

# Möbius Orthogonality for the Zeckendorf Sum-of-Digits Function

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Monday, October 16, 2017

Joint work with Michael Drmota and Lukas Spiegelhofer

# Möbius function

The Möbius function is defined by

$$\mu(n) = \begin{cases} (-1)^k & \text{if } n \text{ is squarefree and} \\ & k \text{ is the number of prime factors} \\ 0 & \text{otherwise} \end{cases}$$

A sequence  $\mathbf{u}$  is **orthogonal to the Möbius function**  $\mu(n)$  if

$$\sum_{n \leq N} \mu(n) u_n = o\left(\sum_{n \leq N} |u_n|\right) \quad (N \rightarrow \infty).$$

Old Heuristic - Mobius Randomness Law

Any "reasonably defined (easy)" bounded sequence independent of  $\mu$  is orthogonal to  $\mu$ .

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# Orthogonality to $\mu$

## Results

- Constant sequences  $\Leftrightarrow$  PNT
- Periodic sequences  $\Leftrightarrow$  PNT in arithmetic Progressions
- Quasiperiodic sequences  $f(n) = F(\alpha n \bmod 1)$  - Davenport
- Nilsequences - Green and Tao
- Horocycle Flows - Bourgain, Sarnak and Ziegler
- Dynamical systems with discrete spectrum

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# Sarnak Conjecture

## Definition

A dynamical system is said to be deterministic, if its topological entropy is 0.

## Conjecture (Sarnak conjecture, 2010)

Every bounded complex sequence  $\mathbf{u} = (u_n)_{n>0}$  that is obtained by a deterministic dynamical system is orthogonal to the Möbius function  $\mu(n)$ .

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# Chowla Conjecture

## Conjecture (Chowla)

Let  $0 \leq a_1 < a_2 < \dots < a_t$  and  $k_1, k_2, \dots, k_t$  in  $\{1, 2\}$  not all even, then as  $N \rightarrow \infty$

$$\sum_{n \leq N} \mu^{k_1}(n + a_1) \mu^{k_2}(n + a_2) \cdots \mu^{k_t}(n + a_t) = o(N).$$

## Theorem (Sarnak)

The Chowla Conjecture implies the Sarnak Conjecture.

## Theorem (Tao)

The logarithmic version of the Sarnak Conjecture implies the logarithmic version of the Chowla Conjecture.

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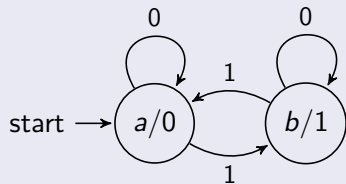
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# Motivation

## Automatic sequence

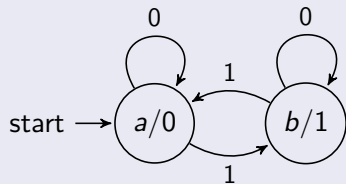


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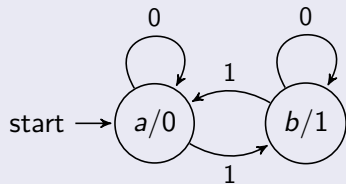


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## Theorem (M., 2016)

Every automatic sequence  $(a_n)_{n \geq 0}$  fulfills the Sarnak Conjecture

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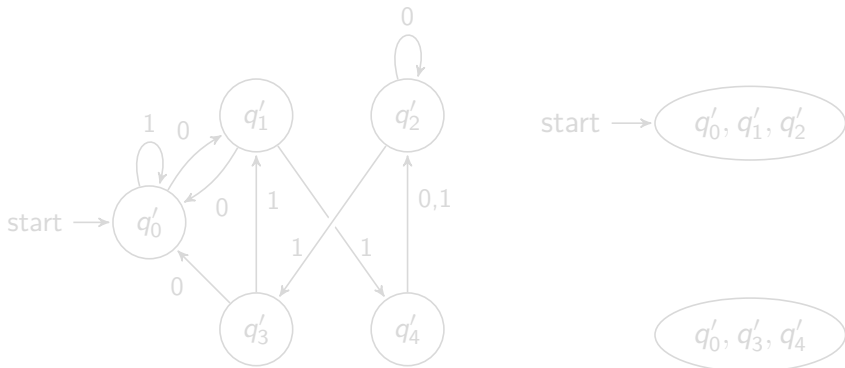
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For every strongly connected automaton  $A$ , there exists a naturally induced transducer  $\mathcal{T}_A$ . All other naturally induced transducers can be obtained by changing the order on the elements of  $Q$ .

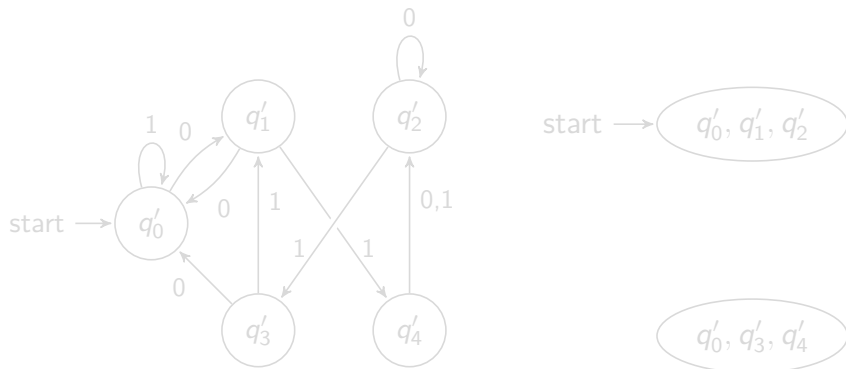
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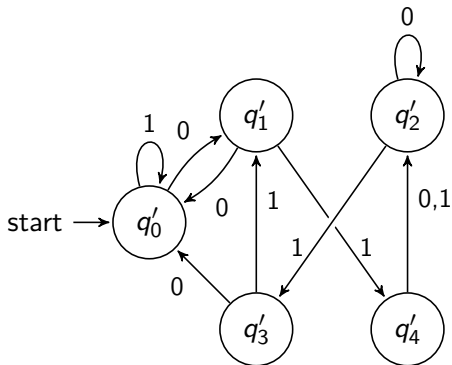




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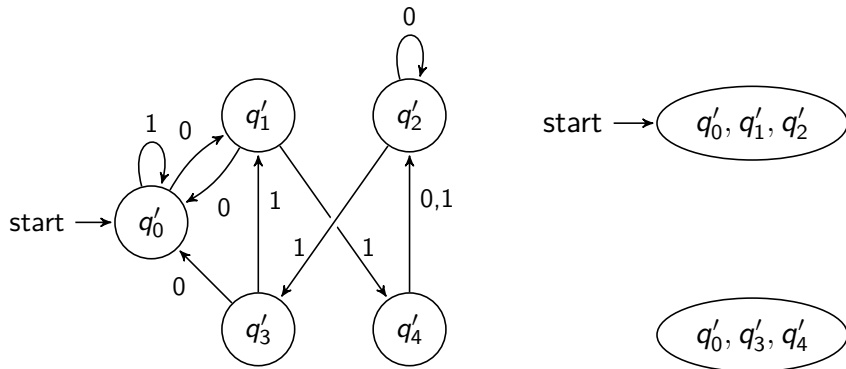
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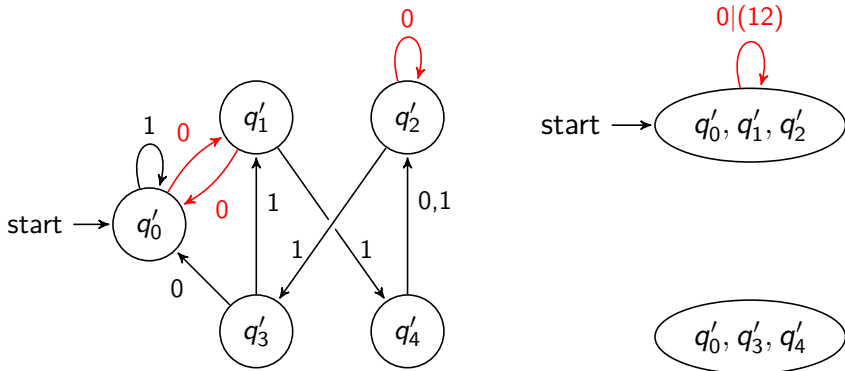
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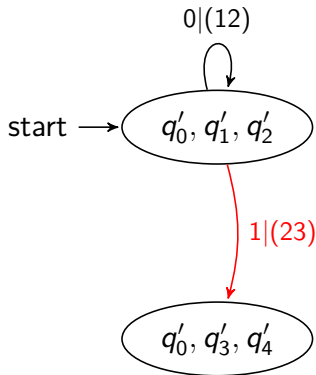
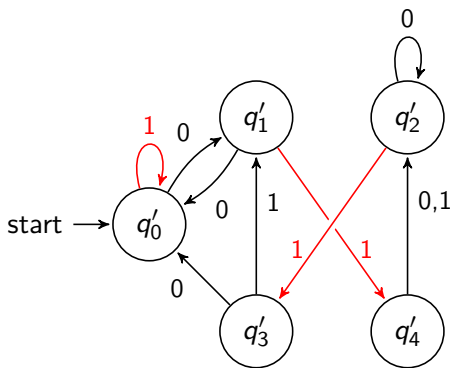
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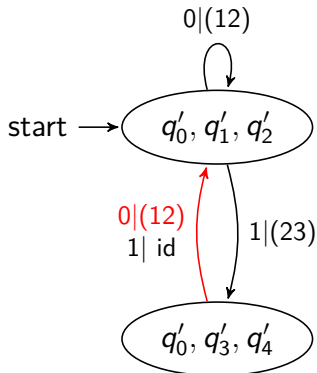
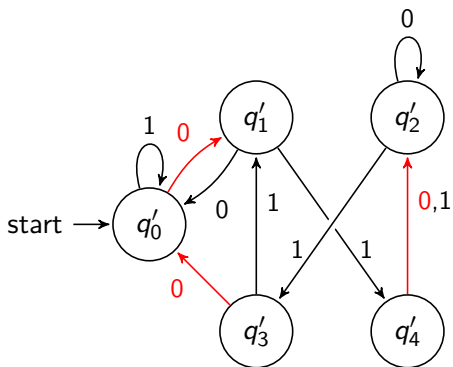
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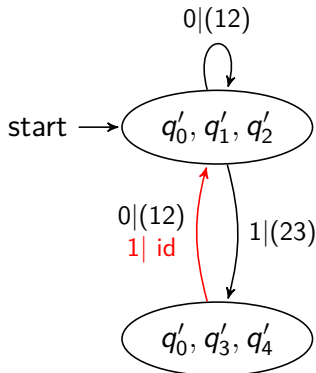
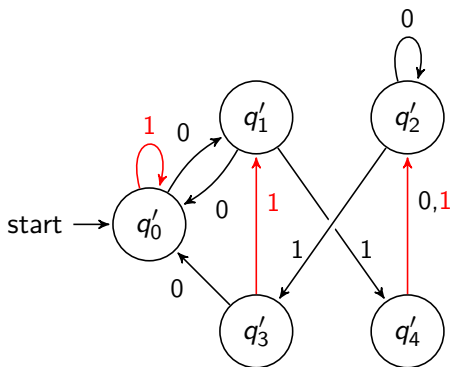
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# Techniques

Use and adopt a framework of Mauduit and Rivat developed for the Rudin-Shapiro sequence.

- Carry Property: The contribution of high and low digits is „independent“.
- Fourier Property:  
We say that  $U$  has the **Fourier property** if there exists  $\eta > 0$  and  $c$  such that for all  $\lambda, \alpha$  and  $t$

$$\left\| \frac{1}{k^\lambda} \sum_{m < k^\lambda} U(mk^\alpha) e(mt) \right\| \leq ck^{-\eta\lambda}.$$

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# Zeckendorf Representation

## Fibonacci numbers

$F_0 = 0, F_1 = 1$  and  $F_{k+2} = F_{k+1} + F_k$  for  $k \geq 0$ .

$$F_n = \frac{\varphi^n - (-\varphi)^{-n}}{\sqrt{5}},$$

where,  $\varphi$  is the golden ratio.

## Zeckendorf Representation

Every positive integer  $n$  admits a unique representation

$$n = \sum_{i \geq 2} \varepsilon_i(n) F_i,$$

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We denote by

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the Zeckendorf sum-of-digits function.

We note that  $s_{\varphi}(n)$  is the least  $k$  such that  $n$  is the sum of  $k$  Fibonacci numbers.

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# Main Result

Theorem (Drmota, M., Spiegelhofer, 2017)

Let  $s_\varphi(n)$  be the Zeckendorf sum-of-digits function and  $m(n)$  a bounded multiplicative function. Then we have

$$\sum_{n < N} (-1)^{s_\varphi(n)} m(n) = o(N) \quad (N \rightarrow \infty).$$

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## Definition

Let  $E$  be a finite set and  $\sigma$  a  $k$ -uniform morphism such that  $\sigma(E) \subseteq E^k$ . Then if  $\mathbf{w}$  is a fixed point of  $\sigma$ , i.e.  $\sigma(\mathbf{w}) = \mathbf{w}$ , then  $\mathbf{w}$  is a  $k$ -automatic sequence.

## Example (Thue-Morse)

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## A Morphism

$$a \mapsto ab$$

$$b \mapsto c$$

$$c \mapsto cd$$

$$d \mapsto a.$$

This gives the sequence  $(-1)^{s_\varphi(n)}$  under the coding  
 $\tau(a) = \tau(d) = 1, \tau(b) = \tau(c) = -1$ .

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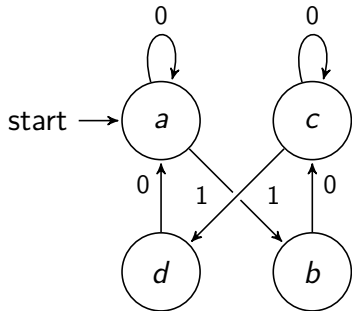
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# A DFAO

We use as input the Zeckendorf representation of  $n$ , i.e.  $\varepsilon_k(n), \dots, \varepsilon_0(n)$ :



# Plan of the Proof

- Use the Kátai Criterion to reduce the problem to

$$\sum_{n \leq N} (-1)^{s_\varphi(pn) + s_\varphi(qn)} = o(N),$$

for all different primes  $p, q$ .

- Use a generating function approach and “quasi-additivity” of  $(-1)^{s_\varphi(pn) + s_\varphi(qn)}$  to reduce this to:

$$s_\varphi(pn_0) \not\equiv s_\varphi(qn_0) \pmod{2} \tag{1}$$

for some  $n_0$ .

- Show (1).

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- Use a generating function approach and “quasi-additivity” of  $(-1)^{s_\varphi(pn) + s_\varphi(qn)}$  to reduce this to:

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# Plan of the Proof

- Use the Kátai Criterion to reduce the problem to

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Suppose that  $(x_n)$  is a bounded complex valued sequence with values in a finite set and that for every pair  $(p, q)$  of different prime numbers we have

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## Definition

We say that  $n_1$  and  $n_2$  are  $r$ -separated at position  $k$  if  $\varepsilon_i(n_1) = 0$  for  $i \geq k - r$  and  $\varepsilon_i(n_2) = 0$  for  $i \leq k + r$ .

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We call a function  $f(n)$  quasi-additive (with respect to the Zeckendorf expansion) if there exists  $r \geq 0$  such that



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for all integers  $n_1, n_2$  that are  $r$  separated.

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## Lemma

Let  $q > p \geq 2$  and  $f(n) = s_\varphi(pn) + s_\varphi(qn)$ . Then  $f(n)$  is quasi-additive with respect to the Zeckendorf expansion.

Proof (Sketch):

It suffices to work with  $s_\varphi(mn)$  as the sum of quasi-additive functions is again quasi-additive.

Choose  $r$  such that  $\varphi^{r-1} < m$ .

$$n_1 < F_{k-r} \Rightarrow mn_1 < F_k.$$

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# Generating Functions Approach

Let  $f$  be a quasi-additive function and

$$H(x, z) := \sum_{k \geq 3} x^k \sum_{F_{k-1} \leq n < F_k} z^{f(n)}.$$

Note that

$$[x^k]H(x, -1) = \sum_{F_{k-1} \leq n < F_k} (-1)^{s_\varphi(pn) + s_\varphi(qn)}.$$

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$$\begin{aligned} H(x, z) &= \frac{1}{1-x} \frac{1}{1 - B(x, z) \frac{x^{2r+1}}{1-x}} B'(x, z) \\ &= \frac{B'(x, z)}{1 - x - x^{2r+1} B(x, z)}. \end{aligned}$$

The dominant singularity of  $H(x, 1)$  is at  $x_0 = \frac{1}{\varphi}$ .

This is due to the fact that  $x = x_0$  is a solution for

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- 1 Other base:  $G_0 = 0, G_1 = 1, G_{k+1} = aG_k + G_{k-1}$ .
- 2 More general bases: Ostrowski numeration.
- 3 Replace the sum-of-digits function by a block-additive function.
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