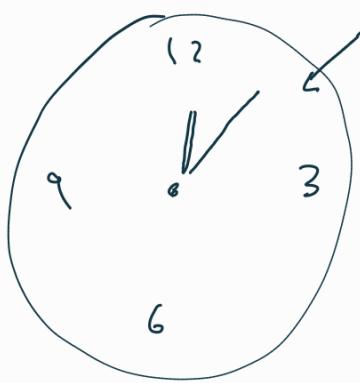


Lyons 28/June/2022

Clock tutorial by M. Woods

(See slides at end of notes)

- V. active area of research. Will aim to cover some of the most relevant results in the field.
- Q. clock emits "ticks" at regular intervals,
- Explain that we will not discuss stopwatches.
- "ticks" represent changes in a classical register. [show clock face or "register" slide]
 $(\text{dim} = N_T + 1)$



"clockwork".
 $(\text{dim} = d)$
"Quantum".

• clocks go "tick-tick" and the register changes between distinguishable states w/ each "tick".

• Reg. classical : $\{ \lvert \psi_R \rangle \}_{R=0}^{N_T}$

- definite register state :

→ associate a particular n^{th} of ticks

$$P_c(t) \otimes |j\rangle_R \xrightarrow{\text{clock has ticked } j \text{ times.}} @ \text{time } t.$$

- We want ensembles of different reg. states

$$\sum_{j=0}^{N_r} P_j(t) \tilde{P}_{j,c}(t) \otimes |j\rangle_R \quad \leftarrow \text{this describes our}$$

- Dynamics:

$$M_{CR \rightarrow CR}^{t_2 - t_1} (\tilde{P}_c \otimes |0\rangle_R) = \sum_{j=0}^{N_r} P_j(t_2) \tilde{P}_{j,c}(t_2) \otimes |j\rangle_R$$

initial state of clearword.

- Point out no zero effect.

Axiomatic Definition of clock

- Unfold to extend register states to integers via

$$|j\rangle_R := |j \bmod N_r + 1\rangle_R \quad \text{for } j \in \mathbb{Z} \setminus \{0, 1, \dots, N_r\}.$$

- Now introduce 4 axioms which restrict the dynamics to something we can think of as a clock.

- 1) the self-timing condition: For all $t_1, t_2 \geq 0$:

$$M_{CR \rightarrow CR}^{t_1 + t_2} = M_{CR \rightarrow CR}^{t_2} \circ M_{CR \rightarrow CR}^{t_1}$$

(Markovian condition).

• Necessary for clock to be autonomous (no external timing)

- 2) The zeroth order condition:

$$M_{CR \rightarrow CR}^0 = I_{CR}$$

$$\lim_{t \rightarrow 0} \| M_{CR \rightarrow CR}^t - I_{CR} \| = 0$$

3) The leading order condition:

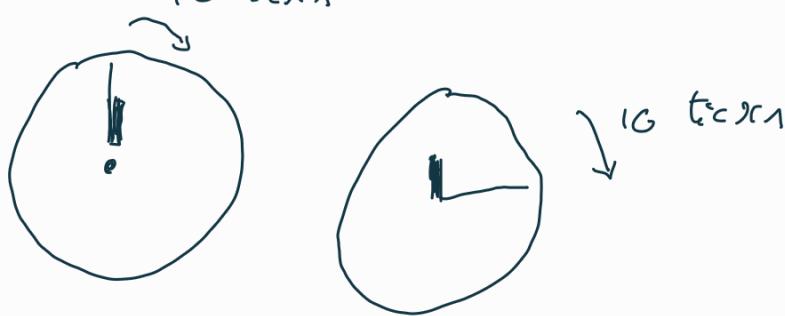
$$P(l, t | \kappa) := \text{tr} \left[M_{CR \rightarrow CR}^t (\mathcal{F}_C^\circ \otimes (\kappa \times \kappa|_R)) |_{l \times l|_R} \right]$$

$$\tilde{P}_l^{(n)}(t)$$

$$\lim_{t \rightarrow 0^+} \frac{\sum_{\ell=0}^{N_T} P(l, t | \kappa)}{P(\kappa+1, t | \kappa)} = 0 \quad \forall t \geq 0, \kappa \in \mathbb{N}.$$

4) Register translating symmetry condition

$$\text{tr}_R \left[M_{CR \rightarrow CR}^t (\mathcal{F}_C^\circ \otimes (\kappa \times \kappa|_R)) \cdot |_{\kappa+l \times \kappa+l|_R} \right] = \\ \text{tr}_R \left[M_{CR \rightarrow CR}^t (\mathcal{F}_C^\circ \otimes (\kappa' \times \kappa'|_R)) \cdot |_{\kappa'+l \times \kappa'+l|_R} \right], \quad \forall t \in \mathbb{R}, \kappa, \kappa', l \in \mathbb{Z}$$



- Using

Structure theorem [Explicit clock representation]: The pair $(\mathcal{F}_C^\circ, M_{CR \rightarrow CR}^t, +_\geq)$ forms a clock (i.e. satisfies axioms 1) to 4) if and only if \exists a Hermitian operator H , and 2 nets of linear operators $(L_j)_{j=1}^{N_L}, (I_j)_{j=1}^{N_I}$ acting on \mathcal{H}_C s.t. $\forall t \geq 0$ and $N_L \in \mathbb{N}$:

$$\tilde{P}_l^{(n)}(t) = \langle L_l(t), \mathcal{F}_C^\circ(t) \rangle$$

Jump op

$$M_{CR \rightarrow CR}(\cdot) = e^{-\cdot}$$

✓ noise

$$\begin{aligned} \dot{\rho}_{CR}(\cdot) &= -i [\tilde{H}, (\cdot)] + \sum_{j=1}^{N_L} \tilde{L}_j(\cdot) \tilde{L}_j^+ - \frac{1}{2} \{ \tilde{L}_j^+ \tilde{L}_j, (\cdot) \} \\ &\quad + \sum_{j=1}^{N_L} \tilde{S}_j(\cdot) \tilde{S}_j^+ - \frac{1}{2} \{ \tilde{S}_j^+ \tilde{S}_j, (\cdot) \} \end{aligned}$$

Jump
Opn
Error

$$\tilde{H} := H_c \otimes \mathbb{1}_R, \quad \tilde{L}_j = L_j \otimes \mathbb{1}_R, \quad \tilde{S}_j = S_j \otimes O_R$$

$$O_R := 10 \times O_R + 12 \times O_R + 13 \times O_R + \dots + 1N_T \times O_{N_T-1} + 10 \times O_{N_T}$$

if
get
cut-off
reg. clock.

- Special form of Lindblad master (also known as G.V. E.L master eq.)

{ G. Lindblad

- V. Gorini, A. Komarowski, and F. Sudarshan.

Clock accuracy

Want to measure how accurate the clock is.

Introduce "delay function of n th tick" $P_{\text{tick}}^{(n)}(t)$:

* $P_{\text{tick}}^{(n)}(t) \cdot \Delta t =$ probability of ticking $n-1$ times in interval $[0, t]$ followed by ticking once in time interval $[t, t+\Delta t]$.

Maths: $P_{CR}^{(n-1)}(t) = 1 \times 1 \times \dots \times 1_R \left[M_{CR \rightarrow CR}^C \left(P_{CR}^0 \otimes O_{O_R} \right) \right] (N-1 \times N-1)_R$

$$P_{\text{ticks}}^{(N)}(t) := \lim_{\delta t \rightarrow 0} \frac{\text{tr} \left[M_{\text{CR} \rightarrow \text{CR}}^{\text{out}} \left(S_{\text{CR}}^{(N-1)}(t) - S_{\text{CR}}^{(N-1)}(t) \right) \right]}{\delta t}$$

- A useful measure of accuracy for the N^{th} tick is R_N :

$$R_N := \left(\frac{\mu_N}{\sigma_N} \right)^2; \quad \mu_N = \int_0^\infty dt \cdot P_{\text{ticks}}^{(N)}(t) \cdot t$$

$$\sigma_N = \sqrt{\int_0^\infty dt \cdot P_{\text{ticks}}^{(N)}(t) \cdot (t - \mu_N)^2}$$

- Note: under $M_{\text{clock}}^{\text{out}} \mapsto M_{\text{CR} \rightarrow \text{CR}}^{\text{out}}$, we find $\begin{cases} \mu_N \mapsto a \mu_N \\ \sigma_N \mapsto a \sigma_N \end{cases}$
- $\therefore R_N \mapsto R_N$

Mapping $t \rightarrow at$ corresponds to a re-scaling of time. We want our maps to be invariant under such re-scaling.

Reset clocks [Slide: structure theorem]

- There are clocks where clockwork is reset to its initial state P_c^0 after every time the clear term:

$$\sum_{j=1}^{N_L} \tau_j \cdot P_c \cdot \tau_j^+ \approx P_c^0 \quad \forall P_c \in S(P_c)$$

w/ P_c^0 = initial state of the clockwork.

- Note that $\underbrace{P_{\text{ticks}}^{(N)}(t_k | t_{N-1}) \delta t}_{\text{Prob. of } N^{\text{th}} \text{ tick occurring during interval } [t_k, t_k + \delta t]} = \underbrace{P_{\text{ticks}}^{(1)}(t_k - t_{N-1}) \delta t}_{\text{Delay function 1st tick shifted in time}}$

Consequences are:

Let $P_{\text{ticks}}(t_1, t_2)$ be prob. of 1st, 2nd ticks at times t_1, t_2

$$P_{\text{ticks}}(t_1, t_2) \approx P^{(1)}(t_2 | t_1) \cdot P_{\text{ticks}}^{(1)}(t_1) = P_{\text{ticks}}^{(1)}(t_2 - t_1) \cdot P_{\text{ticks}}^{(1)}(t_1).$$

$$P_{\text{ticks}}^{(2)}(t_2) = \int_0^{t_2} dt_1 P_{\text{ticks}}(t_1, t_2) = \int_0^{t_2} dt_1 P_{\text{ticks}}^{(1)}(t_2 - t_1) P_{\text{ticks}}^{(1)}(t_1).$$

half way

in

(15 min

break)

by induction:

$$\cdot P_{\text{ticks}}^{(n)}(t_n) = \int_0^{t_n} dt_{n-1} \cdot \int_0^{t_{n-1}} dt_{n-2} \cdots \int_0^{t_3} dt_2 \int_0^{t_2} dt_1$$

$$P_{\text{ticks}}^{(1)}(t_n - t_{n-1}) \cdot P_{\text{ticks}}^{(1)}(t_{n-1} - t_{n-2}) \cdots P_{\text{ticks}}^{(1)}(t_3 - t_2) \cdot P_{\text{ticks}}^{(1)}(t_2).$$

$$=: \left(P_{\text{ticks}}^{(1)} \right)^{*n}(t_n)$$

now calculate mean and standard dev.

$$\left. \begin{array}{l} \mu_n = n \cdot \mu_1 \\ \sigma_n = \sqrt{n} \cdot \sigma_1 \end{array} \right\} \therefore R_n = \left(\frac{n \cdot \mu_1}{\sqrt{n} \cdot \sigma_1} \right)^2 = n \cdot R_1 \quad (\text{for next clock}).$$

Classical clock

Definition: A clock is classical if $\{C_{\ell c}\}_{\ell}$

$$\text{s.t. } \text{tr}_R \left[M_{\text{clock}}^t \left(S_c \otimes I_{\text{clock}} \right) \right] = \sum_m P_{m, \ell}^{(t)} |C_m \times C_{\ell}|_c \quad \forall t \geq 0 \quad \forall \ell \in \mathbb{Z}.$$

\Rightarrow Structure lemma: All classical clocks satisfy: can be viewed at the clock-world suffering decoherence

$$1) M_{\text{clock}}^t \left(\sum_{\ell, m} P_{\ell, m} |C_{\ell} \times C_{\ell}|_c \otimes |m \times m|_c \right) \approx \sum_{\ell, m} P'_{\ell, m} |C_{\ell} \times C_{\ell}|_c \otimes |m \times m|_c.$$

$$2) \underbrace{\text{Diag} \left\{ M_{\text{clock}}^t \left(S_c \otimes I_{\text{clock}} \right) \right\}}_{=: V_{\text{clock}}(t)} = e^{t M} \cdot \underbrace{\text{Diag} \left\{ S_c \otimes I_{\text{clock}} \right\}}_{=: V_{\text{clock}}(0)}$$

$$M = \tilde{N} + \tilde{T}$$

Tick generator
No tick generator

$$\tilde{N} = N \otimes I_R, \quad \tilde{T} = T \otimes O_R$$

$$[N]_{mn} = -\delta_{mn} \langle c_m | \sum_j (L_j^+ L_j + I_j^+ I_j) | c_n \rangle + \sum_j |\langle c_m | L_j | c_n \rangle|^2$$

$$[T]_{mn} = \sum_j |\langle c_m | I_j | c_n \rangle|^2$$

Example: Ladders circuit

[See slides]

- Observe:
 - 1) it's a reset :- only need to calculate delay function of 1st tick.
 - 2) it's classical

$$\text{Diag}_l(S_{CR}^{(0)}(t)) = \text{Diag}_l \left(10 \times 10 \begin{bmatrix} M_{CR \rightarrow CR}^{\delta t} (S_c^{(0)} \otimes 10 \times 10)_R \\ 0 \end{bmatrix} 10 \times 10_R \right)$$

$$= \left[e^{t \tilde{N}} \cdot \text{Diag}(S_{CR}^{(0)}) \right]_l$$

$$P_{l_{\text{ladders}}}^{(0)}(t) = \lim_{\delta t \rightarrow 0} \frac{\text{tr} \left[10 \times 10_R \left(M_{CR \rightarrow CR}^{\delta t} (S_{CR}^{(0)}(t)) - S_{CR}^{(0)}(t) \right) \right]}{\delta t}$$

Since $\tilde{N} = \sum_{l=1}^d [T e^{t \tilde{N}} \text{Diag}(S_{CR}^{(0)})]_l = e^{-t} \cdot \frac{t^{d-1}}{(d-1)!}$

is classical

can plug "top line"
into "bottom line".

Poisson distribution of $d-1$ events in interval $[0, t]$

- Question for audience: what are the "events"?

Aus: each event is one "hop" up one step of the ladder.

In each infinitesimal time step we either stay on current step or hop to next step.]

$$\therefore n_i = d, \sigma_i = \sqrt{d} \implies R_K = \kappa \cdot d.$$

• 2 results which characterize the accuracy of classical clocks:

Lemma: For every c. clock w/ accuracy R and dim d
 \exists a c. reset clock of dim d and accuracy $R' \geq R$.
 - Explain why important.

Thm:

All c. clocks of dim d satisfy

$$R_K \leq \kappa d, \kappa, d \in \mathbb{N}^+$$

i.e. we see that the simple ladder clock is optimal.

Notion proof: write $R(d+1)$ in terms of $R(d)$.

$$\text{conclude } R(d+1) = R(d) + 1.$$

Q. clock

[see slide for Quasi-ideal clock].

- Point out that it is a reset clock.

\therefore only need to calculate accuracy of 1st tick.

$$P_{t \rightarrow x}^{(1)} (+) = \ln \left[\hat{v} \cdot e^{tG} p_c^\circ e^{-tG^+} \right]$$

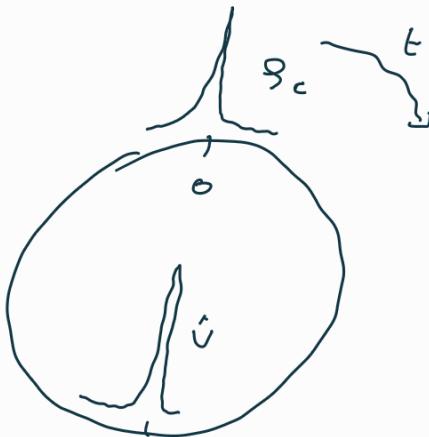
w/ $\begin{cases} \hat{v} := \sum_{j=0}^{d-1} v_j |t_j \times t_j| \geq 0 \\ G := -\hat{v} H - V \end{cases}$

- For intuition, consider $d=\infty$ case. $|t_j| >$

$$\text{choose } H = P, \hat{v} = \int_a^{\tilde{a}} v \sim \tilde{v}$$

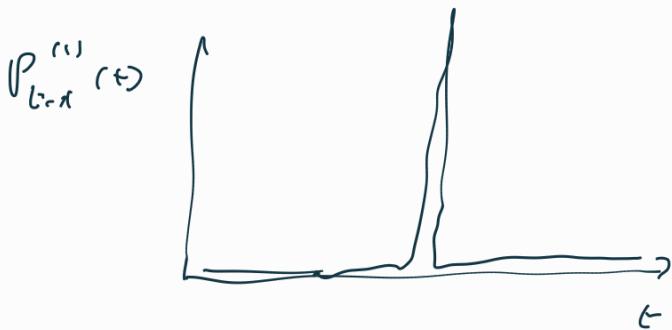
$$\propto \delta(x - \pi)$$

$$\Psi_c = (\psi(x) + i), \quad \langle x | \psi \rangle = \delta(x).$$



ϵ \rightarrow
explain

wave-function: $\langle x | e^{i(-\hat{H} - U(\hat{x}))} | \psi \rangle = e^{- \int_{x-\hbar}^x U(x') dx'} \langle x - \hbar | \psi \rangle$



$$\sigma_x = 0 \quad \therefore R_x = \infty$$

- Want to minimize in finite dim:

Def: $V_j = V_0(j/d)$; $V_0: \mathbb{R} \rightarrow \mathbb{R}$, smooth, 2π periodic.

$$\langle t_k | \psi(t_0) \rangle := \psi(t_0, t_k);$$

$$\psi(t_0, t) = A e^{-\frac{\pi^2}{\sigma^2} (t-t_0)^2} e^{i\pi(t-t_0)}$$

Theorem: $\langle t_k | e^{i(-\hat{H} + U)} | \psi(t_0) \rangle =$

$$= e^{- \int_{k-d\pi/T_0}^{t_k} V_0(x) dx} \cdot \psi(t_0 - \frac{d\pi}{T_0}; k) + \epsilon_U(k, d)$$

$$|\epsilon_U(k, d)| \leq C \cdot \text{poly}(d) e^{-\frac{\pi^2(d)}{4\sigma^2} \cdot \frac{1}{(1 + \frac{k}{d\pi})^2}}$$

$$b \geq \sup_{k \in \mathbb{N}^+} \left(\max_{x \in [0, 2\pi]} |v_0^{(r-1)}(x)| \right)$$

↗ Explain intuition: $\begin{bmatrix} -\alpha \\ -v_0 \end{bmatrix}$
 ↗ Explain proof.

$$\forall k \in \mathbb{N}_d \left(t_0 + kd/\tau \right)$$

Theorem: \exists a sequence of Q. clocks, w/ dim = 1, 2, 3, ...

$$\text{s.t. } R_{jk} \sim \epsilon d^2 \text{ as } d \rightarrow \infty$$

Theorem: All d -dimensional Q. clocks satisfy
 $R_{jk} \leq C_0 \cdot d^2$

[see slide for low dim]

- Other directions: Find the environment and study thermodynamics.

- Physical Realization of Q. clock.

- Connection w/ Lyapunov eq.

Tutorial Overview

References (randomised order):

➤ What is a (quantum) clock?

- Channel with classical register
- Clock Axioms
- Structure theorem
- Examples
- Clock accuracy
- Reset clocks
- Special case of classical clocks

➤ The ladder clock

➤ What is the most accurate classical clock?

- Optimality of reset clocks
- Upper bounds on classical clocks

- *The thermodynamics of clocks*, G. J. Milburn [Contemporary Physics Vol 61, (2020)]
- *Quantum clocks are more precise than classical ones*, Mischa P. Woods, Ralph Silva, Gilles Pütz, Sandra Stupar, and Renato Renner [PRX Quantum (2022)]
- *Autonomous Ticking Clocks from Axiomatic Principles*, Mischa P. Woods [Quantum 5, 381 (2021)]
- *Autonomous Quantum Machines & Finite-Sized Clocks*, Mischa P. Woods, Ralph Silva, Jonathan Openheim, [Annales Henri Poincaré (2019)]
- Y. Yang, L. Baumgartner, R. Silva, and R. Renner, *Accuracy enhancing protocols for quantum clocks*, (2019), [arXiv:1905.09707]
- Y. Yang and R. Renner, *Ultimate limit on time signal generation*, (2020), [arXiv:2004.07857]
- *Autonomous quantum clocks: does thermodynamics limit our ability to measure time?* Paul Erker, Mark T. Mitchison, Ralph Silva, Mischa P. Woods, Nicolas Brunner, and Marcus Huber [Phys. Rev. X 7, 031022 (2017)]
- S. Rankovic, Y. Liang, and R. Renner, *Quantum clocks and their*

➤ The Quasi-ideal quantum clocks

- Delay function of 1st tick
- Analogy with x and p
- Quasi-ideal control theorem
- Accuracy and optimality

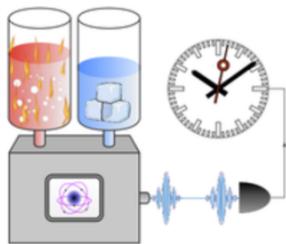
➤ Quantum advantage in low dimensions

- How to compute it

➤ Other interesting things we did not discuss due to time constraints

- synchronization - the Alternate Ticks Game. arXiv:1506.01373, (2015).
- Autonomous Temporal Probability Concentration: Clockworks and the Second Law of Thermodynamics, Emanuel Schwarzhans, Maximilian P. E. Lock, Paul Erker, Nicolai Friis, Marcus Huber, [Phys. Rev. X 11, 011046 (2021)]
- Measuring the thermodynamic cost of timekeeping, A. N. Pearson, Y. Guryanova, P. Erker, E. A. Laird, G. A. D. Briggs, M. Huber, N. Ares, [Phys. Rev. X 11, 021029 (2021)]

Clock Examples:



1) Thermodynamic "heat-engine" clock

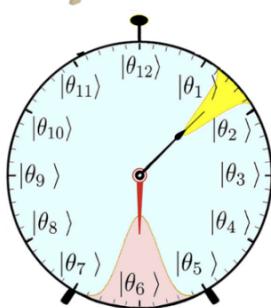
[P. Erker et. al. PRX (2017)]

$$\begin{aligned} L_1 &= \sqrt{\gamma_h} \sigma_h, & L_2 &= \sqrt{\gamma_h e^{-\beta_h E_h}} \sigma_h^\dagger, \\ L_3 &= \sqrt{\gamma_c} \sigma_c, & L_4 &= \sqrt{\gamma_c e^{-\beta_c E_c}} \sigma_c^\dagger, \\ J_1 &= \sqrt{\Gamma} |0\rangle\langle d-1|_w, & J_2 = J_3 = J_4 &= 0, \end{aligned}$$



2) Ladder clock [S. Stupar et.al, arXiv:1806.08812 (2018)]

$$\begin{aligned} L_j &= |c_{j+1}\rangle\langle c_j|, & J_j &= 0, \quad j = 0, 1, 2, \dots, d-1 \\ L_d &= 0, & J_d &= |c_1\rangle\langle c_d|, \end{aligned}$$



3) Quasi-Ideal clock [M. Woods et al, PRX Quantum (2019)]

$$L_j = 0, \quad J_j = \sqrt{2V_j} |\psi_C\rangle\langle t_j|, \quad j = 0, 1, 2, \dots, d$$

[A. Peres AJP (1980)]
[M. Woods et al, Annales Henri Poincaré (2018)]

Ladder Clock

- $N_L = d$ with

$$\begin{aligned} L_j &= |c_{j+1}\rangle\langle c_j|, & J_j &= 0, \quad j = 1, 2, 3, \dots, d-1 \\ L_d &= 0, & J_d &= |c_1\rangle\langle c_d|, \\ H &= 0, \end{aligned}$$

$$\rho_C^0 = |c_1\rangle\langle c_1|$$

- No tick and tick generators:

$$\mathcal{N} = \begin{bmatrix} -1 & 0 & \dots & & \dots & 0 \\ 1 & -1 & 0 & \dots & & \vdots \\ 0 & 1 & -1 & 0 & \dots & \\ \vdots & \ddots & \ddots & \ddots & \ddots & \\ & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & & 0 & 1 & -1 & 0 \\ 0 & \dots & \dots & 0 & 1 & -1 \end{bmatrix}, \quad [\mathcal{T}]_{i,j} = \begin{cases} 1 & \text{if } i = 1, j = d \\ 0 & \text{otherwise} \end{cases}$$

Low dimensional accuracy of quasi-ideal clocks

