

# Graphical Approach to the Drinfeld Centre

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# Plan

## Part 1. Categorical considerations

- Graphical calculus and Drinfeld centre
- Tube category
- Equivalence with Drinfeld centre

## Part 2. Applications in conformal field theory

- Modular invariants as representations of the tube category
- Extending the trace of a module category

# Graphical Calculus

## Objects and Morphisms

Let  $\mathcal{C}$  be a monoidal category. Within the graphical calculus of  $\mathcal{C}$ :

objects are depicted  
as labelled strands  
(except the tensor identity)

$X$

morphisms are depicted  
as labelled boxes

$X$   
 $\alpha$   
 $Y$

composition is depicted  
by vertical juxtaposition

$\alpha$   
 $\beta$

tensor product is depicted  
by horizontal juxtaposition

$\alpha$        $\beta$

We adopt the convention that diagrams are read top to bottom.  
Even when not using graphical calculus we often omit the ' $\otimes$ ' symbol.

# Graphical Calculus

## Half-braidings

A *half-braiding* on  $X$  in  $\mathcal{C}$  is a collection of natural isomorphisms

$$\tau_G: GX \rightarrow XG \quad \text{such that} \quad \begin{array}{c} G H X \\ \tau_{GH} \\ \hline X G H \end{array} = \begin{array}{c} G H X \\ \tau_H \\ \hline \tau_G \\ X G H \end{array}.$$

for all  $G, H$  in  $\mathcal{C}$ . From a graphical perspective the condition that  $\tau_G$  is natural allows us to 'push' morphisms through  $\tau$ :

$$\begin{array}{c} G X \\ \alpha \\ \hline \tau_H \\ X H \end{array} = \begin{array}{c} G X \\ \tau_G \\ \hline \alpha \\ X H \end{array}.$$

Note that this only works when pushing morphisms from the top-left to the bottom-right, endomorphisms of  $X$  cannot be pushed through.

# Graphical Calculus

## Drinfeld Centre

The *Drinfeld centre* of  $\mathcal{C}$  categorifies the notion of the centre of a group or monoid. Whereas the classical centre consists of all elements which commute with everything, in the categorical context we have to specify *how* they commute. This is encoded by a half-braiding.

- Objects in  $Z(\mathcal{C})$  are pairs  $(X, \tau)$  where  $X$  is in  $\mathcal{C}$  and  $\tau$  is a half-braiding on  $X$ .
- $\text{Hom}_{Z(\mathcal{C})}((X, \tau), (Y, \nu))$  is given by the subspace  $\text{Hom}_{\mathcal{C}}(X, Y)$  defined by the condition that  $\alpha \in \text{Hom}_{\mathcal{C}}(X, Y)$  satisfies

$$\begin{array}{c} G \quad X \\ | \quad \square \alpha \\ \nu_G \\ | \\ Y \quad G \end{array} = \begin{array}{c} G \quad X \\ | \quad \square \tau_G \\ \alpha \\ | \\ Y \quad G \end{array} \quad \forall G \text{ in } \mathcal{C}.$$

In other words  $\text{Hom}_{Z(\mathcal{C})}((X, \tau), (Y, \nu))$  is the space of morphisms which *can* be pushed through top-right to bottom-left.

**Applications:** Mathematical physics and the representation theory of monoidal categories in general.

# Graphical Calculus

## Dual Objects

We now suppose that  $\mathcal{C}$  is rigid, i.e. for every  $X$  in  $\mathcal{C}$  there is an object  $X^\vee$  and  ${}^\vee X$  together with morphisms

$$X \xrightarrow{\quad} X^\vee, \quad X^\vee \xrightarrow{\quad} X, \quad {}^\vee X \xrightarrow{\quad} X \quad \text{and} \quad X \xrightarrow{\quad} {}^\vee X$$

such that the 'S-relations' holds:

$$\begin{array}{c} X \\ \curvearrowright \end{array} = \begin{array}{c} X \\ | \end{array} = \begin{array}{c} X \\ \curvearrowleft \end{array}.$$

Furthermore we chose a *pivotal structure* in  $\mathcal{C}$  i.e. a natural isomorphism  ${}^\vee X \rightarrow X^\vee$ . This allows us to identify left and right duals. We can now consider the left and right dimensions of an object  $X$ ,

$$d_r(X) = X \bigcirc X^\vee \quad \text{and} \quad d_l(X) = {}^\vee X \bigcirc X.$$

Finally we suppose that our pivotal structure is *spherical*, i.e.

$$d_l(X) = d_r(X) =: d(X).$$

# Graphical Calculus

## Braidings and Factorizability

At times we will also suppose  $\mathcal{C}$  is equipped with a *braiding*  $\sigma_{XY}: XY \rightarrow YX$  which is depicted by the over-crossing,



Naturality of the braiding allows morphisms to pass over and under strands,

$$\begin{array}{ccc} X & Y \\ \boxed{\alpha} & \end{array} = \begin{array}{ccc} X & Y \\ & \boxed{\alpha} \end{array} \quad \text{and} \quad \begin{array}{ccc} X & Y \\ & \boxed{\nu} \end{array} = \begin{array}{ccc} X & Y \\ \nu & \end{array}.$$

A choice of braiding on  $\mathcal{C}$  determines a functor

$$\Phi: \mathcal{C} \boxtimes \overline{\mathcal{C}} \rightarrow Z(\mathcal{C})$$

$$X \boxtimes Y \mapsto (XY, \begin{array}{c} GXY \\ \diagup \diagdown \end{array}).$$

The braiding is called *factorisable* if  $\Phi$  is an equivalence.

# Graphical Calculus

## Fusion Categories

For the rest of this talk we assume that  $\mathcal{C}$  is a *linear category* over an algebraically closed field (i.e. enriched over Vect). We also assume that  $\mathcal{C}$  is *semisimple* and admits a finite complete set of simple objects  $\mathcal{I}$ .

Such categories are equivalent to  $\mathcal{RC} := \text{Fun}^{\text{op}}(\mathcal{C}, \text{Vect})$ . Indeed, the functor

$$\flat: \mathcal{RC} \rightarrow \mathcal{C}$$

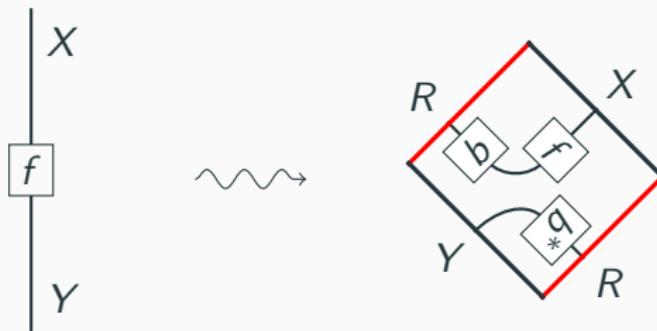
$$F \mapsto \bigoplus_S F(S) \cdot S$$

(where the sum ranges over  $\mathcal{I}$ ) is a inverse to the Yoneda embedding  $\sharp$ . When such a category is equipped with a rigid monoidal product it is known as a *fusion category*. Its dimension is given by  $d(\mathcal{C}) = \sum_S d(S)^2$ .

# The Tube Category

## An Introduction

Let  $\mathcal{C}$  be a spherical fusion category and let  $f$  be in  $\text{Hom}_{\mathcal{C}}(X, Y)$ .



$$\text{Hom}_{\mathcal{TC}}(X, Y) = \bigoplus_R \text{Hom}_{\mathcal{C}}(RX, YR)$$

$\mathcal{TC}$  shares the same objects as  $\mathcal{C}$  but has more morphisms.

Morphisms in  $\mathcal{TC}$  are described by diagrams in  $\mathcal{C}$  drawn on a cylinder.

# The Tube Category

## Representations of $\mathcal{TC}$

We consider the following morphism in  $\mathcal{TC}$

$$c_{G,X} = \begin{array}{c} G \quad X \\ \diagdown \quad \diagup \\ G \quad X \\ \diagup \quad \diagdown \\ G \quad G \end{array}$$

Any morphism in  $\mathcal{TC}$  may be re-written in the following form:

$$\begin{array}{c} G \quad X \\ \diagdown \quad \diagup \\ Y \quad G \\ \alpha \end{array} = \begin{array}{c} G \quad Y \\ \diagdown \quad \diagup \\ X \quad G^\vee \\ \beta \end{array} \circ \begin{array}{c} Y \quad G \\ \diagdown \quad \diagup \\ G \quad G^\vee \\ \gamma \end{array} \circ \begin{array}{c} Y \\ G^\vee \quad G \end{array}$$

Therefore a (contravariant) representation  $\mathcal{F}: \mathcal{TC} \rightarrow \underline{\text{Vect}}$  is simply a representation of  $\mathcal{C}$  together with an action of the  $c_{G,X}$  morphisms i.e. a collection of isomorphisms  $\kappa_{G,X}: \mathcal{F}(GX) \rightarrow \mathcal{F}(XG)$  which are natural in  $G$  and  $X$  and satisfy  $\kappa_{GH,X} = \kappa_{H,XG} \circ \kappa_{G,HX}$ .

# The Tube Category

## Equivalence of Representations with $\mathcal{Z}(\mathcal{C})$

$\mathcal{Z}(\mathcal{C})$  is a category with objects  $(F, \tau)$  where  $F \in \mathcal{C}$  and  $\tau: (- \otimes F) \rightarrow (F \otimes -)$  is a half-braiding i.e. satisfies:

$$\tau_{GH}^\sharp = (\tau_G \otimes \text{id}_H)^\sharp \circ (\text{id}_G \otimes \tau_H)^\sharp.$$

$$\kappa_{GH,X} = \kappa_{H,XG} \circ \kappa_{G,HX}$$

As  $\mathcal{C}$  is fusion the Yoneda embedding  $F \mapsto F^\sharp$  is an equivalence. This induces an equivalence  $(F, \tau) \mapsto (F^\sharp, \tau^\sharp)$  between  $\mathcal{Z}(\mathcal{C})$  and  $\mathcal{Z}(\mathcal{RC})$  where

$$\begin{array}{ccc} \tau_{G^\vee}^\sharp: (G^\vee \otimes F)^\sharp & \rightarrow & (F \otimes G^\vee)^\sharp \\ \parallel & & \parallel \end{array}$$

$$F^\sharp \circ (G \otimes -) \qquad \qquad F^\sharp \circ (- \otimes G)$$

$$\kappa_{G,X}: F^\sharp(GX) \rightarrow F^\sharp(XG)$$

An object in  $\mathcal{Z}(\mathcal{C}) = \mathcal{Z}(\mathcal{RC})$  is an object in  $\mathcal{RC}$  together with together with isomorphisms  $\kappa_{G,X}$  as above, i.e. an object in  $\mathcal{RC}$ .

# The Tube Category

## Primitive Idempotents when $\mathcal{C}$ is Modular

We now suppose that  $\mathcal{C}$  is equipped with a *factorisable* braiding, in particular, we have a canonical equivalence

$$\Phi: \mathcal{C} \boxtimes \overline{\mathcal{C}} \rightarrow \mathcal{Z}(\mathcal{C}) = \mathcal{RT}\mathcal{C}.$$

Therefore elements of the form  $\Phi(I \boxtimes J)$  for  $I, J \in \mathcal{I}$  form a complete set of simples in  $\mathcal{RT}\mathcal{C}$ . Helpfully, the simple representation  $\Phi(I \boxtimes J)$  is represented by the idempotent

$$\epsilon_I^J = \frac{1}{d(\mathcal{C})} \sum_S d(S) \begin{array}{c} \text{Diagram of a } 3 \times 3 \text{ grid with red lines forming a diagonal path from top-left to bottom-right, with labels } S \text{ in the cells.} \\ \text{The path is: } (S, I) \rightarrow (I, J) \rightarrow (J, S) \end{array} \in \text{End}_{\mathcal{TC}}(IJ).$$

Therefore every representation of  $\mathcal{TC}$  is represented by an idempotent. This effectively allows us to translate any computation in  $\mathcal{Z}(\mathcal{C})$  into a computation of idempotents in  $\mathcal{TC}$ .

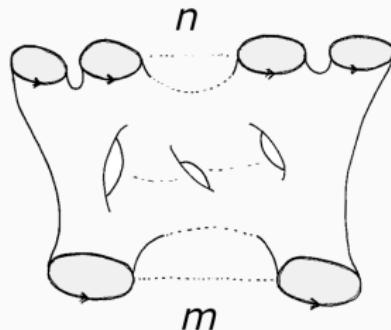
# Conformal Field Theory

The cobordism category  $\underline{\text{Seg}}^{\text{op/cl}}$

$$\text{Obj}(\underline{\text{Seg}}^{\text{cl}}) = \left\{ \underline{n} := \bigsqcup_n S_1 \mid n \in \mathbb{N} \right\}$$

$$\underline{n} = \underbrace{\circ \circ \circ \dots \circ}_{n}$$

$\text{Hom}_{\underline{\text{Seg}}^{\text{cl}}}(\underline{n}, \underline{m}) = \{[X] \mid \begin{array}{l} X \text{ is a Riemann surface with } n \text{ incoming} \\ \text{boundaries and } m \text{ outgoing boundaries.} \end{array}\}$



There is an analogous category  $\underline{\text{Seg}}^{\text{op}}$  where the objects are 1-dimensional manifolds with two boundary points labelled by *boundary conditions*.

## Conformal Field Theory

## The Algebraic Data

## Full 2D Conformal Field Theory

$$\mathcal{F}: \mathbf{Seg}^{\text{op}/\text{cl}} \rightarrow \mathbf{\underline{Vect}}$$

4

## Open CFT

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## Closed CFT

The *vertex operator algebra* (VOA) structure axiomatises (one chiral half of) the genus zero part of  $\mathcal{F}$ . The category of modules  $\mathcal{C}$  over a (rational) VOA is a *modular tensor category* i.e. a linear category which is:

- spherical
- fusion
- equipped with a factorisable braiding

### Extra data:

## Module category over $\mathcal{C}$      Modular invariant of $\mathcal{C}$

# Conformal Field Theory

## Module Categories

A module category over  $\mathcal{C}$  is a linear category  $\mathcal{B}$  together with a monoidal functor

$$\mathcal{M}: \mathcal{C} \rightarrow \text{End}(\mathcal{B}).$$

In other words it is simply a categorification of a module over a ring.

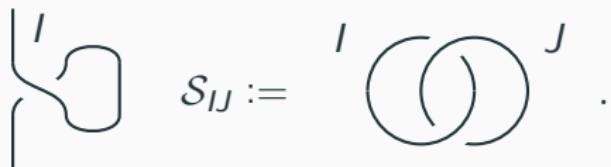
# Conformal Field Theory

## Modular Invariants

The value of a closed CFT  $\mathcal{F}$  on a torus is given by the *partition function*  $Z$ .

$$\mathcal{F}(\mathbb{T}_\tau) =: Z(\tau) = \sum_{I,J \in \mathcal{I}} Z_{IJ} \chi_I(\tau) \chi_J(\tau)^* \in \mathbb{C}$$

where  $\tau \in \mathbb{H}$ ,  $T_\tau = \mathbb{C}/(\mathbb{Z} + \mathbb{Z}\tau)$  and  $\chi_I$  is the character of  $I$ . Conformal invariance implies invariance of  $Z$  under the action of the mapping class group on  $\mathbb{H}$ . MTCs come with a special representation of  $\mathrm{PSL}_2(\mathbb{Z})$  given by their *modular data*:

$$\mathcal{T}_{IJ} := \delta_{I,J} \quad \text{and} \quad \mathcal{S}_{IJ} := \begin{cases} I & \text{if } I = J \\ 1 & \text{if } I \neq J \end{cases}.$$


Conformal invariance of  $Z(\tau)$  is equivalent to the  $\mathcal{I} \times \mathcal{I}$ -matrix of multiplicities  $(Z_{IJ})$  commuting with the modular data. We define a *modular invariant* to be any such non-negative integer matrix  $Z$ .

# Conformal Field Theory

## Modular Invariants as Representations of $\mathcal{TC}$

As described before, a complete set of simple objects in  $\mathcal{RTC}$  is indexed by  $\mathcal{I} \times \mathcal{I}$  so the isomorphism classes in  $\mathcal{RTC}$  are in bijection with non-negative integer  $\mathcal{I} \times \mathcal{I}$ -matrices.

$$Z: \mathcal{RTC} \rightarrow \text{Mat}_{\mathcal{I} \times \mathcal{I}}(\mathbb{N})$$
$$F \mapsto \left( \dim \text{Im } F(\epsilon_i^j) \right)$$

We may therefore assign to any  $F \in \mathcal{RTC}$  a *candidate* modular invariant by simply considering its isomorphism class. In fact we can go further and *define* a modular invariant to be a representation of  $\mathcal{TC}$  satisfying its own notion of invariance under the modular group.

# Conformal Field Theory

## Recovering the Mapping Class Group

Under what condition should an object  $F \in \mathcal{RTC}$  be a modular invariant?

$$\text{tr } F \left( \begin{array}{ccccc} & & G & & X \\ & & \alpha & & \\ & & \alpha & & \\ & & G & & X \\ X & & & & X \\ & & & & \end{array} \right)$$

**Theorem (to appear, H.)**

*The isomorphism class  $Z(F)$  gives a modular invariant if and only if  $\text{tr } F$  is invariant under the action of  $\text{SL}_2(\mathbb{Z}) = \text{MCG}(\mathbb{T}^2)$ .*

$$\text{tr } F \left( \begin{array}{ccccc} & & G & & X \\ & & \alpha & & \\ & & \alpha & & \\ & & G & & X \\ X & & & & X \\ & & & & \end{array} \right) = \text{tr } F \left( \begin{array}{ccccc} & & G & & X \\ & & \alpha & & \\ & & \alpha & & \\ & & G & & X \\ X & & & & X \\ & & & & \end{array} \right)$$

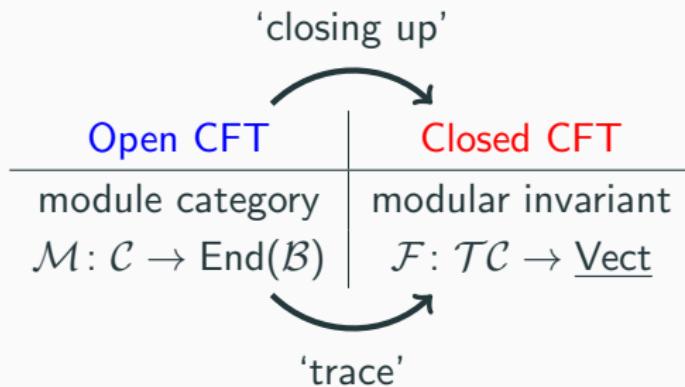
For  $\alpha = \epsilon_I^J$  this is equivalent to  
 $(\mathcal{T}^{-1}Z(F)\mathcal{T})_{IJ} = Z(F)_{IJ}$

$$\text{tr } F \left( \begin{array}{ccccc} & & G & & X \\ & & \alpha & & \\ & & \alpha & & \\ & & G & & X \\ X & & & & X \\ & & & & \end{array} \right) = \text{tr } F \left( \begin{array}{ccccc} X^V & & G & & \\ & & \alpha & & \\ & & \alpha & & \\ & & G & & X^V \\ G & & & & X^V \\ & & & & \end{array} \right)$$

For  $\alpha = \epsilon_I^J$  this is equivalent to  
 $(S^{-1}Z(F)\mathcal{S})_{IJ} = Z(F)_{IJ}$

# Conformal Field Theory

## A Unified Approach to Open/Closed CFT



There is a notion of trace for monoidal functors. However it simply gives the following representation of  $\mathcal{C}$ :

$$\text{Tr } \mathcal{M}: \mathcal{C} \rightarrow \underline{\text{Vect}}$$

$$X \mapsto \text{Hom}_{\text{End}(\mathcal{B})}(\mathbf{1}, \mathcal{M}(X)).$$

# Conformal Field Theory

## Extending the Trace

To extend  $\text{Tr } \mathcal{M}$  to a representation of  $\mathcal{TC}$ , which we denote  $\mathcal{TM}$ , we have to specify the value of  $\mathcal{TM}$  on the  $c_{G,X}$  morphisms. Setting

$$\mathcal{TM} \left( \begin{array}{c} G & X \\ \diagup & \diagdown \\ G & X \\ \diagdown & \diagup \\ G & X \end{array} \right) : \begin{array}{c} G & X \\ \text{---} & \text{---} \\ \alpha & \end{array} \mapsto \begin{array}{c} X & G \\ \text{---} & \text{---} \\ \alpha & \end{array}$$

defines a unique object in  $\mathcal{RTC}$  (where the blue colour denotes evaluation under  $\mathcal{M}$ ) when  $\mathcal{M}$  ‘induces a pivotal structure’.

**Theorem (Theorem 6.7, H.)**

$\mathcal{TM}$  is a symmetric, commutative Frobenius algebra in  $\mathcal{RTC}$ .

This result together with the work of Kong and Runkel [KR, Theorem 3.4] implies that  $\mathcal{TM}$  is a modular invariant if and only if  $d(\mathcal{TM}) = d(\mathcal{C})$ .



Leonard Hardiman.

**Extending the trace of a pivotal monoidal functor.**

To appear in *Comm. Math. Phys.*

<https://arxiv.org/abs/1911.09024>.



Liang Kong and Ingo Runkel.

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*Comm. Math. Phys.*, 292(3):871912, 2009.

<https://arxiv.org/abs/0807.3356>.