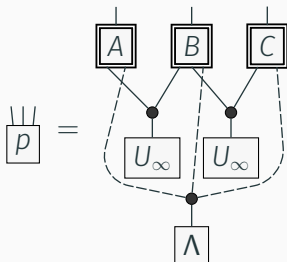
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$\boxed{p} \in \mathcal{I}_{SR} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \nearrow \quad \nwarrow \quad \nearrow \quad \nwarrow \\ \textcircled{1} \quad \textcircled{2} \end{array} \right)$ iff there exist \boxed{A} , \boxed{B} , \boxed{C} , $\boxed{\Lambda}$
such that



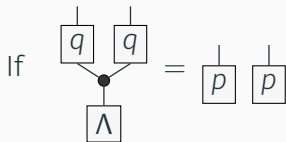
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such that



Convex but **too permissive**. Idea: forbid the agent to use Λ ?

Main lemma



Main lemma

If $\begin{array}{c} \text{---} \quad \text{---} \\ | \quad | \\ \boxed{q} \quad \boxed{q} \\ \diagdown \quad / \\ \bullet \\ | \\ \boxed{\Lambda} \end{array} = \begin{array}{c} \text{---} \quad \text{---} \\ | \quad | \\ \boxed{p} \quad \boxed{p} \end{array}$ then* $\forall \lambda : \begin{array}{c} \text{---} \\ | \\ \boxed{q} \\ | \\ \lambda \end{array} = \begin{array}{c} \text{---} \\ | \\ \boxed{p} \end{array}$

Main lemma

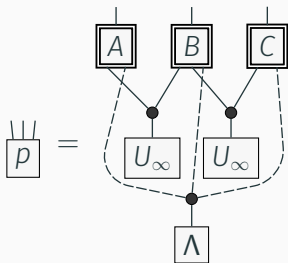
If $\begin{array}{c} \square q \quad \square q \\ \diagdown \quad / \\ \bullet \\ | \\ \square \Lambda \end{array} = \begin{array}{c} \square p \quad \square p \end{array}$ then* $\forall \lambda : \begin{array}{c} \square q \\ | \\ \lambda \end{array} = \begin{array}{c} \square p \end{array}$

*Robust version: $\int d\lambda \begin{array}{c} \lambda \\ \square \Lambda \end{array} \left\| \begin{array}{c} \square p \\ - \\ \begin{array}{c} \square q \\ | \\ \lambda \end{array} \right\|_2^2 \leq 3 \left\| \begin{array}{c} \square p \quad \square p \\ - \\ \begin{array}{c} \square q \quad \square q \\ \diagdown \quad / \\ \bullet \\ | \\ \square \Lambda \end{array} \end{array} \right\|_1$

Step 1: Convexification

$$\boxed{p} \in \mathcal{I}_{\text{SR}} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

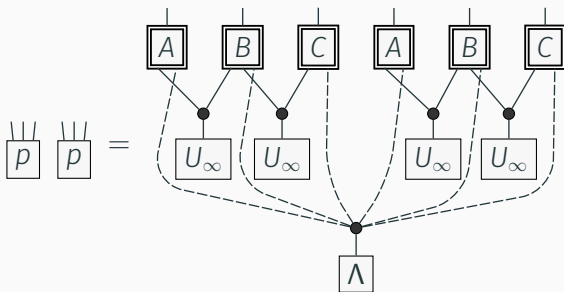
such that



Step 1: Convexification

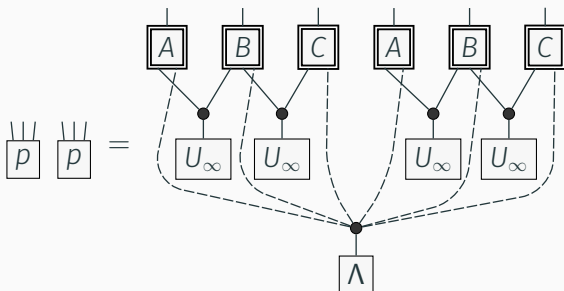
$$\boxed{p} \in \mathcal{L} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

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such that

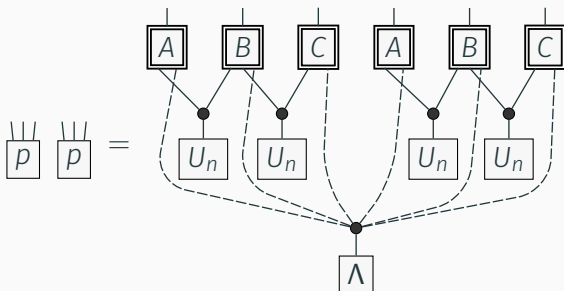


Equivalent to the original problem thanks to the main lemma

Step 2: Restricting the source output cardinality

$$\boxed{p} \in \mathcal{I}_{\text{restr}}^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

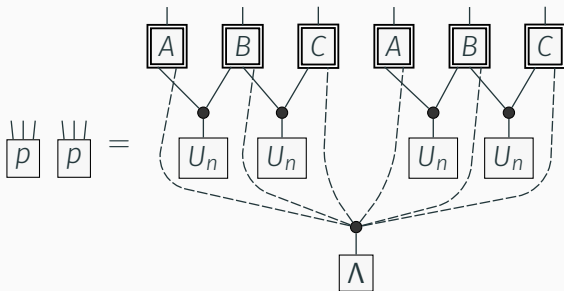
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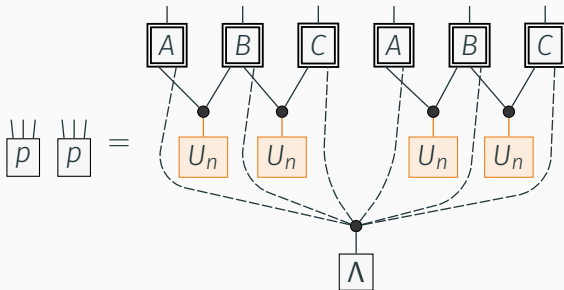
such that



Inner approximation $\mathcal{I}_{\text{restr}}^{(n)} \subseteq \mathcal{L}$, but also $\mathcal{I}_{\text{restr}}^{(n)} \xrightarrow{n \rightarrow \infty} \mathcal{L}$

Step 2: Restricting the source output cardinality

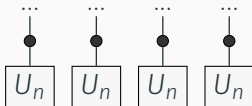
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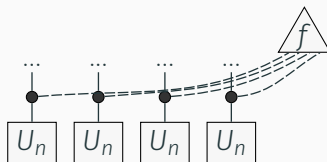
Want outer approximations \rightarrow let the agents use Λ a little bit?

Postselection for additional correlations

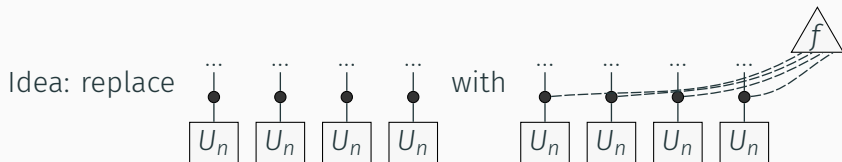
Idea: replace



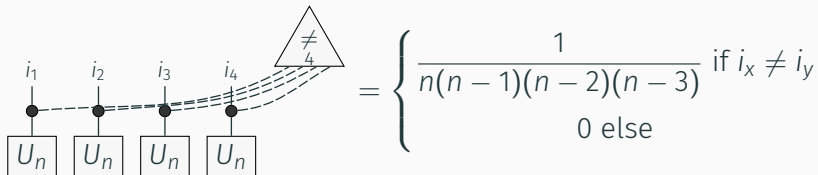
with



Postselection for additional correlations



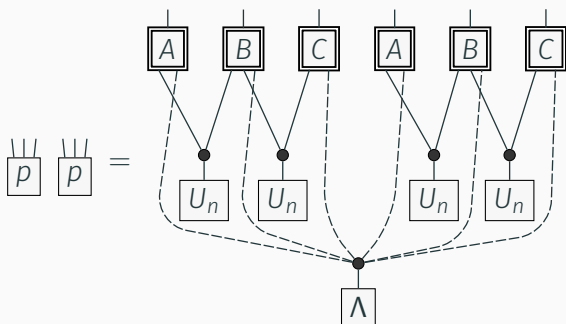
For instance, if $n \geq 4$,



Step 3: Adding postselection

$$\boxed{p} \in \mathcal{I}_{\text{restr}}^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

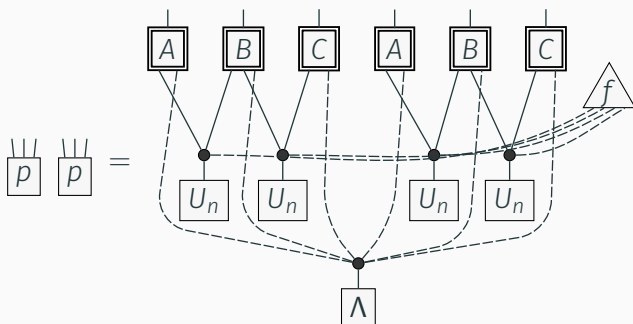
such that



Step 3: Adding postselection

$$\boxed{p} \in \mathcal{I}_f^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

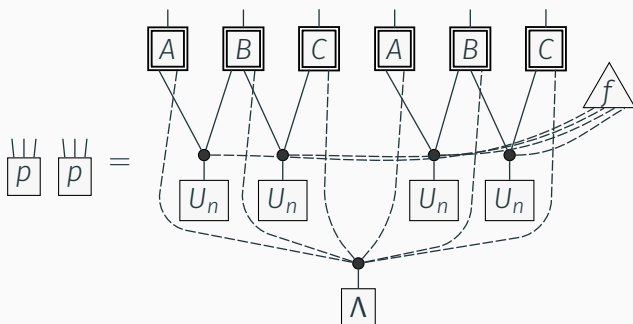
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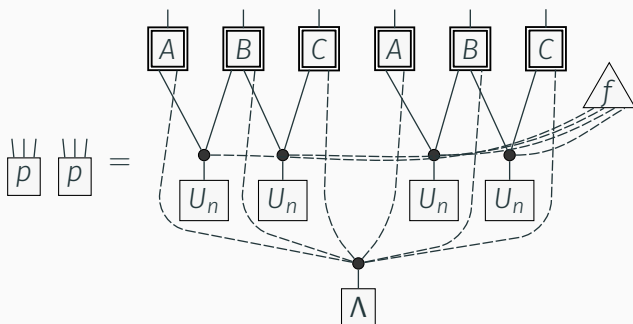


$\mathcal{I}_f^{(n)}$ and \mathcal{L} are **incomparable** for general f . But if $f \rightarrow$ “trivial” and $n \rightarrow \infty$ then $\mathcal{I}_f^{(n)} \rightarrow \mathcal{L}$

Step 3: Adding postselection

$$\boxed{p} \in \mathcal{I}_f^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

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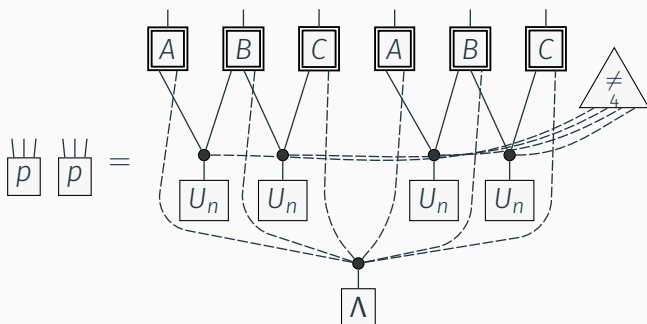


How to choose f ? The agents need some **asymmetry** to use “different parts” of Λ

Step 4: Fixing the postselection

$$\boxed{p} \in \mathcal{I}_{\neq}^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

such that

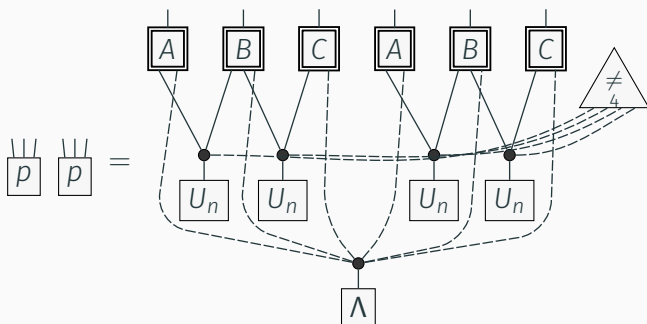


Answer: provide the agents with **distinct** values from the U_n sources

Step 4: Fixing the postselection

$$\boxed{p} \in \mathcal{I}_{\neq}^{(n)} \left(\begin{array}{c} \textcircled{A} \quad \textcircled{B} \quad \textcircled{C} \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \textcircled{1} \quad \textcircled{2} \end{array} \right) \text{ iff there exist } \boxed{A}, \boxed{B}, \boxed{C}, \boxed{\Lambda}$$

such that



Outer approximation: $\mathcal{L} \subseteq \mathcal{I}_{\neq}^{(n)}$ (left as an exercise).

Furthermore, $n \rightarrow \infty$ implies “ \neq ” \rightarrow “trivial”, so $\mathcal{I}_{\neq}^{(n)} \xrightarrow{n \rightarrow \infty} \mathcal{L}$

- Definition of $\mathcal{I}_{\neq}^{(n)}(\mathcal{N})$ for arbitrary network \mathcal{N}

Results

- Definition of $\mathcal{I}_{\neq}^{(n)}(\mathcal{N})$ for arbitrary network \mathcal{N}
- Proofs of inclusions: $\mathcal{L}(\mathcal{N}) \subseteq \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$ and $\mathcal{I}_{\neq}^{(n+1)}(\mathcal{N}) \subseteq \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$

Results

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- Proofs of inclusions: $\mathcal{L}(\mathcal{N}) \subseteq \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$ and $\mathcal{I}_{\neq}^{(n+1)}(\mathcal{N}) \subseteq \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$
- **Convergence:** if $\begin{array}{c} \text{---} \\ \diagdown \diagup \\ \square \\ p \end{array} \in \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$, then

$$\inf_{\begin{array}{c} \text{---} \\ \diagdown \diagup \\ \square \\ q \end{array} \in \mathcal{L}(\mathcal{N})} \left\| \begin{array}{c} \text{---} \\ \diagdown \diagup \\ \square \\ p \end{array} - \begin{array}{c} \text{---} \\ \diagdown \diagup \\ \square \\ q \end{array} \right\|_1 \leq \frac{c(\mathcal{N})}{\sqrt{n}} + \mathcal{O}\left(\frac{1}{n\sqrt{n}}\right)$$

Results

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- Drawback: the linear program testing $\begin{array}{c} \text{---} \\ \diagdown \quad \diagup \\ \square \\ p \end{array} \in \mathcal{I}_{\neq}^{(n)}(\mathcal{N})$ takes $e^{\text{poly}(n)}$ time/memory

Inflation correspondence

Inflation correspondence

- Two equivalent formulations: $\mathcal{I}_{\neq}^{(n)}(\mathcal{N}) = \mathcal{I}_{\text{std}}^{(n)}(\mathcal{N})$

Inflation correspondence

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Standard
inflation

Postselected
inflation

Is an outer
approximation:

Converges:

Inflation correspondence

- Two equivalent formulations: $\mathcal{I}_{\neq}^{(n)}(\mathcal{N}) = \mathcal{I}_{\text{std}}^{(n)}(\mathcal{N})$

Standard
inflation

Postselected
inflation

Is an outer
approximation:



Converges:

Navascuès, Wolfe, arXiv:1707.06476

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	Standard inflation	Postselected inflation
Is an outer approximation:		
Converges:		

- The present proof of convergence extends the original proof to same-strategy networks

Outlook

- Further developments to obtain new tools and solve open problems in classical causal compatibility? E.g., still lack noise-robust proofs of “genuine” multipartite nonlocality

Gisin, arXiv:1708.05556

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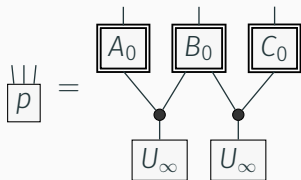
- We still lack convergent outer approximations for **quantum** causal compatibility, although some progress has been made recently

Ligthart, Gachechiladze, Gross, arXiv:2110.14659

Appendix

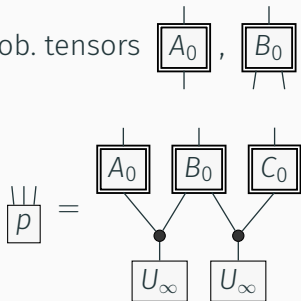
Solution: proof that $\mathcal{L} \subseteq \mathcal{I}_{\neq}^{(n)}$ for $n = 4$

Let \boxed{p} $\in \mathcal{L}$ with prob. tensors $\boxed{A_0}$, $\boxed{B_0}$, $\boxed{C_0}$ such that



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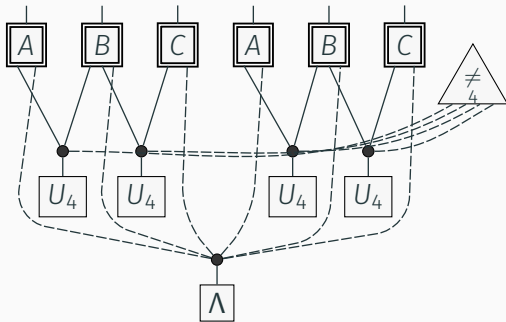


We let

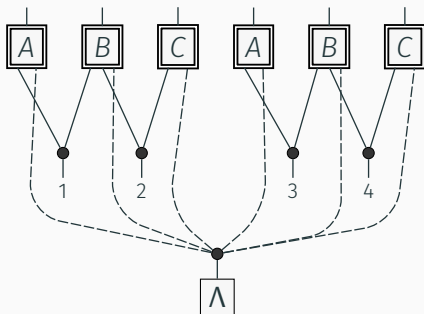
$$\boxed{\Lambda}^{(x_1, x_2, x_3, x_4)} := \begin{array}{c} x_1 \\ \boxed{U_\infty} \end{array} \begin{array}{c} x_3 \\ \boxed{U_\infty} \end{array} \begin{array}{c} x_3 \\ \boxed{U_\infty} \end{array} \begin{array}{c} x_4 \\ \boxed{U_\infty} \end{array}$$

$$\begin{array}{c} a \\ \boxed{A} \\ i \end{array} (x_1, \dots, x_4) := \begin{array}{c} a \\ \boxed{A_0} \\ x_i \end{array}, \quad \begin{array}{c} b \\ \boxed{B} \\ ij \end{array} (x_1, \dots, x_4) := \begin{array}{c} b \\ \boxed{B_0} \\ x_i \quad x_j \end{array}, \quad \begin{array}{c} c \\ \boxed{C} \\ j \end{array} (x_1, \dots, x_4) := \begin{array}{c} c \\ \boxed{C_0} \\ x_j \end{array}$$

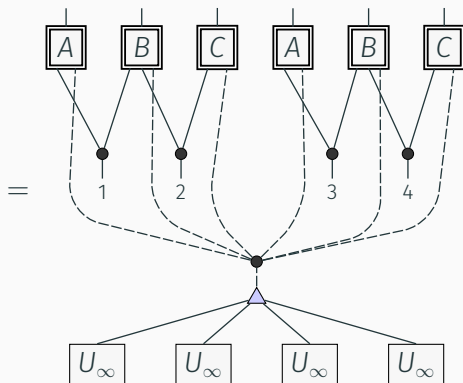
Solution: proof that $\mathcal{L} \subseteq \mathcal{I}_{\neq}^{(n)}$ for $n = 4$



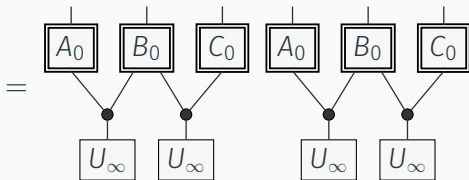
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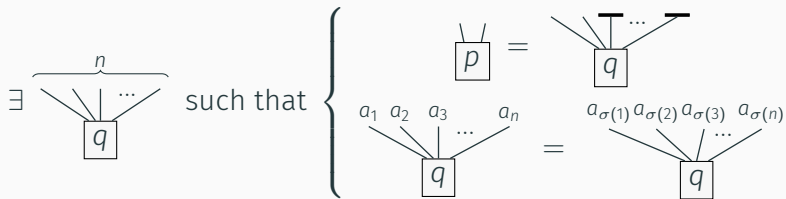


Solution: proof that $\mathcal{L} \subseteq \mathcal{I}_{\neq}^{(n)}$ for $n = 4$

$$= \begin{array}{|c|} \hline \text{|||} \\ \hline p \\ \hline \end{array} \begin{array}{|c|} \hline \text{|||} \\ \hline p \\ \hline \end{array}$$

Inflation correspondence: the case of de Finetti

For any distribution \boxed{p} ,



Inflation correspondence: the case of de Finetti

For any distribution \boxed{p} ,

$$\exists \begin{array}{c} \overbrace{\quad\quad\quad}^n \\ \diagup \quad \diagdown \quad \dots \\ \boxed{q} \end{array} \text{ such that } \left\{ \begin{array}{l} \boxed{p} = \begin{array}{c} \overline{\quad\quad\quad} \quad \dots \quad \overline{\quad\quad\quad} \\ \diagup \quad \diagdown \quad \dots \quad \diagup \quad \diagdown \\ \boxed{q} \end{array} \\ \begin{array}{c} a_1 \quad a_2 \quad a_3 \quad \dots \quad a_n \\ \diagup \quad \diagdown \quad \dots \quad \diagup \quad \diagdown \\ \boxed{q} \end{array} = \begin{array}{c} a_{\sigma(1)} \quad a_{\sigma(2)} \quad a_{\sigma(3)} \quad \dots \quad a_{\sigma(n)} \\ \diagup \quad \diagdown \quad \dots \quad \diagup \quad \diagdown \\ \boxed{q} \end{array} \end{array} \right.$$

if and only if

$$\exists \boxed{A}, \boxed{\Lambda} \text{ such that } \boxed{p} = \begin{array}{c} \triangle \neq_2 \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \\ \diagup \quad \diagdown \quad \diagup \quad \diagdown \\ \boxed{U_n} \quad \boxed{U_n} \quad \boxed{\Lambda} \end{array}$$

Inflation correspondence: the case of de Finetti

For any distribution \boxed{p} ,

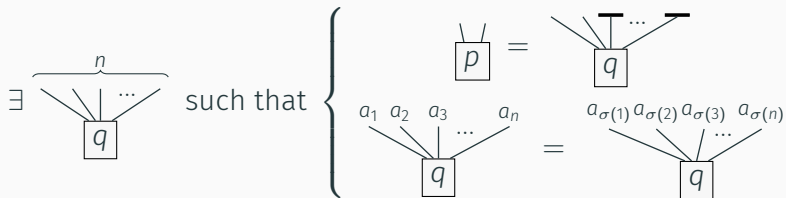
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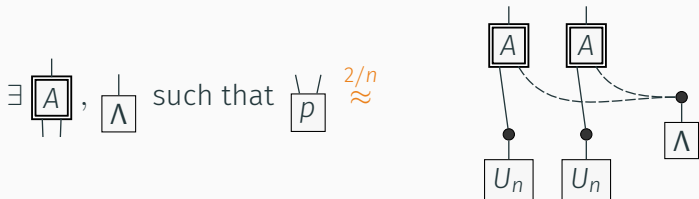
$$\exists \boxed{A}, \boxed{\Lambda} \text{ such that } \boxed{p} = \begin{array}{c} \begin{array}{c} \boxed{A} \quad \boxed{A} \\ \diagdown \quad \diagup \end{array} \\ \dots \\ \begin{array}{c} \bullet \quad \bullet \\ \diagdown \quad \diagup \end{array} \\ \boxed{U_n} \quad \boxed{U_n} \end{array} \quad \boxed{\Lambda}$$

Inflation correspondence: the case of de Finetti

For any distribution \boxed{p} ,



if and only if



De Finetti equivalence

For any distribution \boxed{p} , given $q(\{A_i = \cdot\}_{i=1}^n)$ such that

$$\boxed{p} = q(A_1 = \cdot, A_2 = \cdot)$$

and $\forall \{a_i\}_{i=1}^n, \forall \sigma \in S_n : q(\{A_{\sigma(i)} = a_i\}_{i=1}^n) = q(\{A_i = a_i\}_{i=1}^n)$,

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define first $q =: \sum_{\lambda} \boxed{\Lambda}^{\lambda} q^{(\lambda)}$ where the $q^{(\lambda)}$ are deterministic distributions,

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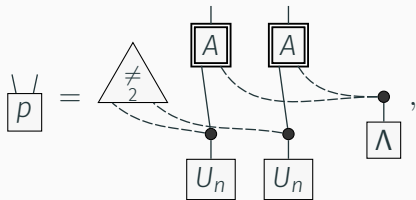
and $\forall \{a_i\}_{i=1}^n, \forall \sigma \in S_n : q(\{A_{\sigma(i)} = a_i\}_{i=1}^n) = q(\{A_i = a_i\}_{i=1}^n)$,

define first $q =: \sum_{\lambda} \boxed{\Lambda}^{\lambda} q^{(\lambda)}$ where the $q^{(\lambda)}$ are deterministic distributions, and then

$$\boxed{A}^a := q^{(\lambda)}(A_i = a)$$

De Finetti equivalence

For any distribution \boxed{p} , given \boxed{A} and $\boxed{\Lambda}$ such that



define q such that

$$q(\{A_i = a_i\}_{i=1}^n) = \sum_{\lambda} \boxed{\Lambda}^{\lambda} \frac{1}{n!} \sum_{\sigma \in S_n} \prod_{i=1}^n \boxed{A}^{a_{\sigma(i)}}_{i \lambda}$$