

# Fully Packed Loop Configurations in a triangle

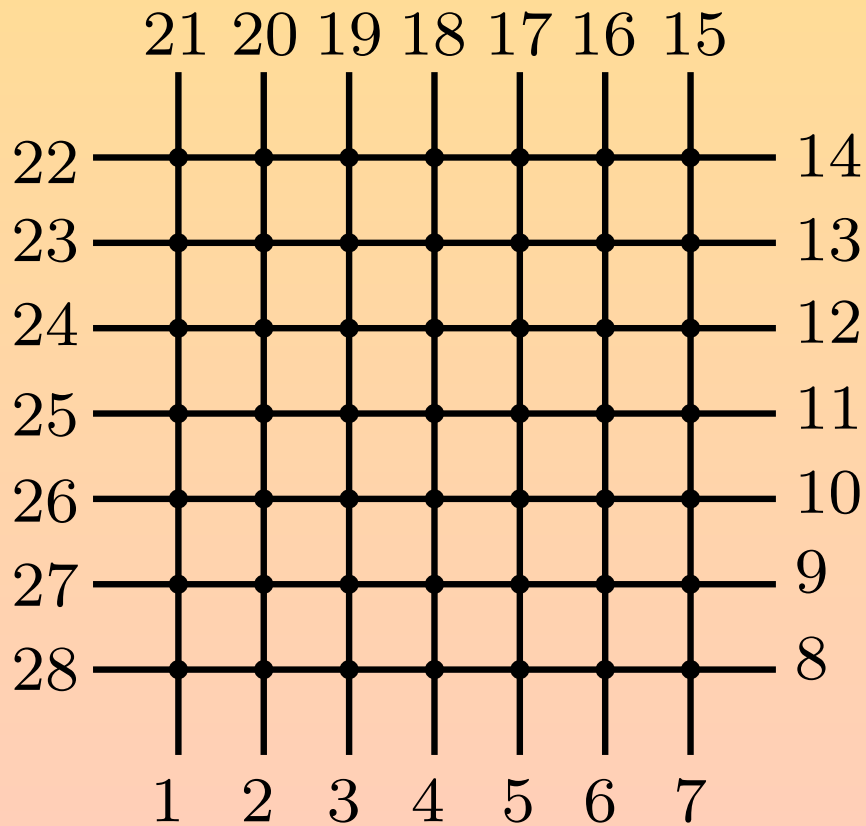
Philippe Nadeau

Faculty of Mathematics, University of Vienna

Paris, IHP, October 2009

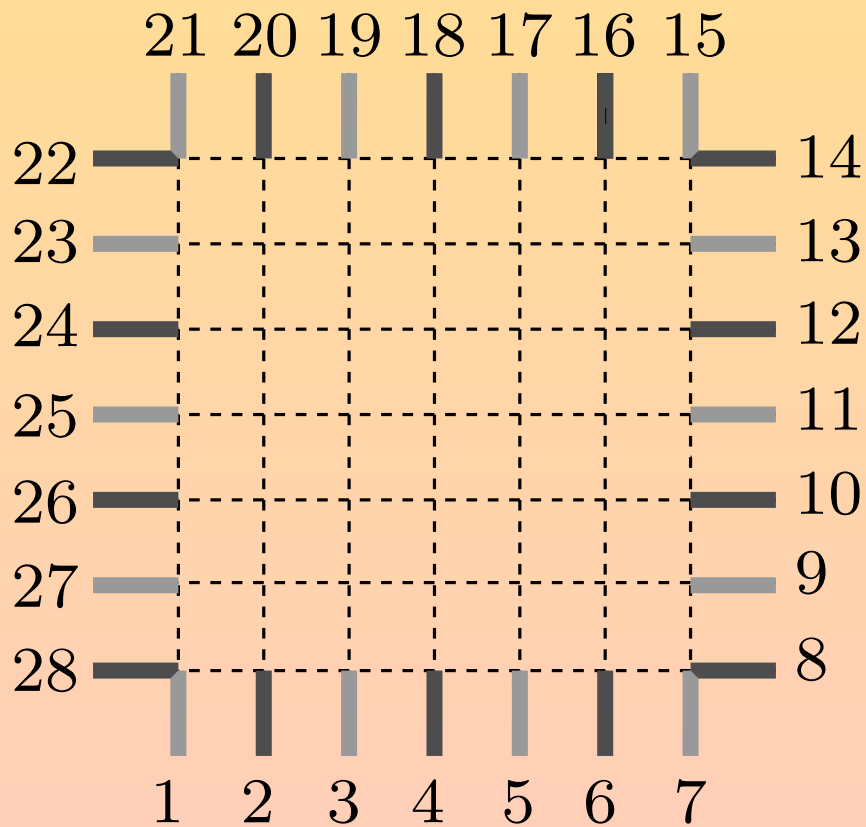
# FPL configurations : Definition

Start with the **square grid**  $G_n$  with  $n^2$  vertices and  $4n$  external edges. In the example, we have  $n = 7$ .



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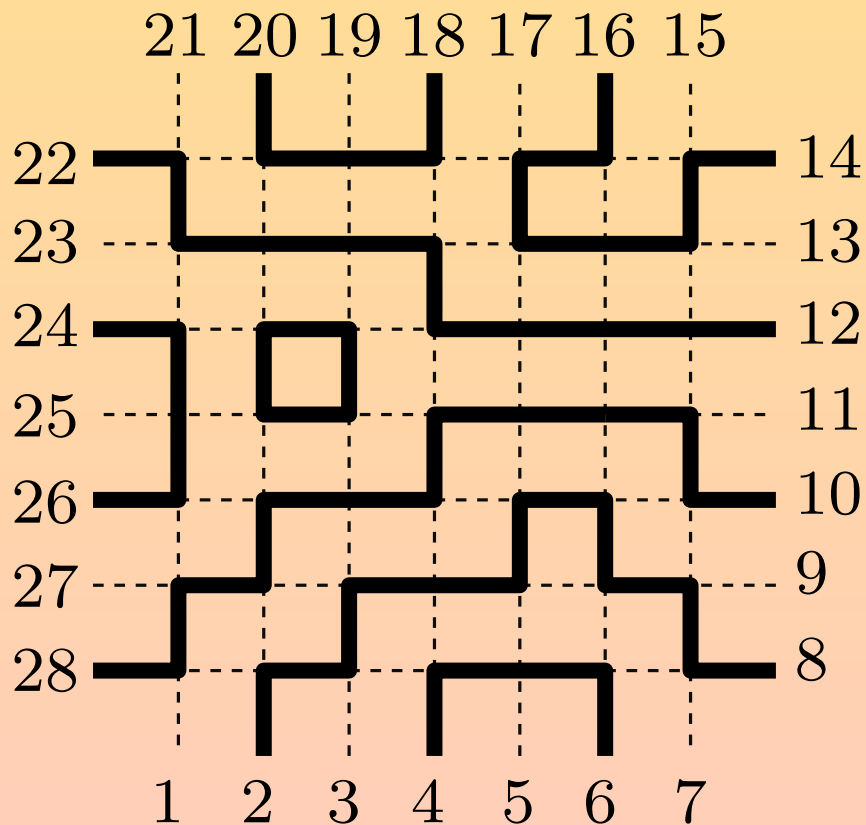


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(1) containing every other external edge, i.e. contains either all odd edges or all even edges.

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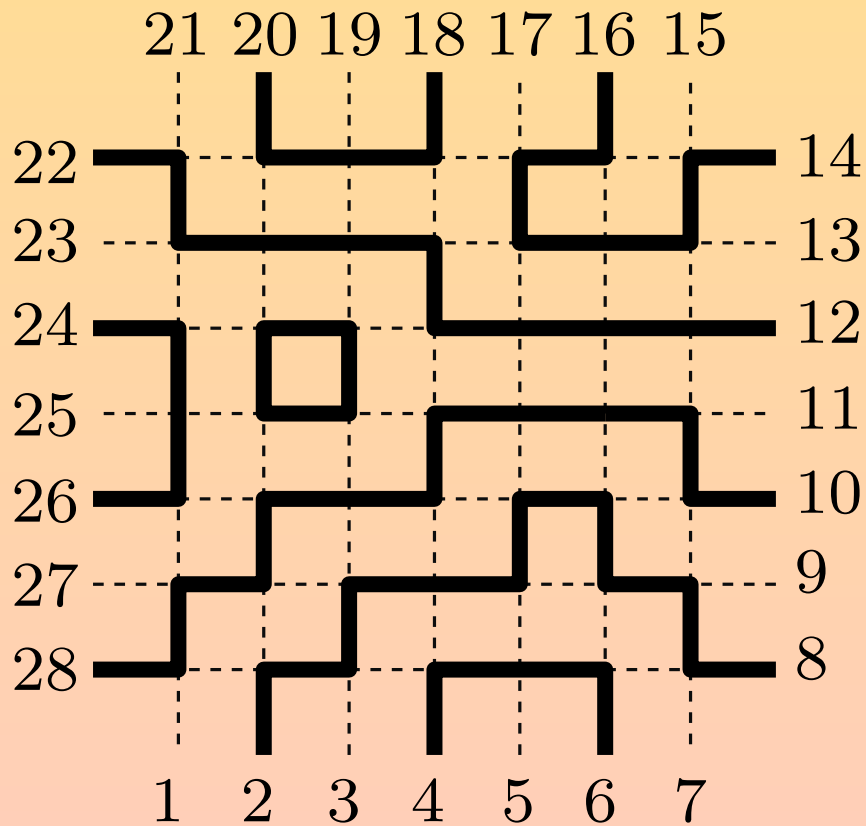


A **FPL configuration of size  $n$**  is a subgraph of the grid  $G_n$

- (1) containing every other external edge ;
- (2) such that around each vertex of  $G_n$ , 2 edges out of 4 are selected.

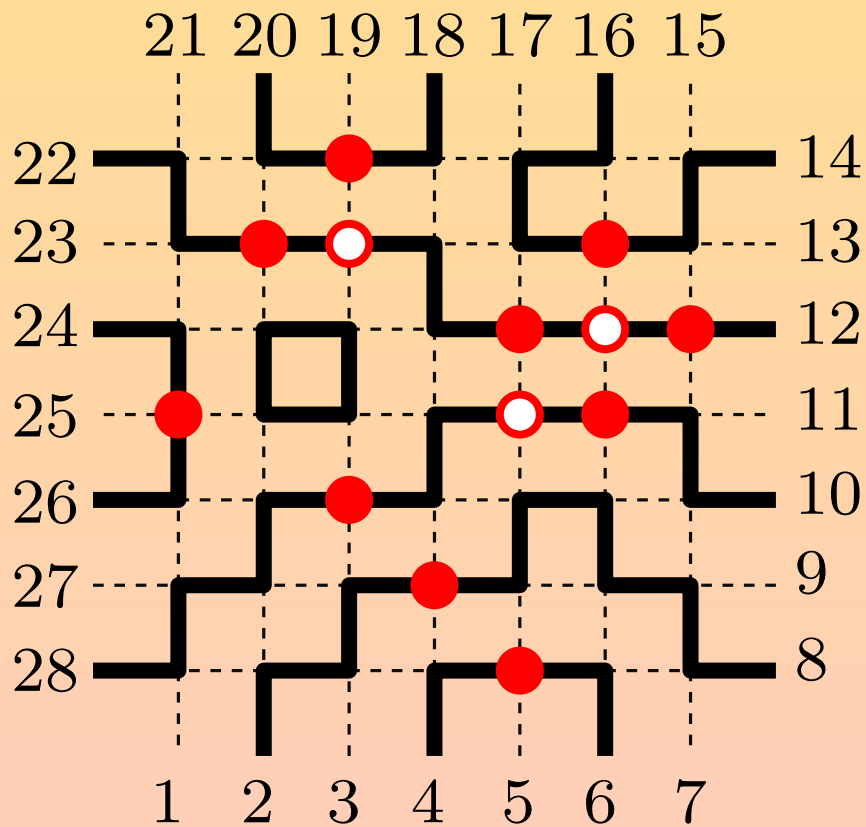
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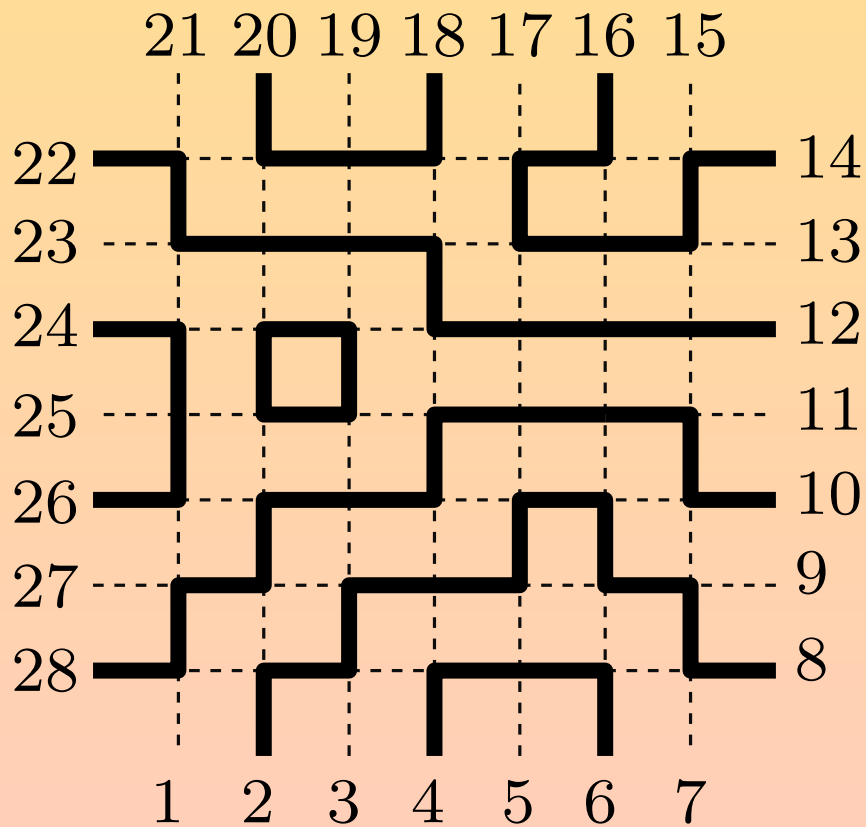
Alternating sign matrices of size  $n$

[ ASM = matrix with coefficients in  $\{1, 0, -1\}$  such that on each row or column  $1$  and  $-1$  alternate, and the sum is  $1$ .]

Here  $1 \rightarrow \bullet$  and  $-1 \rightarrow \circ$

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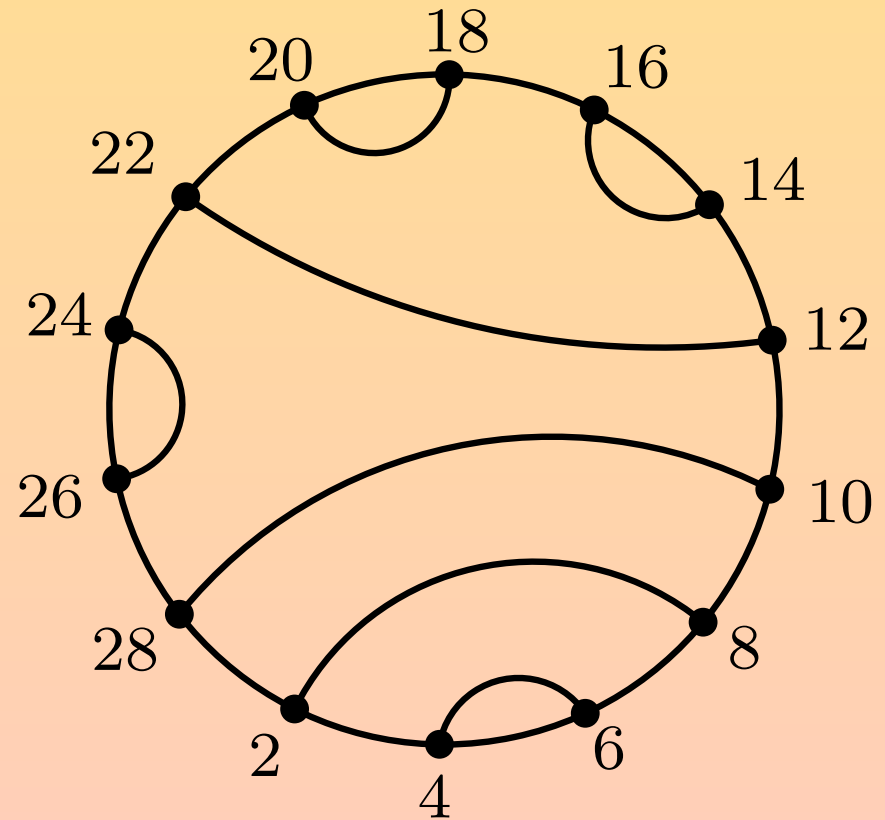
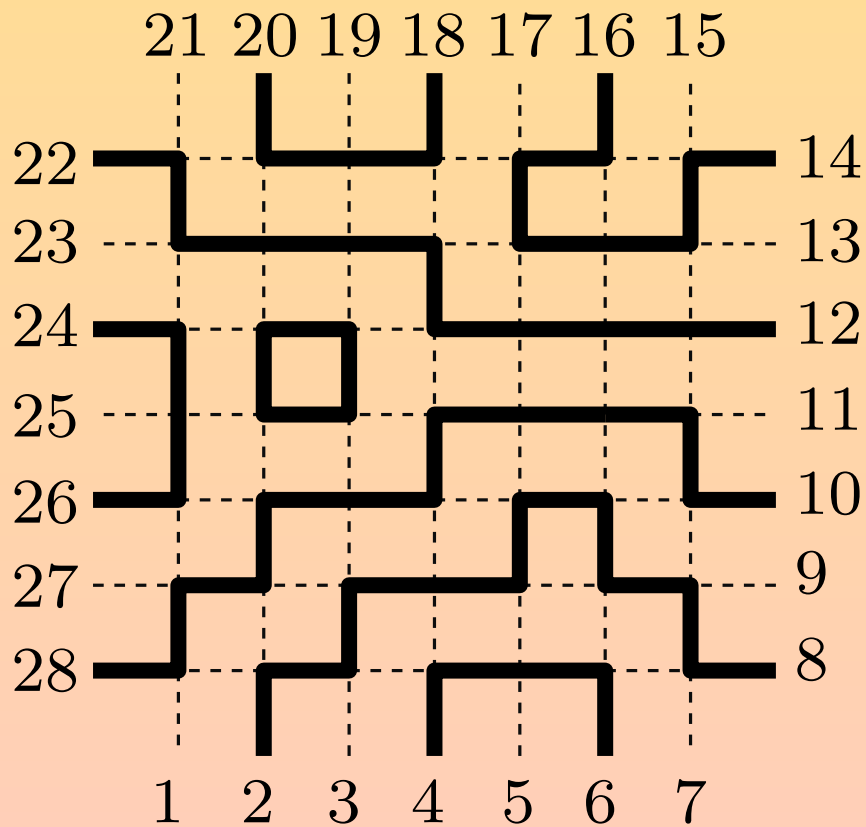
Alternating sign matrices of size  $n$

$$A_n = \prod_{i=0}^{n-1} \frac{(3i+1)!}{(n+i)!}$$

[Zeilberger '96, Kuperberg '96]

# FPL configurations : Refined enumeration

Every FPL configuration determines a **link pattern** on the odd or even external edges of the grid  $G_n$ .





# FPL configurations : Refined enumeration

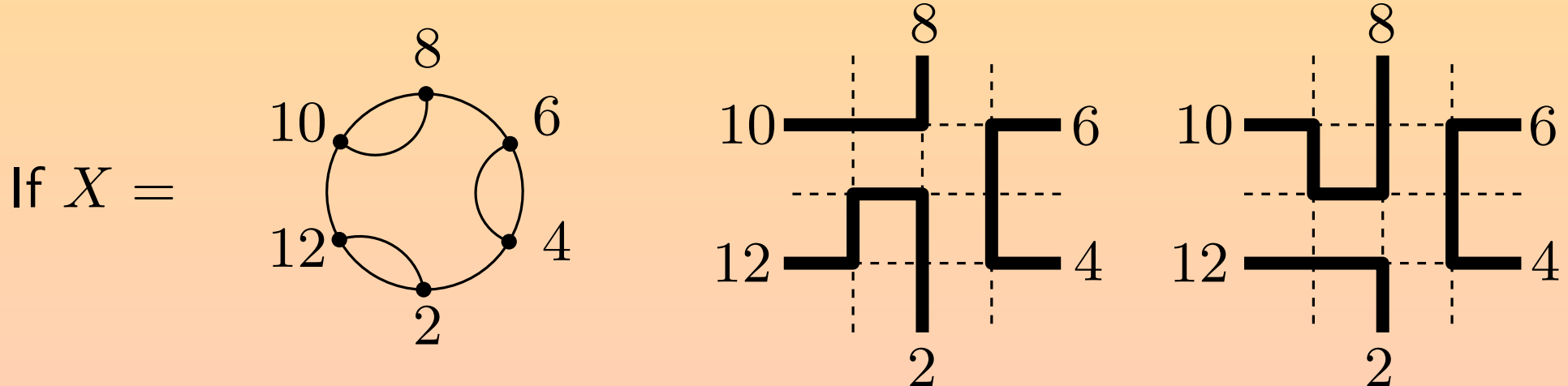
Now if we are given a pairing  $X$  of odd (or even) external edges, our main question will be : **how many FPL configurations respect the link pattern  $X$  ?**

**Definition** We note  $A_X$  the number of FPL configurations which induce the link pattern  $X$ .

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For this link pattern we have  $A_X = 2$ .

# FPL configurations : Refined enumeration

Now given a link pattern  $X$ , let  $X'$  be defined by

$$(i, j) \in X' \Leftrightarrow (i - 1, j - 1) \in X$$

**Theorem [Wieland '00]**

$$A_X = A_{X'}$$

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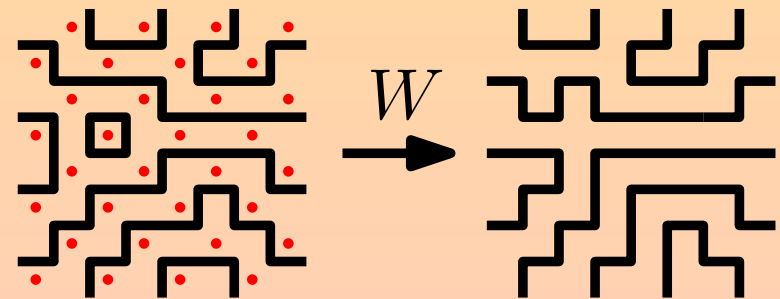
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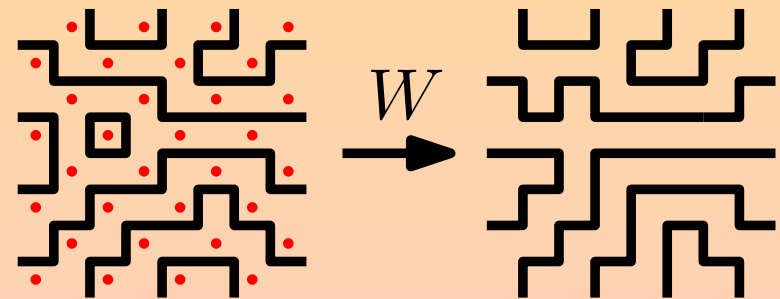
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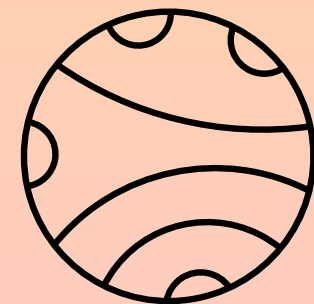
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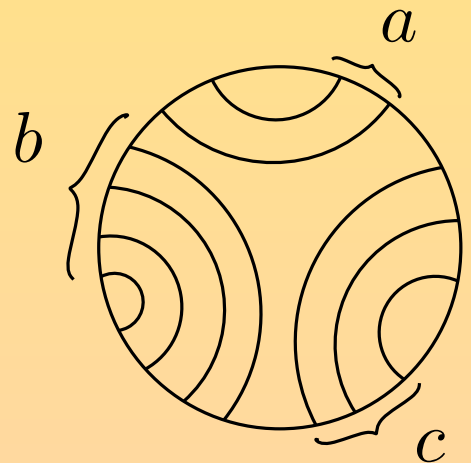


For enumeration purposes, we can then use **unlabeled** link patterns :



# Outline of the talk

Known enumerations for the numbers  $A_X$  are


$$= \prod_{i=1}^a \prod_{j=1}^b \prod_{k=1}^c \frac{i + j + k - 1}{i + j + k - 2}$$


$$= \text{Complicated determinant formulas}$$

+ certain variants of these.

These results are due to Zinn-Justin, Zuber, Di Francesco, Caselli, Krattenthaler,...

# Outline of the talk

For a given link pattern  $X$  of size  $n$ , there exist numerous instances in [Zuber '04] of conjectured identities of the form

$$A_X = \sum_{c_{XX'} \in \mathbb{Z}} c_{XX'} A_{X'} \text{ where } X' \text{ are link patterns of size } n - 1.$$

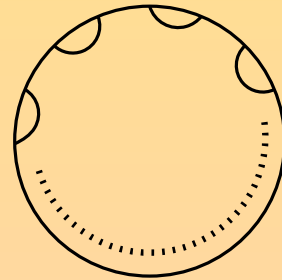
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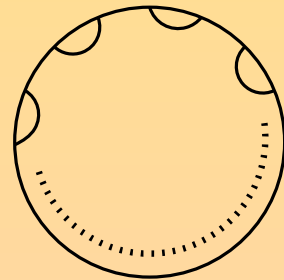
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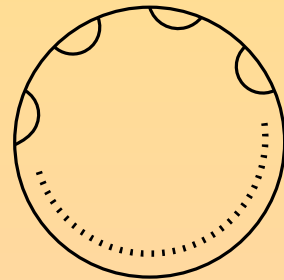
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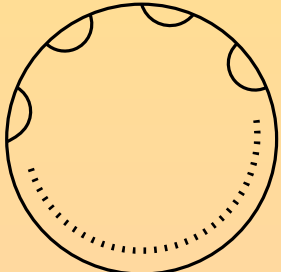
These coefficients are defined with respect to certain [FPL configurations in a triangle](#), and we will focus on enumerating these configurations in certain special cases.

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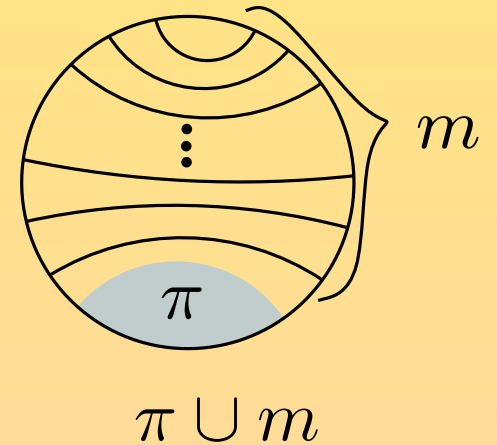
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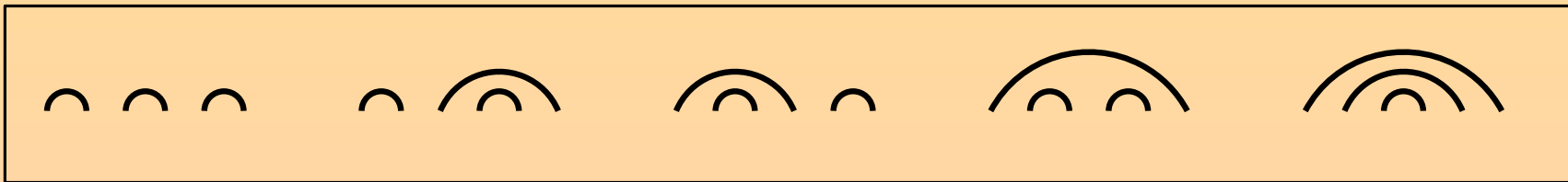
In one such case, we will show that the answer is given by the famous [Littlewood-Richardson coefficients](#).

# Link patterns with nested arches

We consider now integers  $n, m \geq 0$ , and link patterns with  $m$  nested arches, and  $\pi$  is a **noncrossing matching** with  $n$  arches.

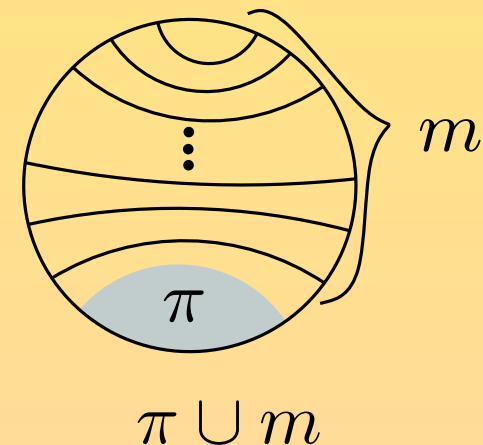


For instance if  $n = 3$ , there are 5 possible  $\pi$  :

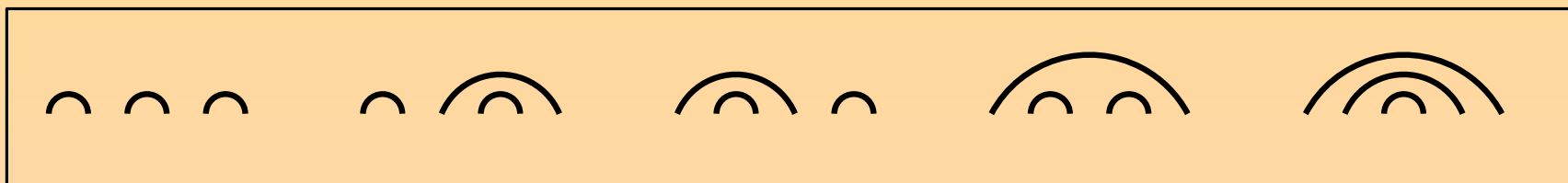


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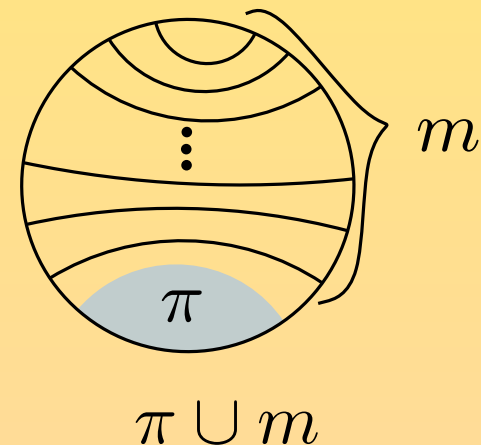
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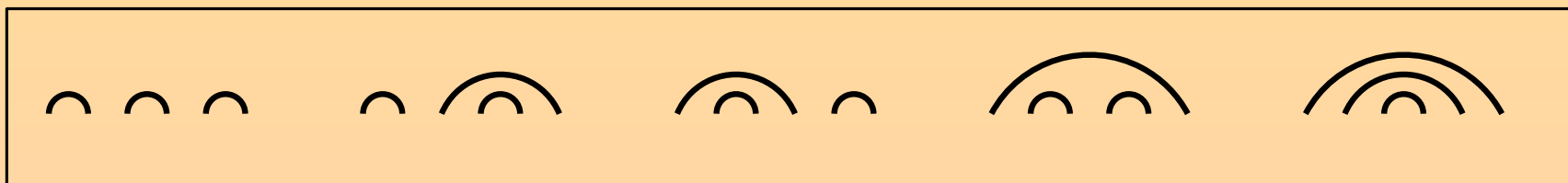
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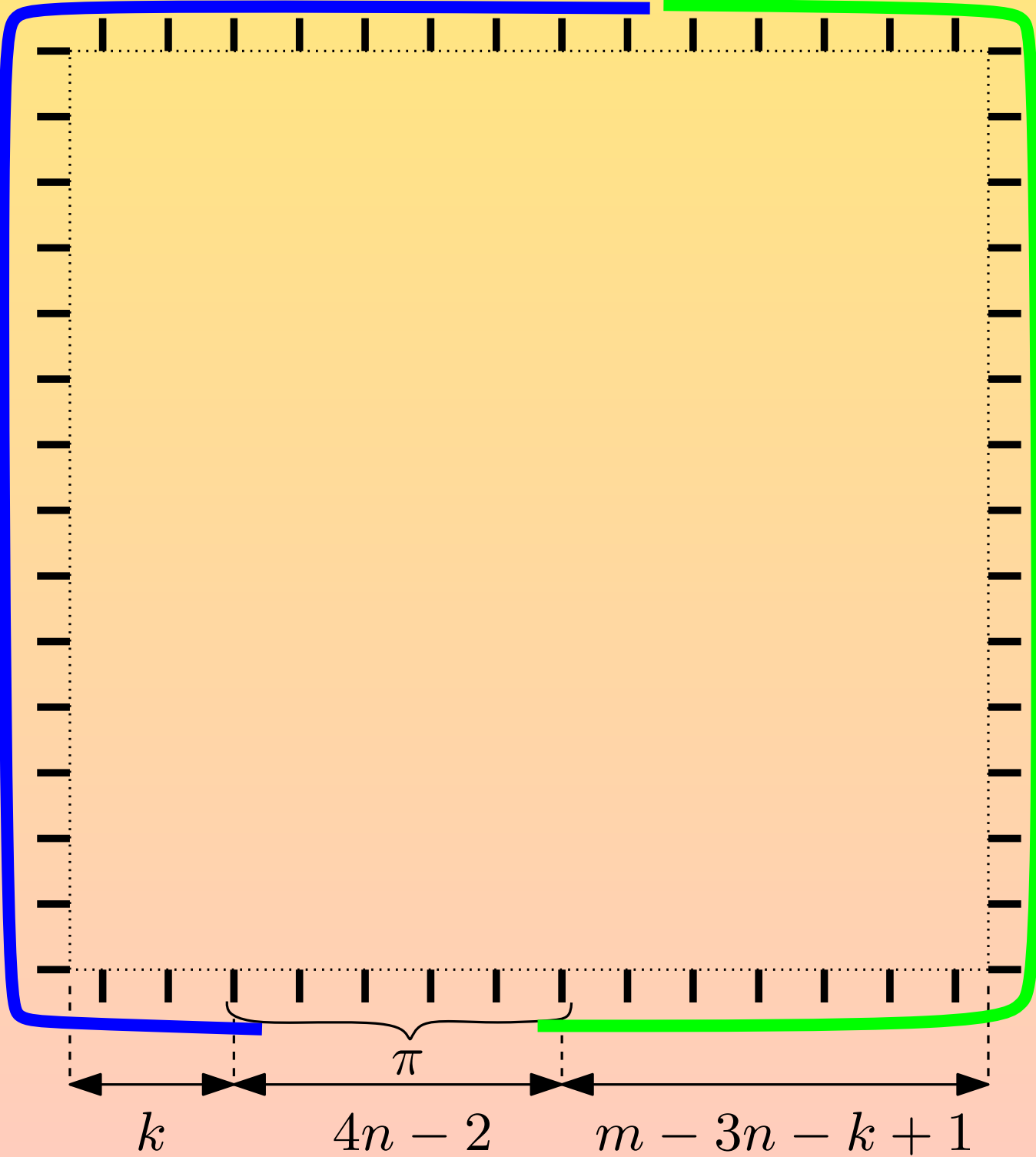
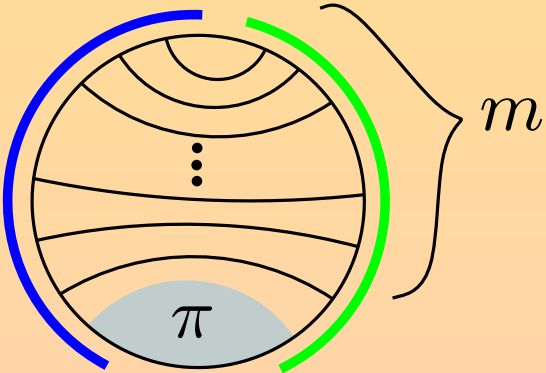


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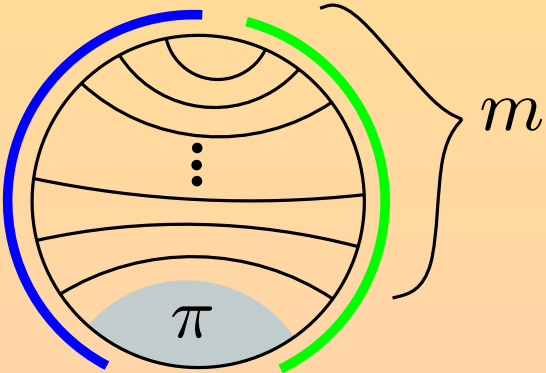
**Theorem [Caselli, Krattenthaler, Lass, N. '05]**

$A_{\pi}(m)$  is a **polynomial** function of  $m$ .

We suppose  $m \geq 3n - 1$ ,  
 and choose  $k$  such that  
 $0 \leq k \leq m - (3n - 1)$ .

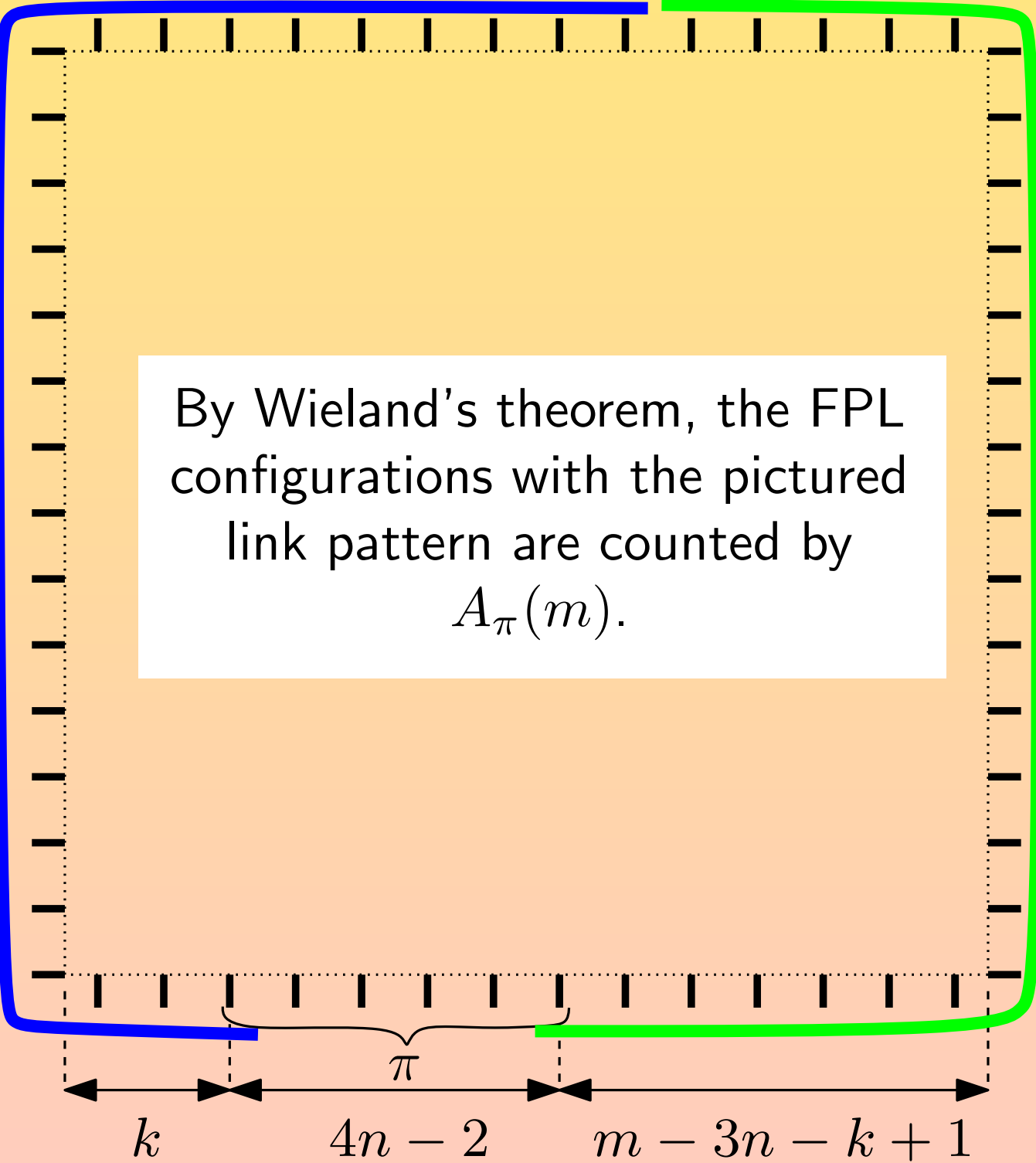


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By Wieland's theorem, the FPL configurations with the pictured link pattern are counted by

$$A_{\pi}(m).$$

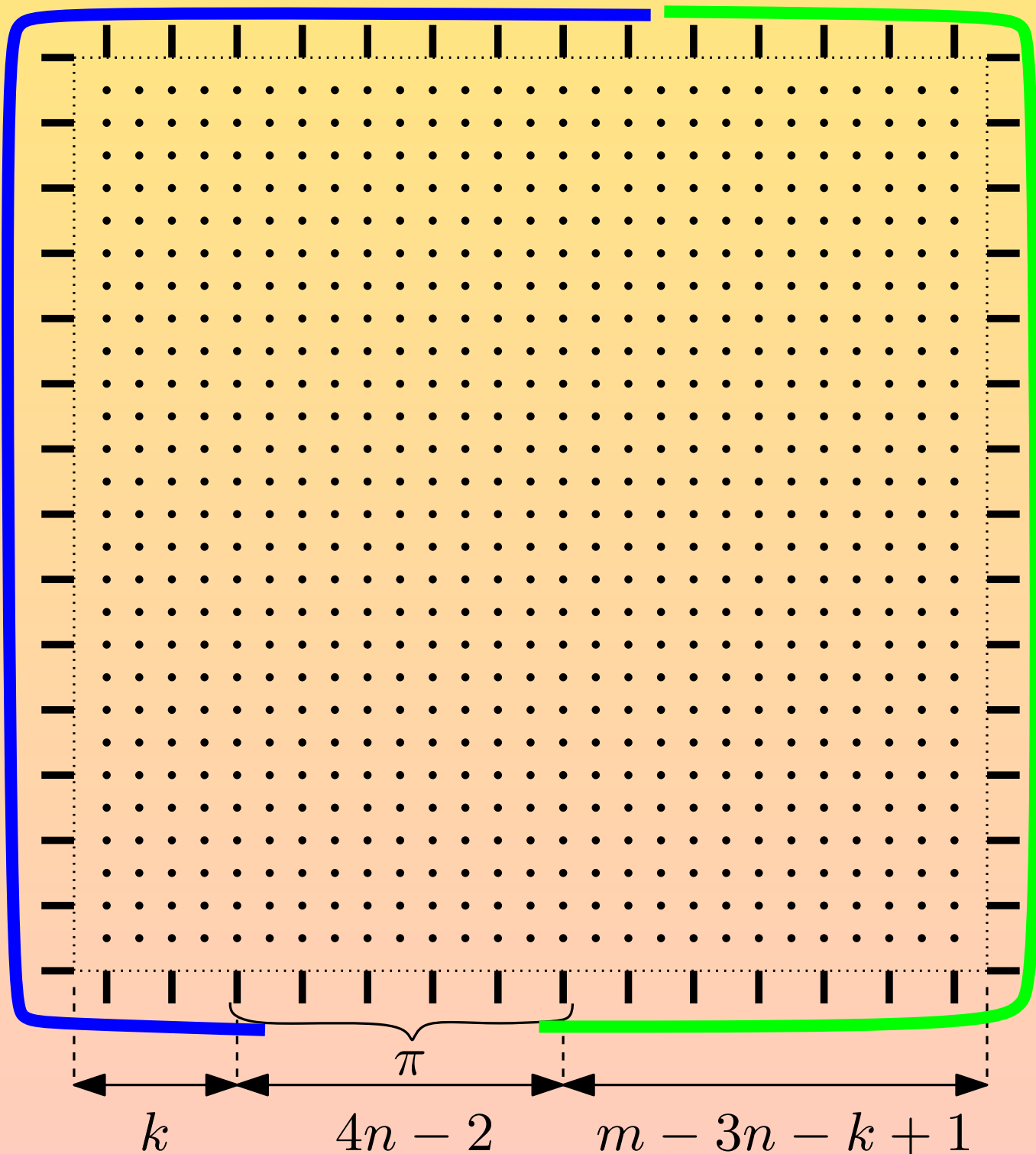




Many edges of the grid belong to every FPL configuration respecting the link pattern.

⇒ "Fixed edges"

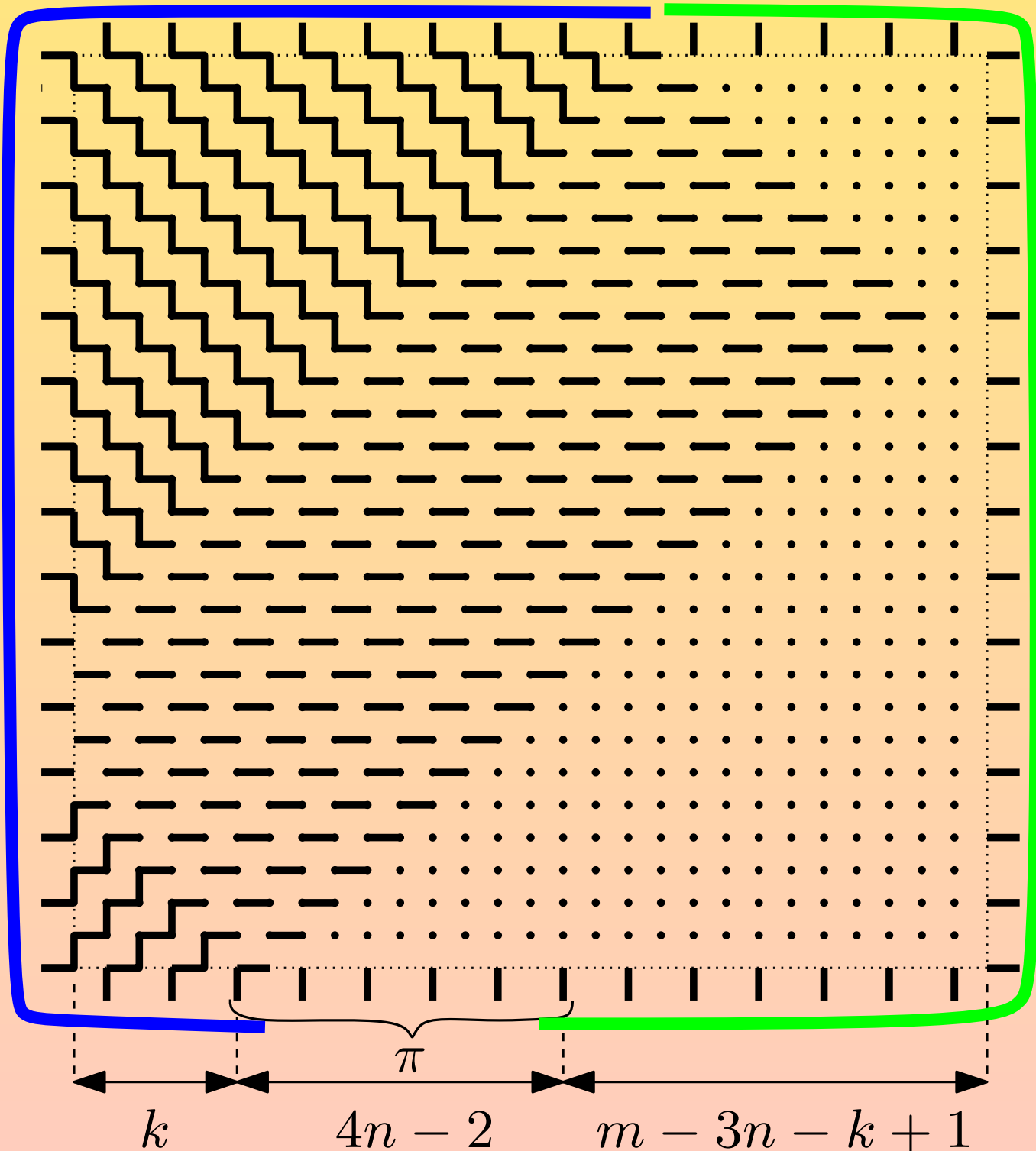
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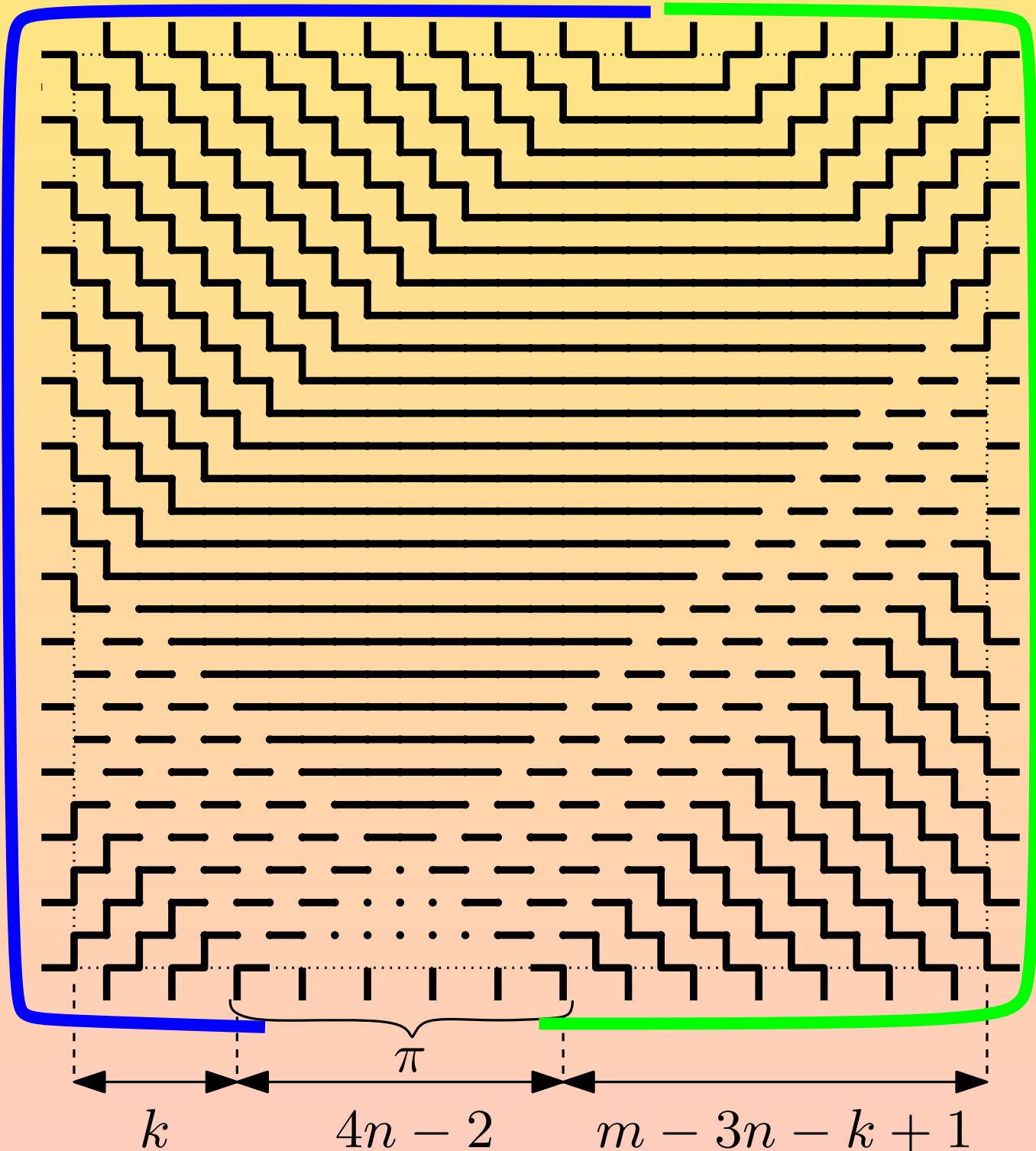
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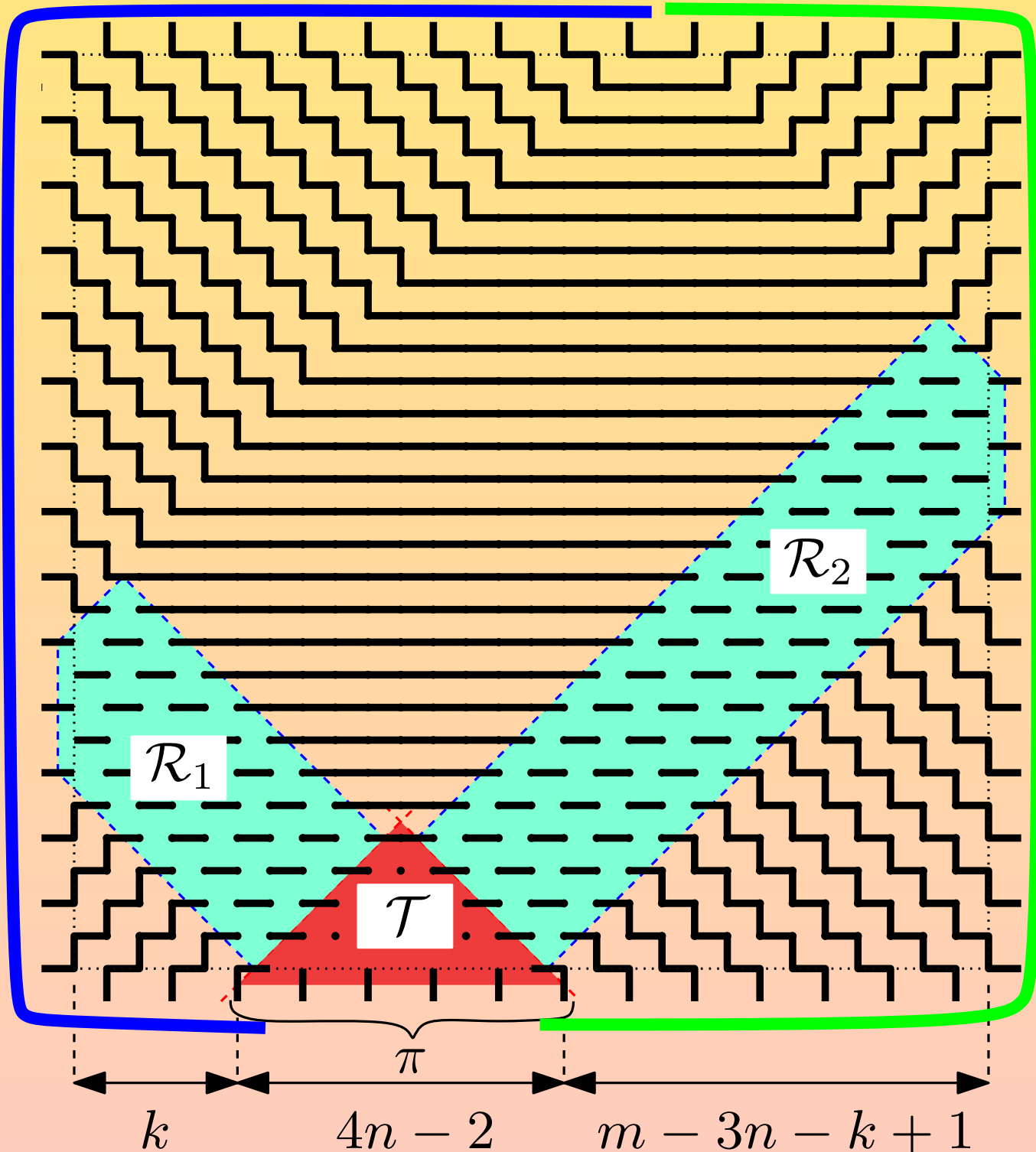
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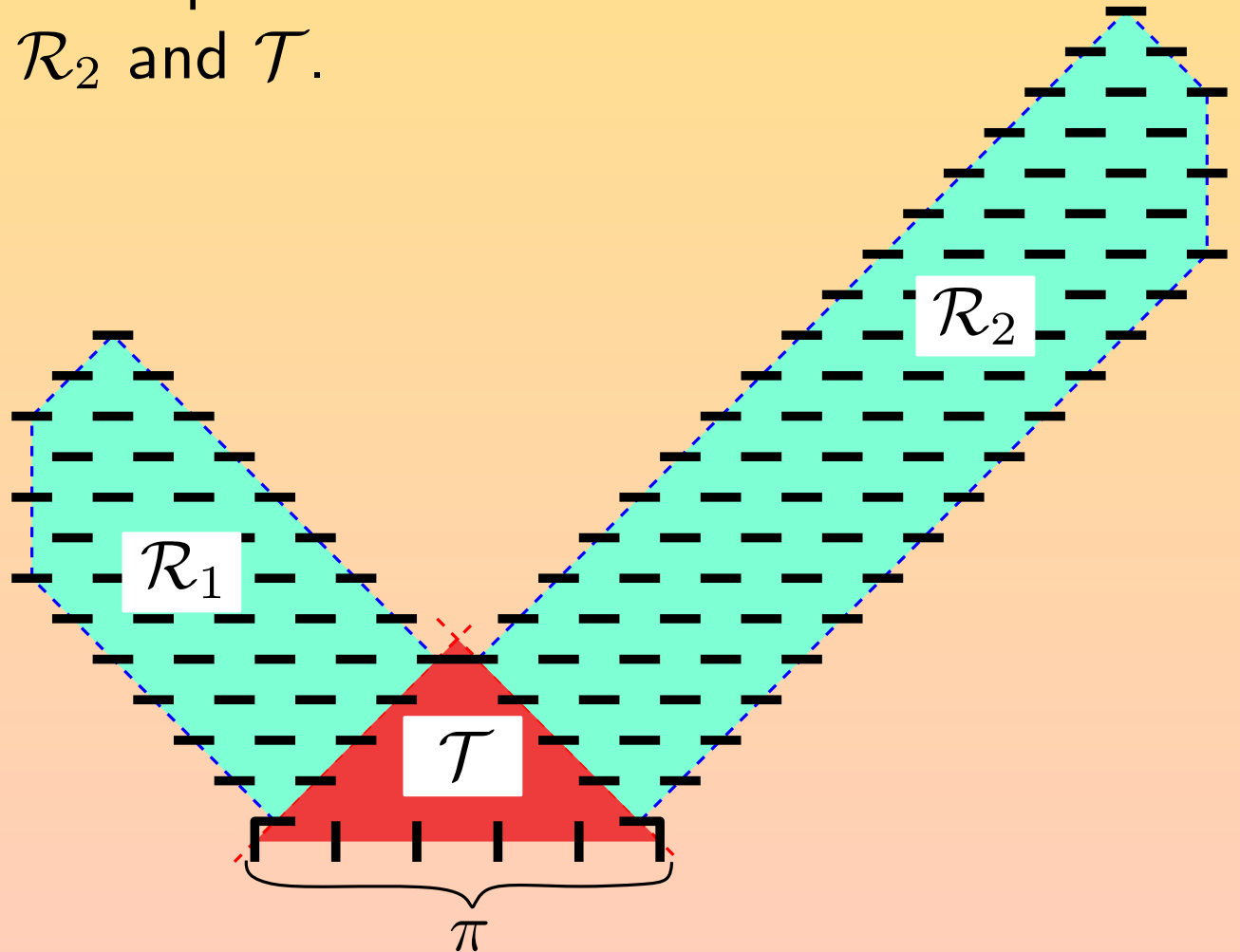
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To compute the numbers  $A_\pi(m)$ , we will count FPL configurations separately in  $\mathcal{R}_1, \mathcal{R}_2, \mathcal{T}$ .

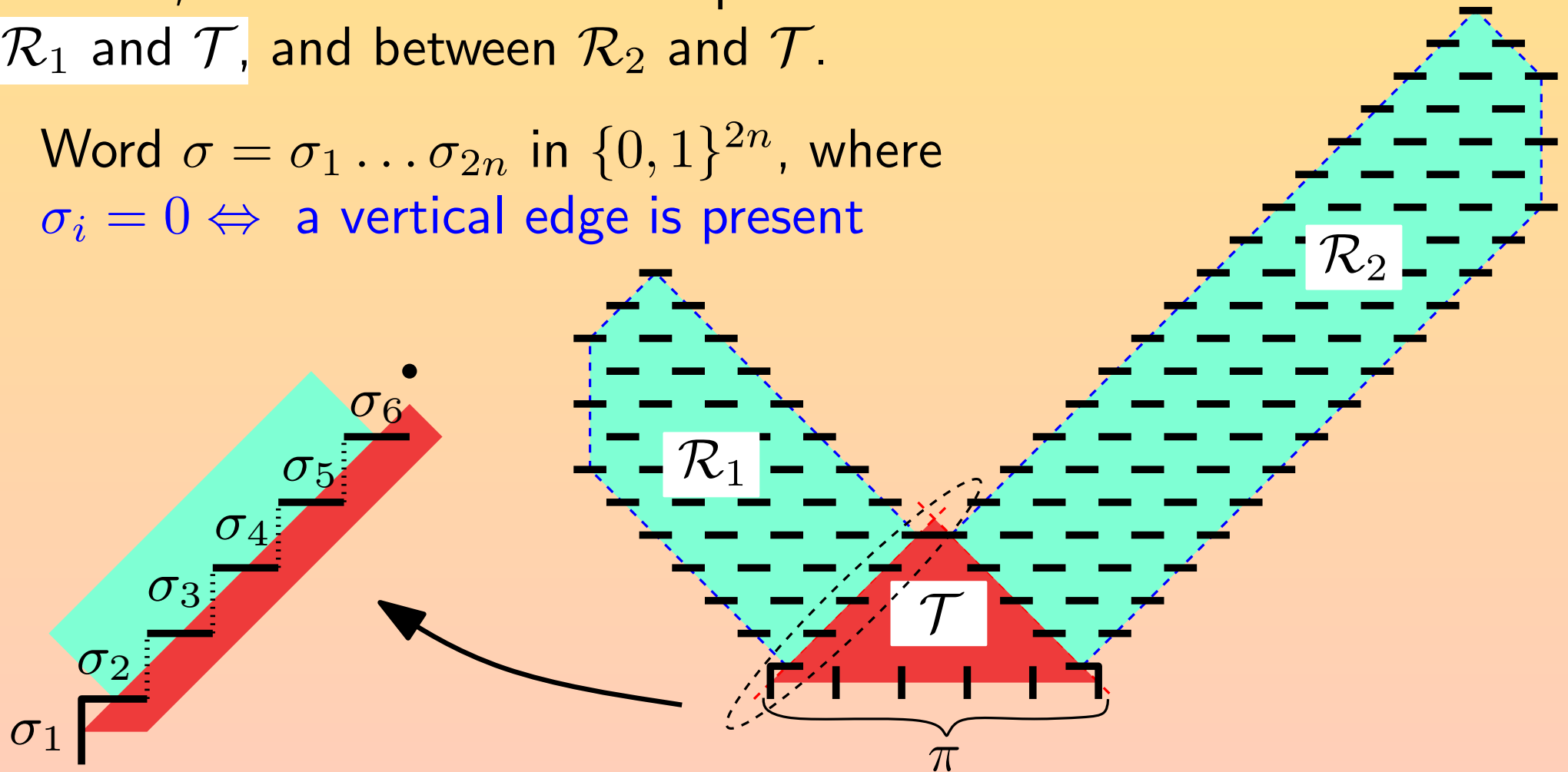
For this, we need to encode the possible boundaries between  $\mathcal{R}_1$  and  $\mathcal{T}$ , and between  $\mathcal{R}_2$  and  $\mathcal{T}$ .



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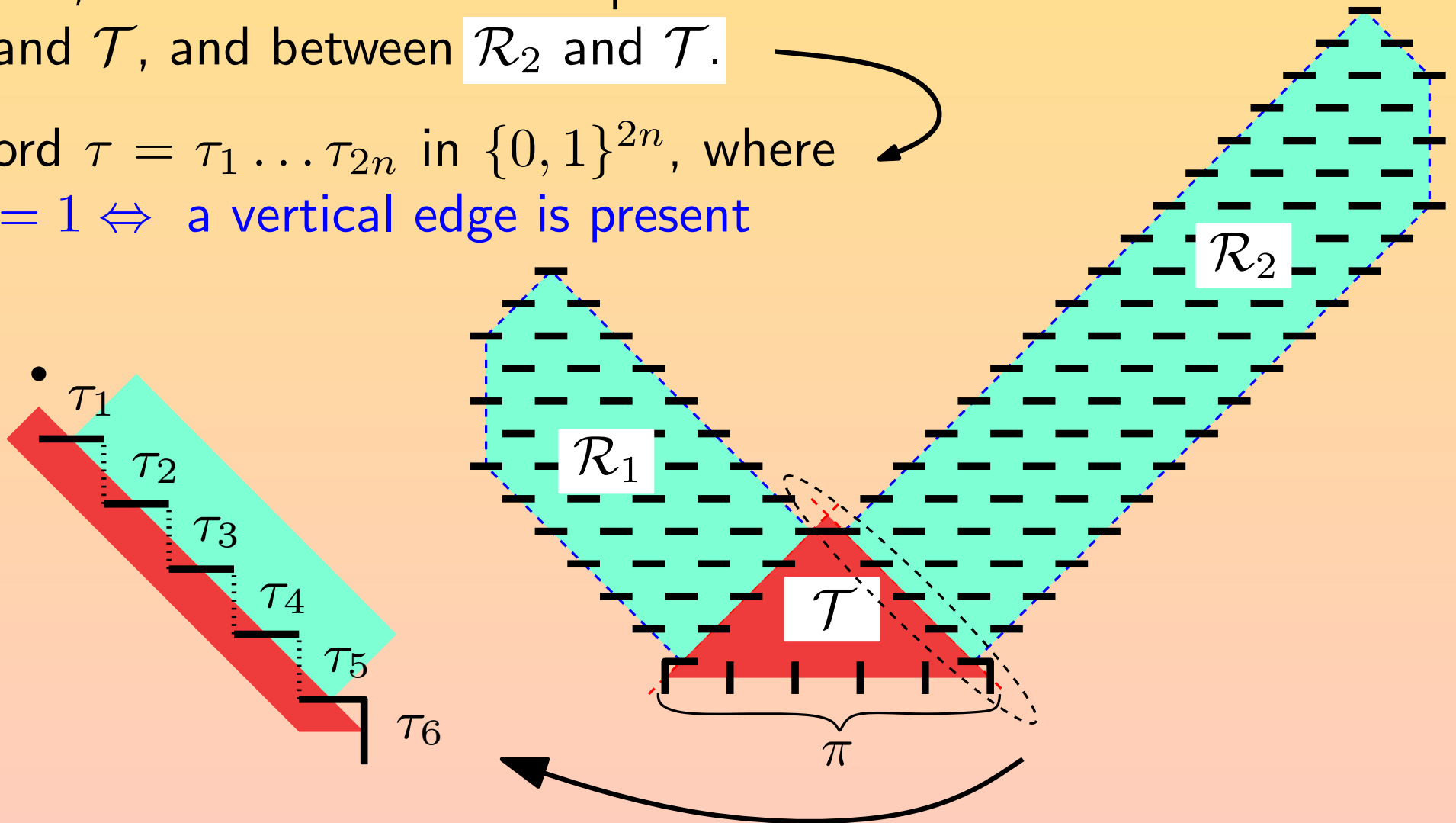
Word  $\sigma = \sigma_1 \dots \sigma_{2n}$  in  $\{0, 1\}^{2n}$ , where  
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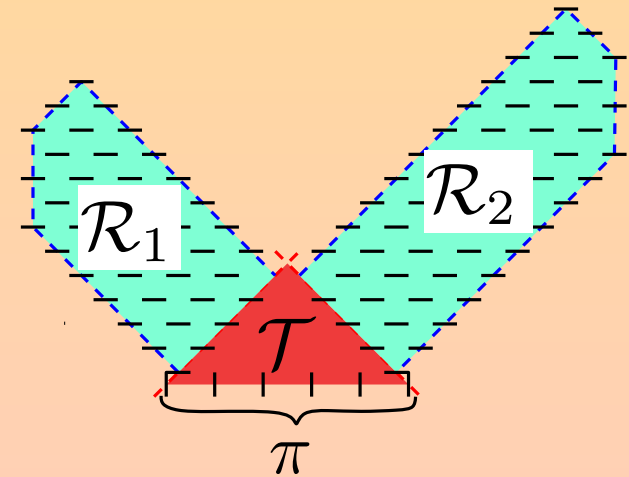
# Putting things together

We can then write, for  $m \geq 3n - 1$  and  $0 \leq k \leq m - (3n - 1)$

$$A_\pi(m) = \sum_{\sigma, \tau} |\mathcal{R}_1(\sigma, k)| \times t_{\sigma, \tau}^\pi \times |\mathcal{R}_2(\tau, m - 3n - k + 1)|$$

where

- $\sigma, \tau$  are words of length  $2n$  on  $\{0, 1\}$  ;
- $\mathcal{R}_1(\sigma, \cdot), \mathcal{R}_2(\tau, \cdot)$  are the sets of FPL configurations in the regions  $\mathcal{R}_1$  and  $\mathcal{R}_2$  with boundaries  $\sigma, \tau$  respectively ;
- $t_{\sigma, \tau}^\pi$  is the number of FPL configurations in the triangle  $\mathcal{T}$  with boundary data  $\{\sigma, \pi, \tau\}$ .







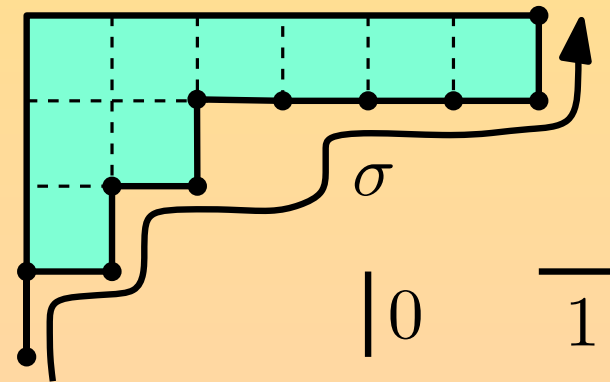
# Words and Shapes

Let  $\sigma = \sigma_1 \dots \sigma_p$  be a word in  $\{0, 1\}^p$ ; we write  $|\sigma| := p$ .

We will identify words and Ferrers shapes in a box.

$$\sigma = 0101011110$$

$$|\sigma| = 10, |\sigma|_0 = 4, |\sigma|_1 = 6$$



$$\ell(\sigma) = 9$$

**Length**  $\ell(\sigma)$  := the number of boxes in the diagram  $\sigma$ .

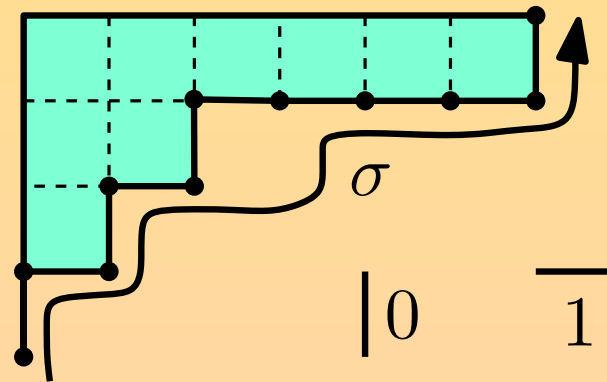
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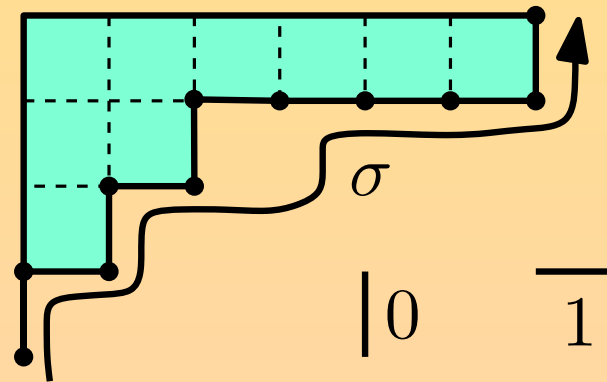
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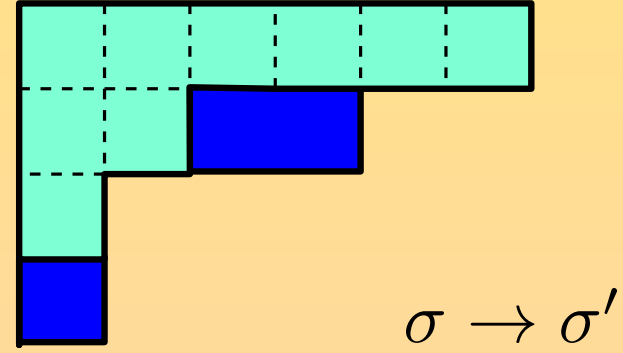
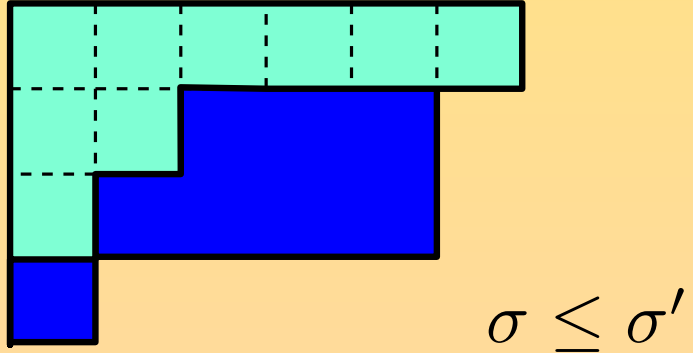
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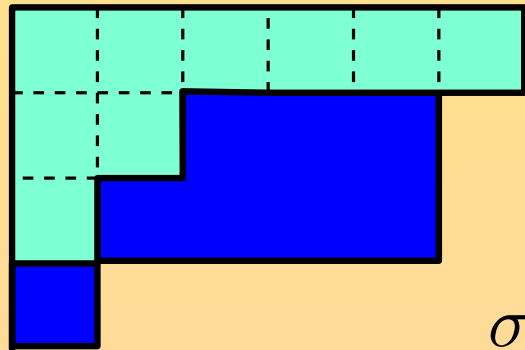
For two words  $\sigma, \sigma'$  with  $|\sigma|_0 = |\sigma'|_0$  and  $|\sigma|_1 = |\sigma'|_1$  we define :

- $\sigma \leq \sigma'$  if, as shapes,  $\sigma$  is **included** in  $\sigma'$ .
- $\sigma \rightarrow \sigma'$  if  $\sigma \leq \sigma'$ , and  $\sigma'$  has at most one more box in each column;  $\sigma, \sigma'$  form a **horizontal strip**.

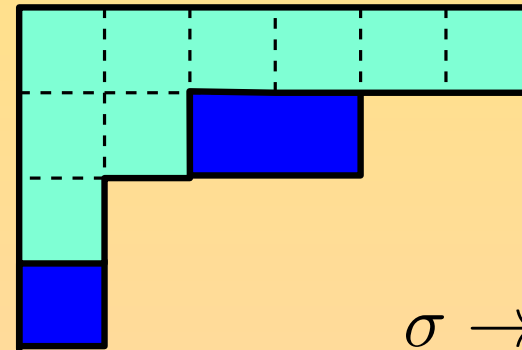
# Words and Shapes



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## Definition

A **semi standard Young tableau** of shape  $\sigma$  and entries bounded by  $N$  is a filling of the shape  $\sigma$  by integers in  $\{1, \dots, N\}$  such that entries are strictly increasing in columns and weakly increasing in rows.

Such a tableau can be equivalently defined by a sequence of shapes

$$\emptyset = \sigma_0 \rightarrow \sigma_1 \rightarrow \dots \rightarrow \sigma_N = \sigma$$

# Words and Shapes

Given a box  $u$  in a Ferrers diagram, in the  $i$ th row from the top and  $j$ th column from the left, we define

- the **content**  $c(u) := j - i$ ;
- the **hook-length**  $h(u)$  as the number of boxes below it, or to its right, including the  $u$  itself.

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-1	0	

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## Theorem [Stanley]

The number of semistandard Young tableaux of shape  $\lambda$  and entries bounded by  $N$  is given by

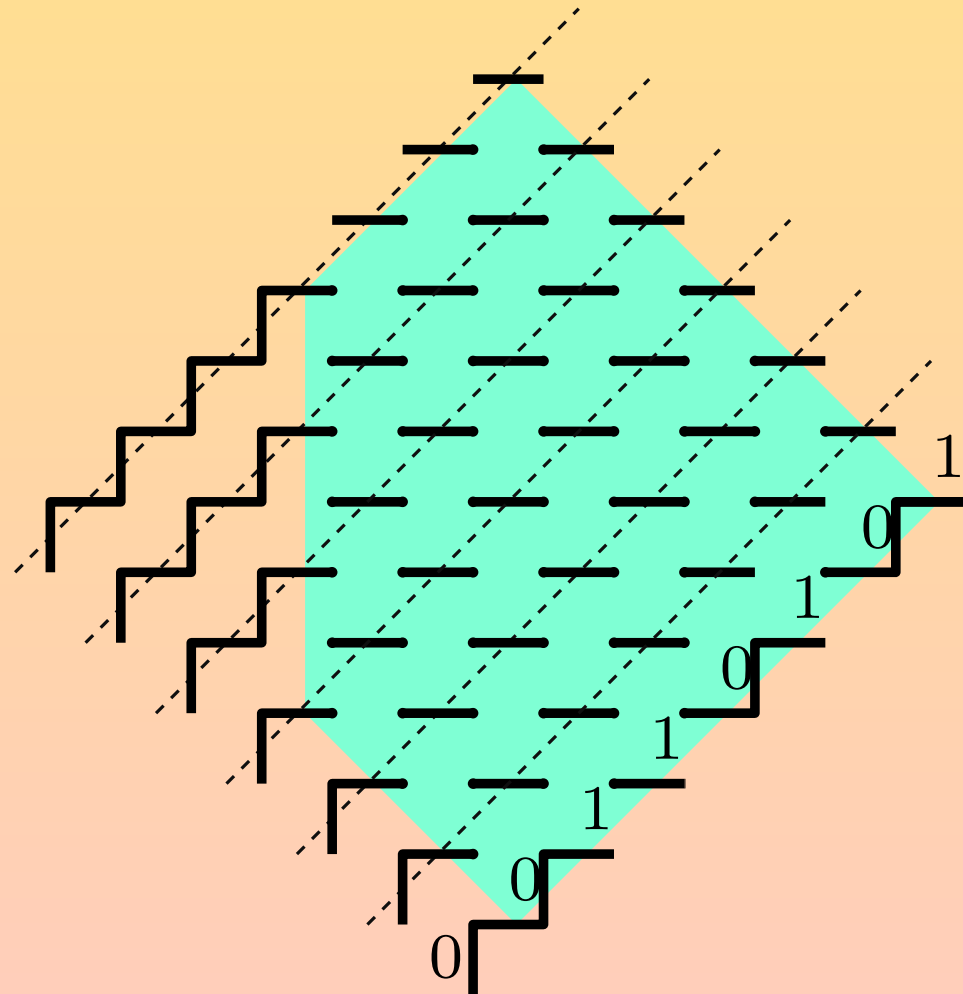
$$SSYT(\lambda, N) = \prod_{u \in \lambda} \frac{N + c(u)}{h(u)}$$

Polynomial of with leading term  $\frac{1}{h(\lambda)} N^{\ell(\lambda)}$



# Regions $\mathcal{R}_1$ and $\mathcal{R}_2$

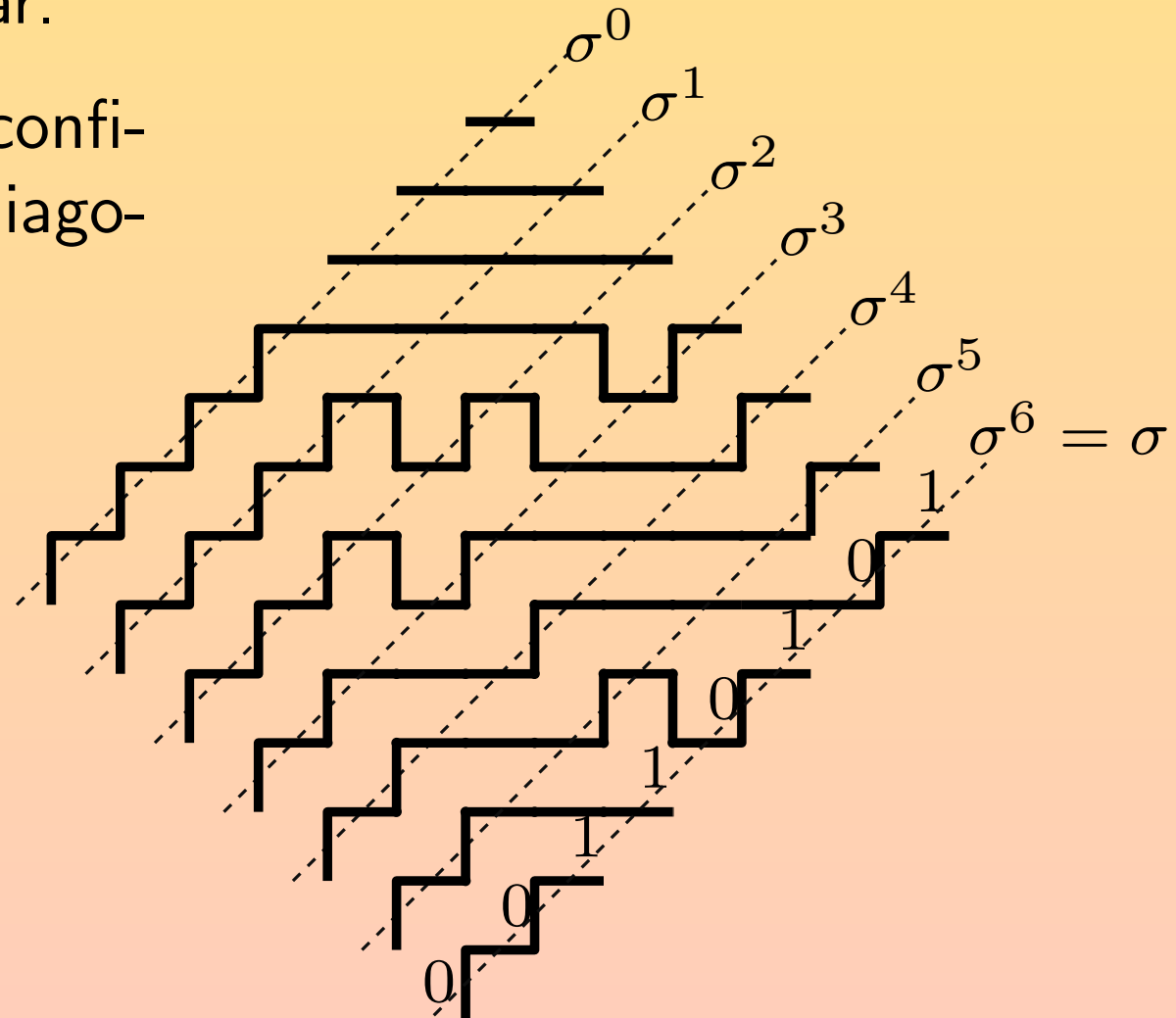
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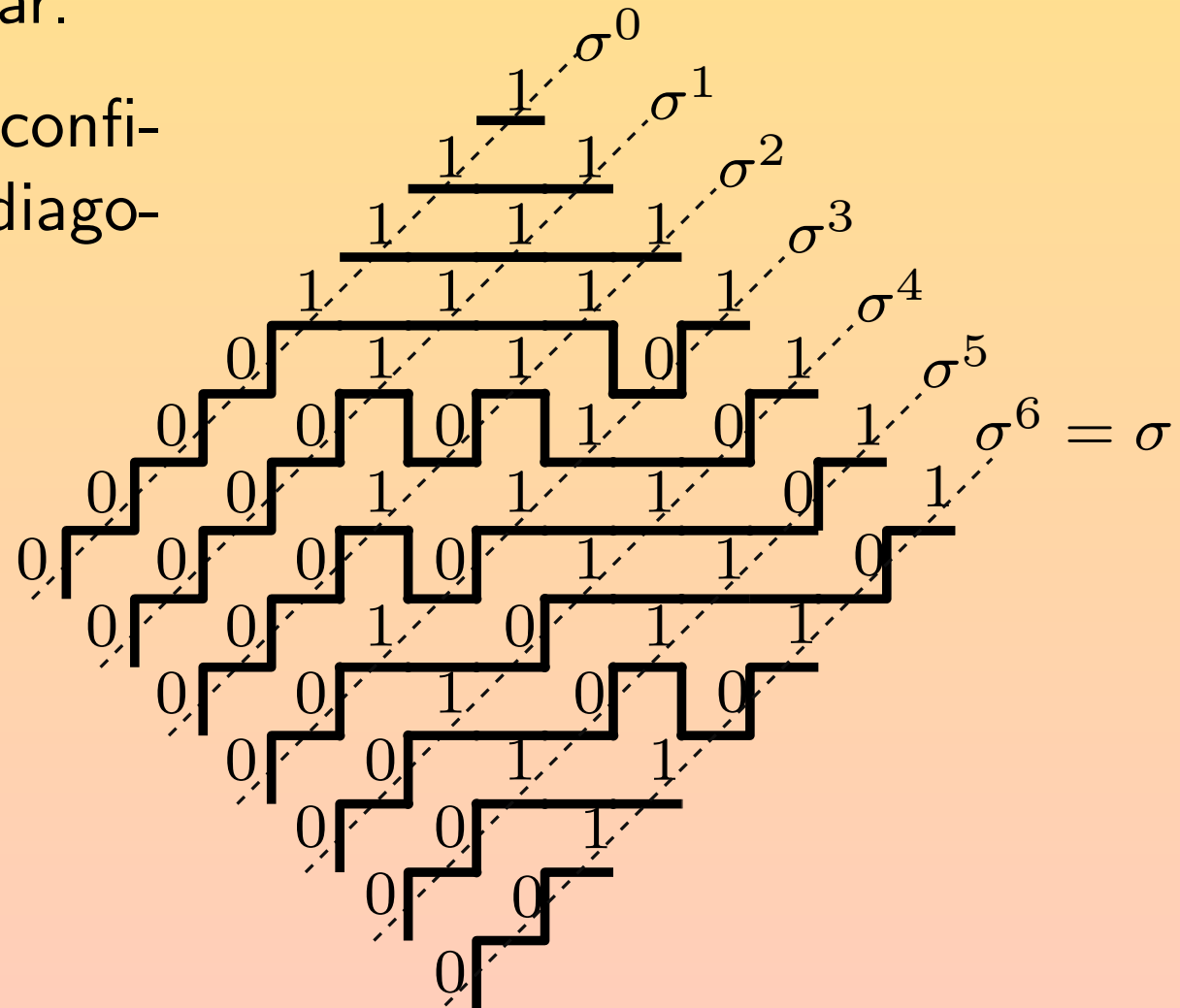
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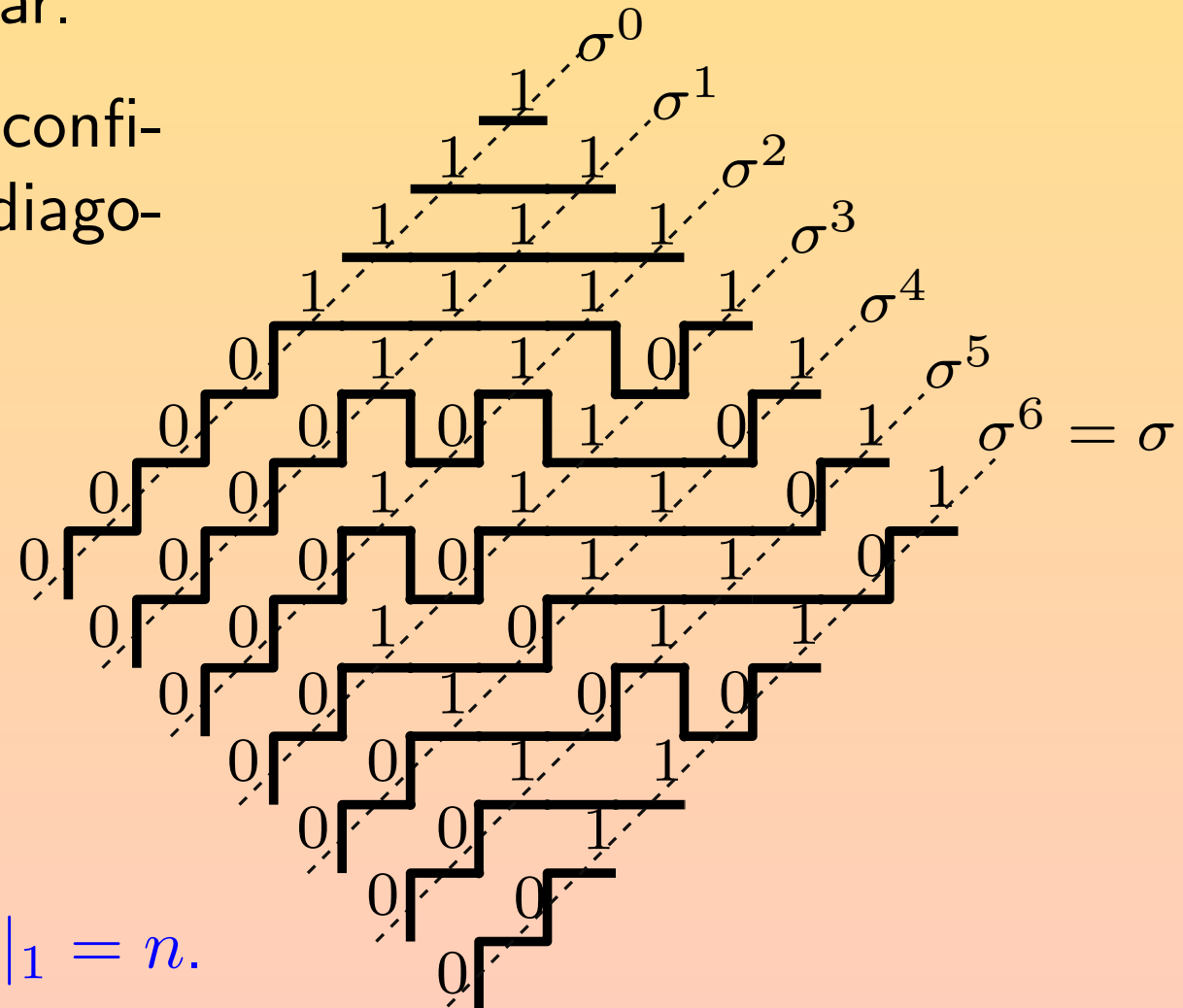
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All words  $\sigma^i$  verify  $|\sigma^i|_0 = |\sigma^i|_1 = n$ .

# Regions $\mathcal{R}_1$ and $\mathcal{R}_2$

## Proposition [CKLN '05]

For any FPL configuration in  $\mathcal{R}_1$ , the sequence of shapes  $\sigma^0, \sigma^1, \dots, \sigma^{n+k}$  form a **semistandard Young tableau**.

This is a bijection between  $\mathcal{R}_1(\sigma, k)$  and tableaux of shape  $\sigma$  and length  $n + k$ .

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$$\begin{aligned} \text{So } A_\pi(m) &= \sum_{\sigma, \tau} |\mathcal{R}_1(\sigma, 0)| \cdot t_{\sigma, \tau}^\pi \cdot |\mathcal{R}_2(\tau, m - 3n + 1)| \\ &= \sum_{\sigma, \tau} SSYT(\sigma, n) \cdot t_{\sigma, \tau}^\pi \cdot SSYT(\tau^*, m - 2n + 1) \end{aligned}$$

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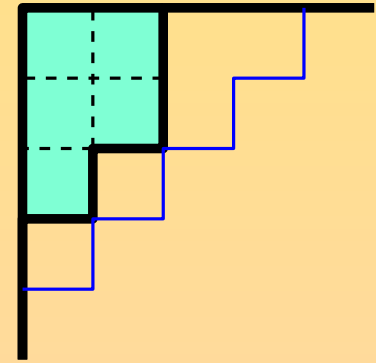
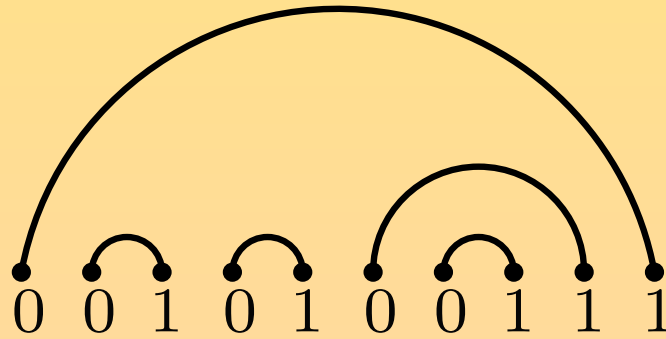
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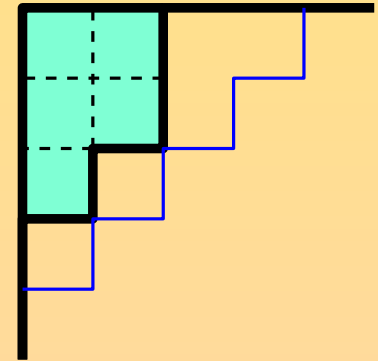
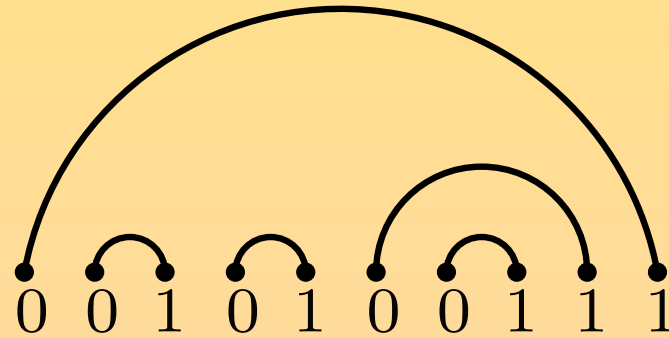
In fact,  $A_\pi(m)$  is given by the same polynomial for  $m < 3n - 1$  [CKLN '05].

Given a noncrossing matching  $\pi$  of size  $n$ , we can associate to it a word, and thus a Ferrers shape :





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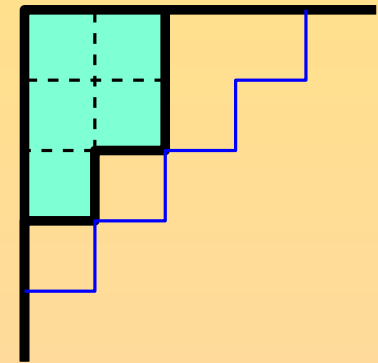
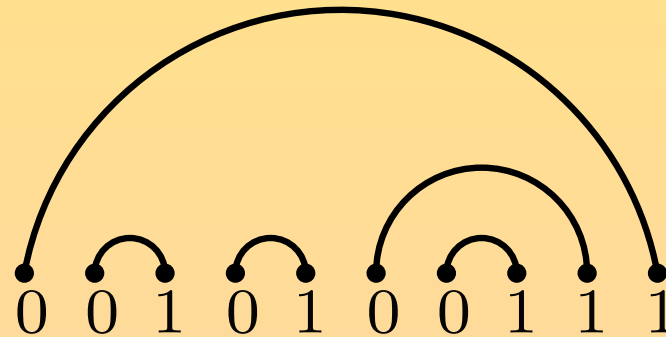


**Definition** We note  $\mathcal{D}_n$  the words  $w$  such that  $|w|_0 = |w|_1 = n$  and which are smaller than  $(01)^n$ .

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### Theorem [CKLN '04]

$t_{\sigma, \tau}^{\pi} = 0$  unless  $\sigma \leq \pi$ . Moreover,  $t_{\pi, \mathbf{0}_n}^{\pi} = 1$  and  $t_{\pi, \tau}^{\pi} = 1$  if  $\tau \neq \mathbf{0}_n$ .

### Corollary

- The formula for  $A_{\pi}(m)$  can be restricted to words  $\sigma, \tau \in \mathcal{D}_n$ ,
- The polynomial  $A_{\pi}(m)$  has leading term  $\frac{1}{h(\pi)} t^{\ell(\pi)}$ .

# The decomposition formula

We want to write  $A_\pi(m)$  as a  $\mathbb{Z}$ -linear combination of polynomials  $A_\alpha(m-1)$ , where  $\alpha, \pi$  are in  $\mathcal{D}_n$ .

# The decomposition formula

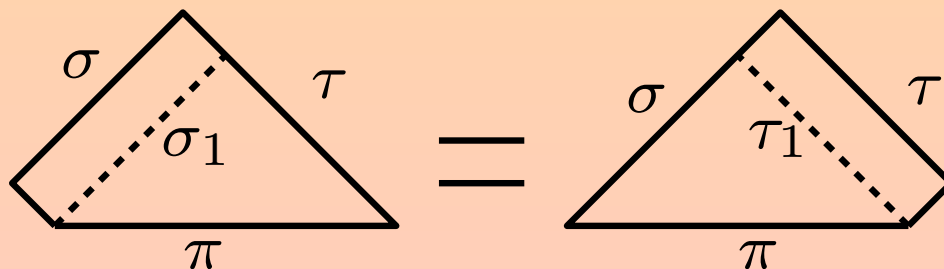
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**Theorem** [N. '09] (conjectured in [Thapper '07]).

Let  $\sigma, \tau, \pi$  be elements of  $\mathcal{D}_n$ . Then we have the equality :

$$\sum_{\substack{\sigma_1 \in \mathcal{D}_n \\ \sigma \rightarrow \sigma_1}} t_{\sigma_1, \tau}^\pi = \sum_{\substack{\tau_1 \in \mathcal{D}_n \\ \tau^* \rightarrow \tau_1^*}} t_{\sigma, \tau_1}^\pi.$$

In terms of diagrams, this means precisely that

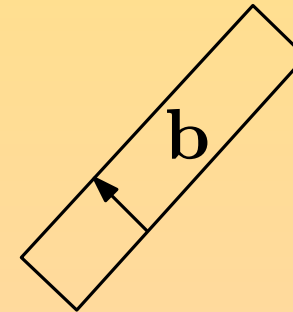


The proof is an application of Wieland's rotation.

# The decomposition formula

We now define certain matrices endomorphisms  $\mathbf{b}$ ,  $\tilde{\mathbf{b}}$ ,  $\mathbf{t}^\pi$  acting on the complex **vector space with distinguished basis**  $\mathcal{D}_n$ .

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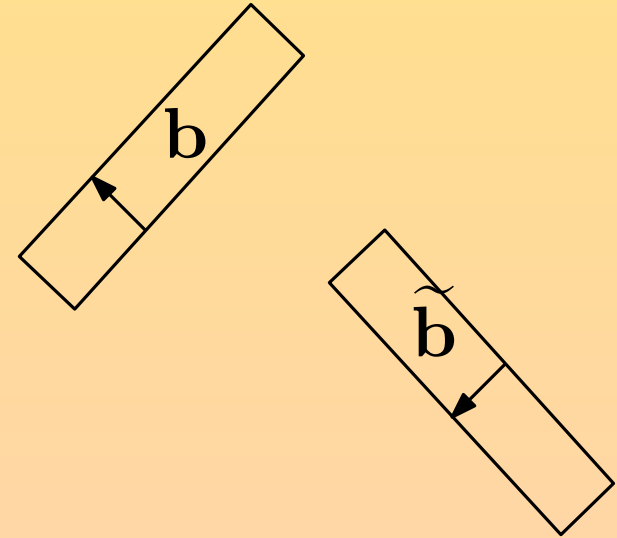


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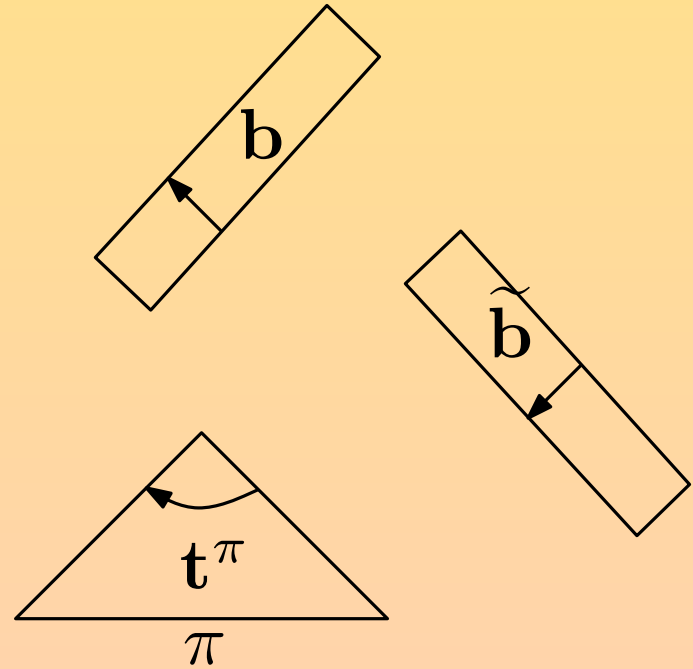
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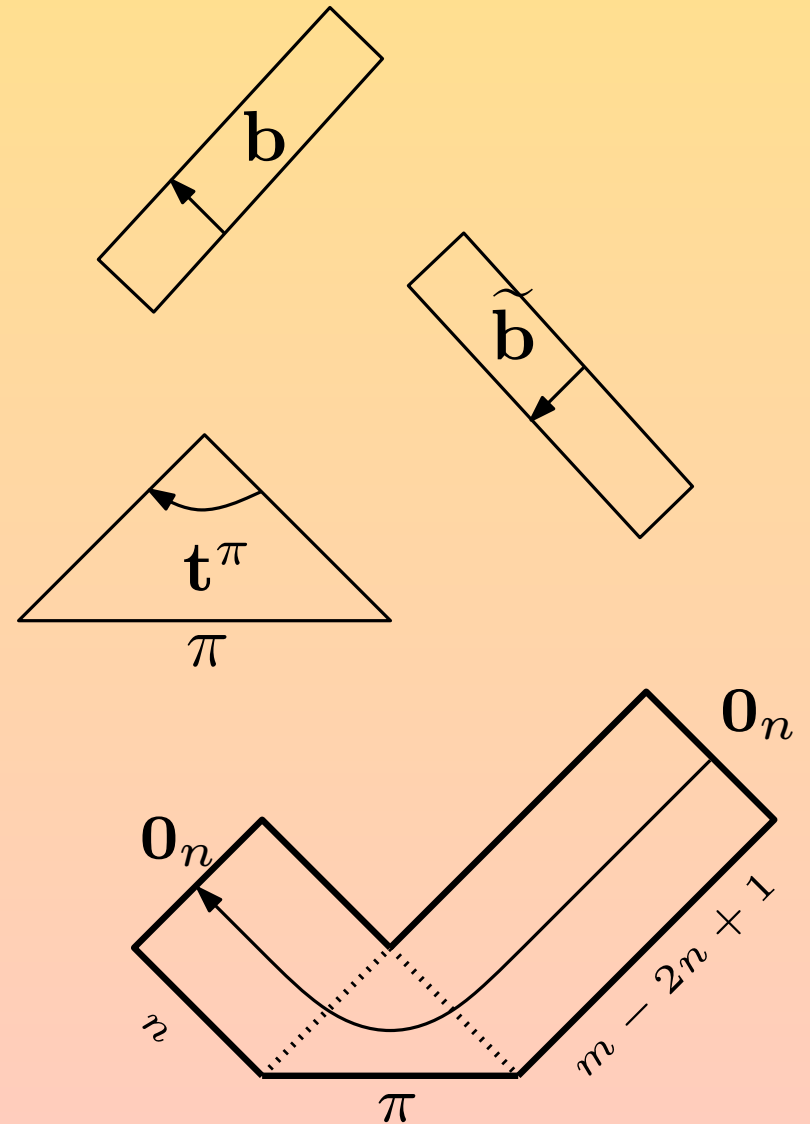
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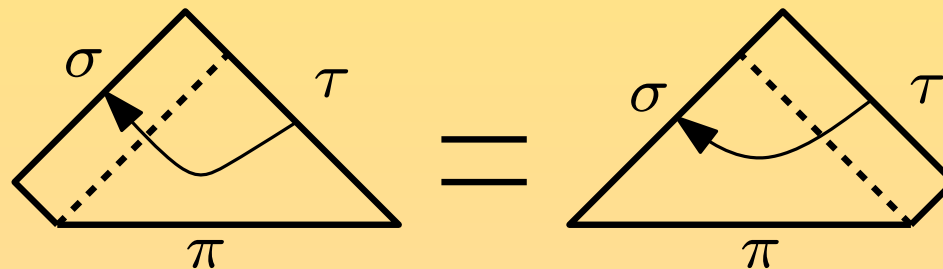
Putting these pieces together we get

$$A_\pi(m) = \begin{pmatrix} \mathbf{b}^n \mathbf{t}^\pi \tilde{\mathbf{b}}^{m-2n+1} \\ \mathbf{0}_n \mathbf{0}_n \end{pmatrix}$$



# The decomposition formula

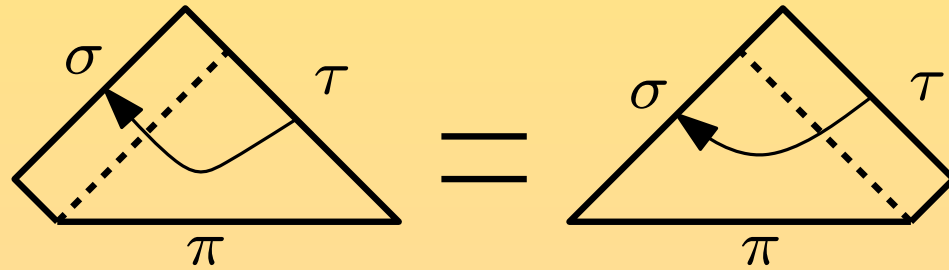
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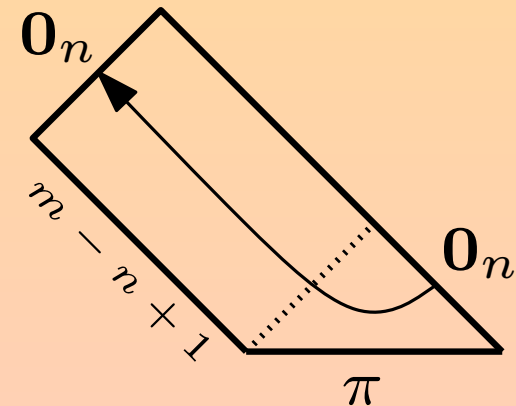
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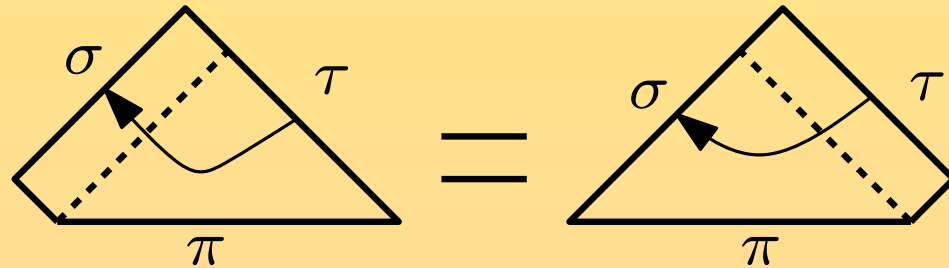
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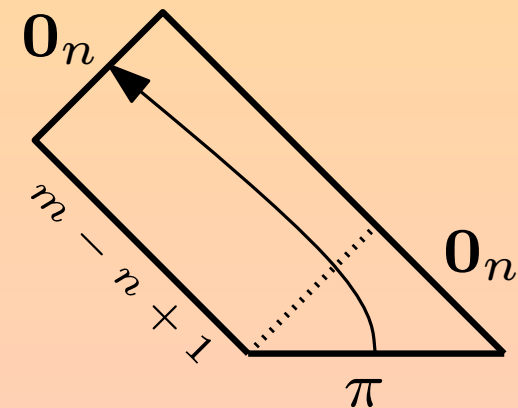
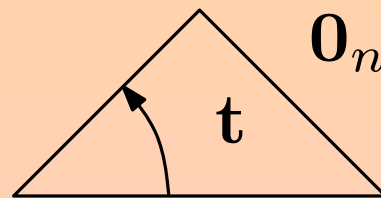


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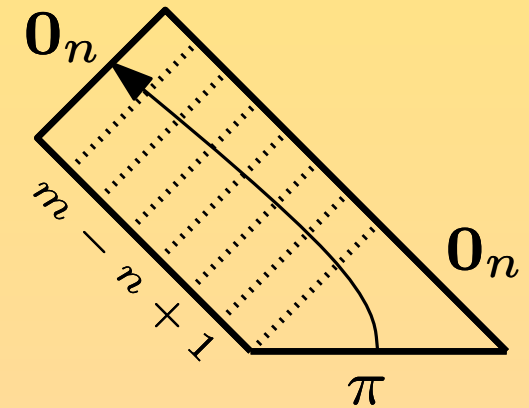
we can rewrite this as  $A_\pi(m) = (\mathbf{b}^{m-n+1}\mathbf{t})_{\mathbf{0}_n\pi}$

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## Proposition [N. '09]

For  $\pi \in \mathcal{D}_n$ , and  $m$  an integer

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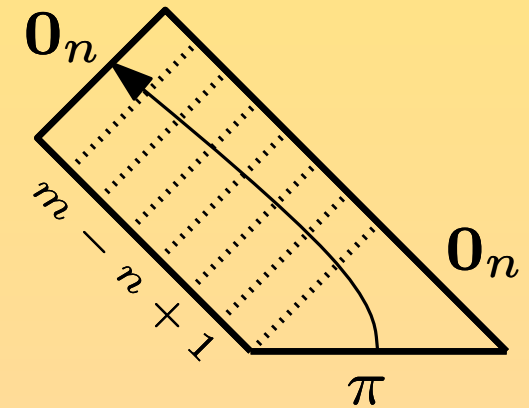


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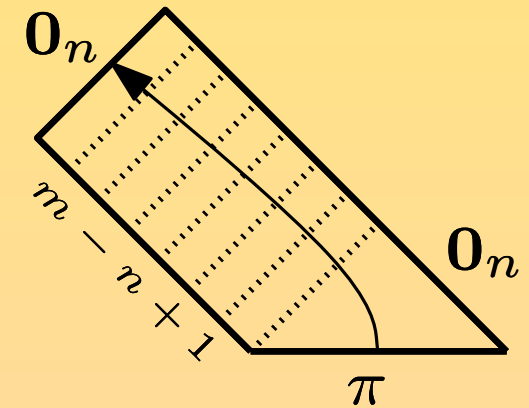
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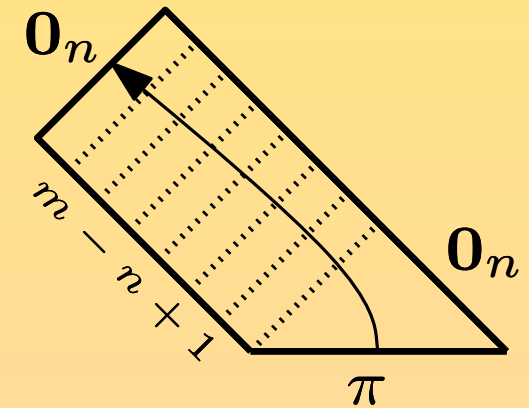
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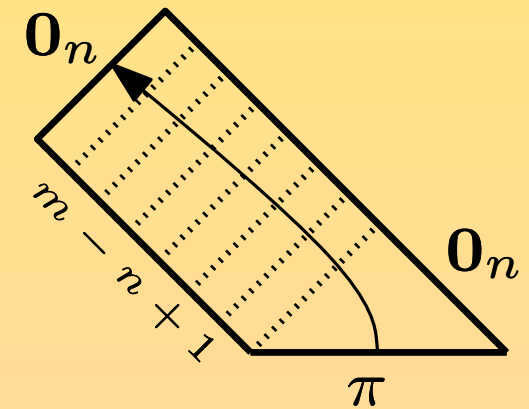


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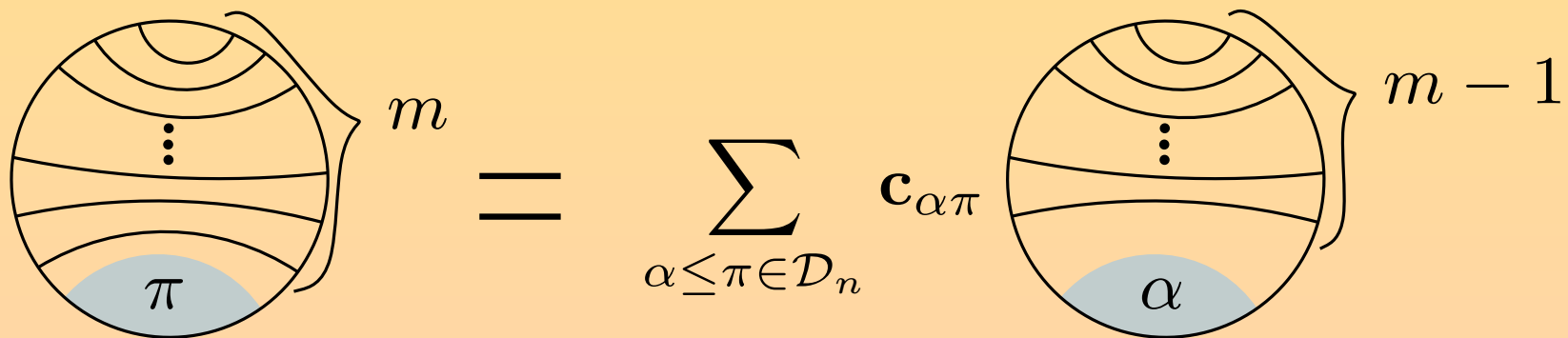
$$A_\alpha(m-1) \\ //$$

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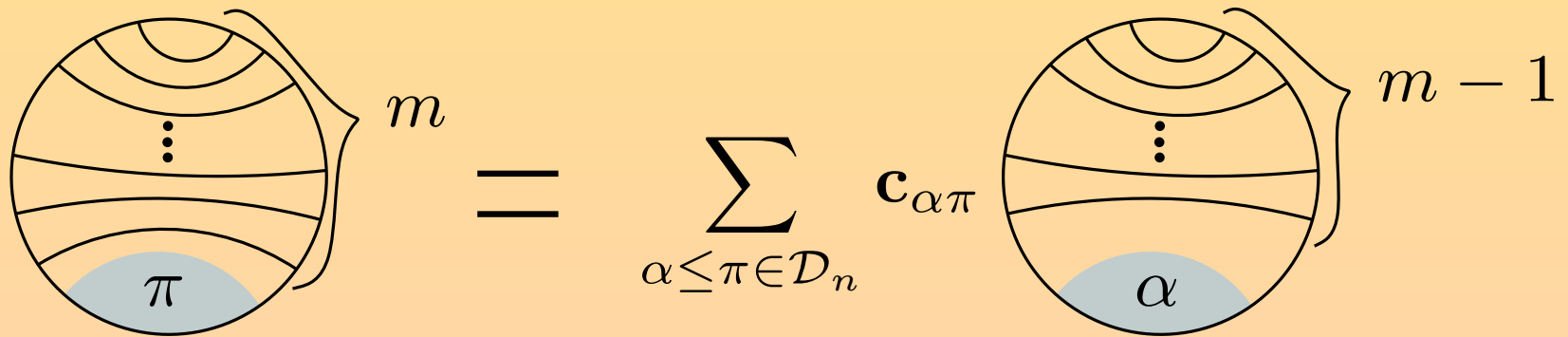
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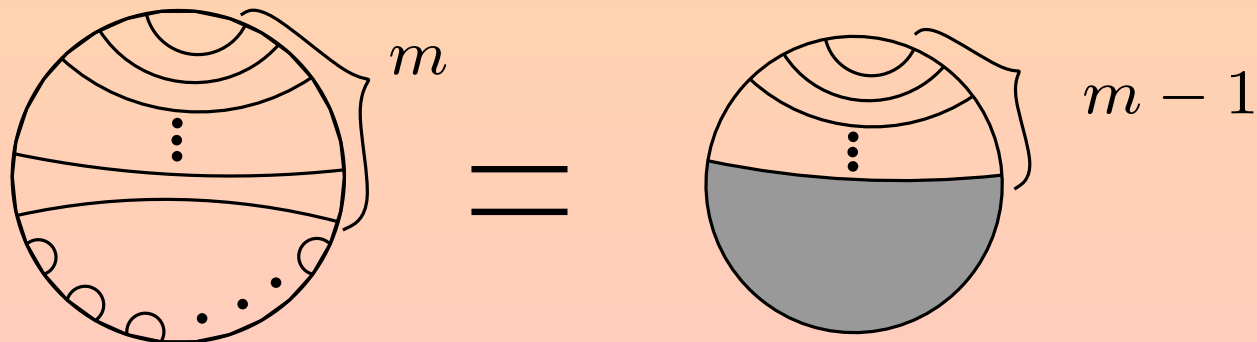
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## Conjecture [Thapper]

If  $\pi = \mathbf{1}_n$ , then  $\mathbf{c}_{\alpha\mathbf{1}_n} = 1$  for any  $\alpha \in \mathcal{D}_n$ .

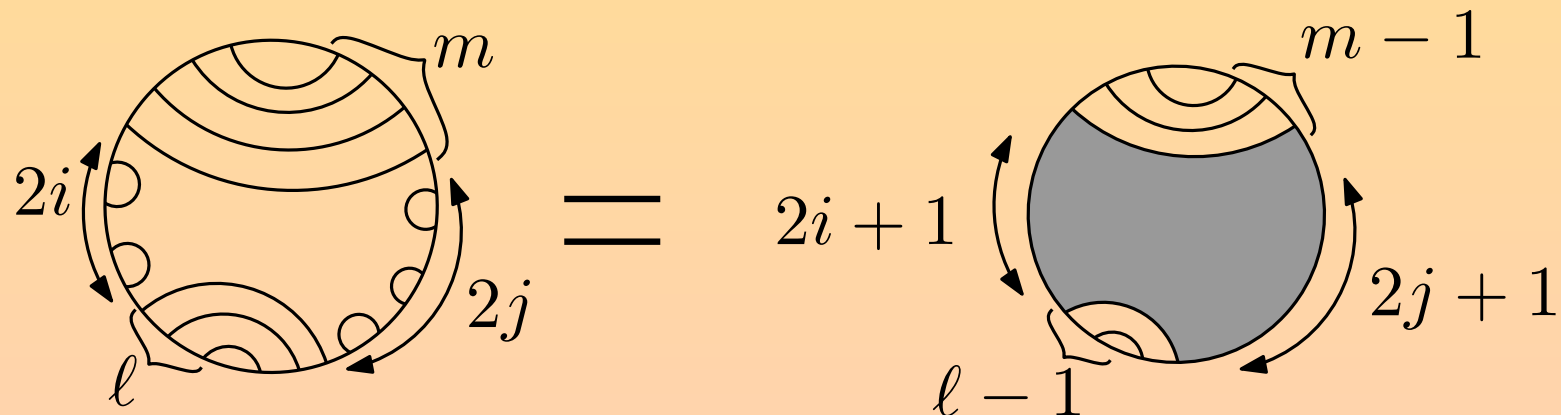
This implies



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Some remarks on the coefficients  $c_{\alpha\pi}$  : there are of course not the unique numbers such that the previous theorem holds. But, based on data for small  $n$ , these numbers **conjecturally** :

- give nice decomposition formulas, for instance :

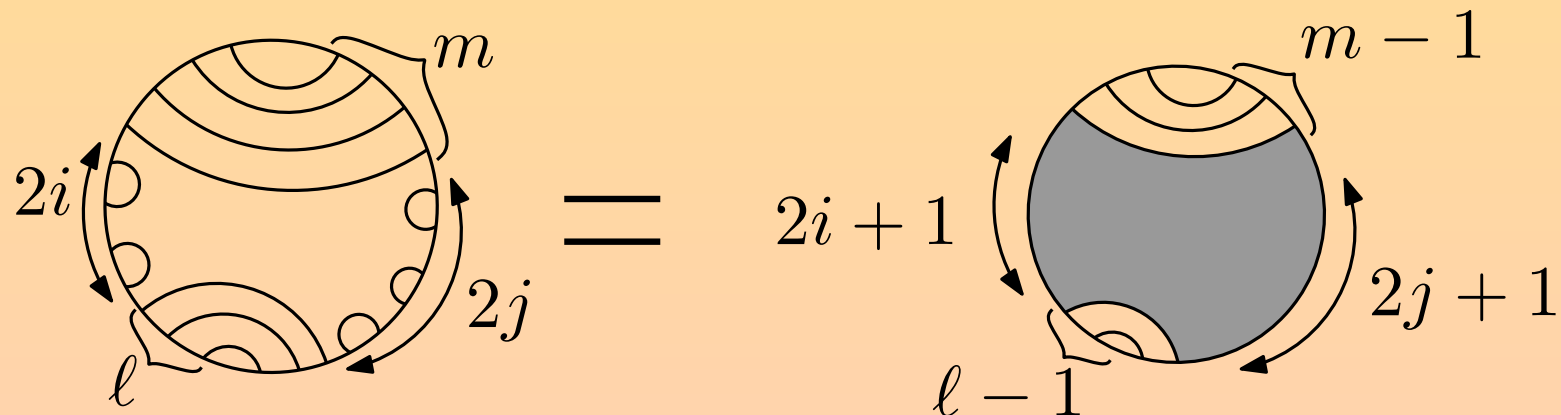


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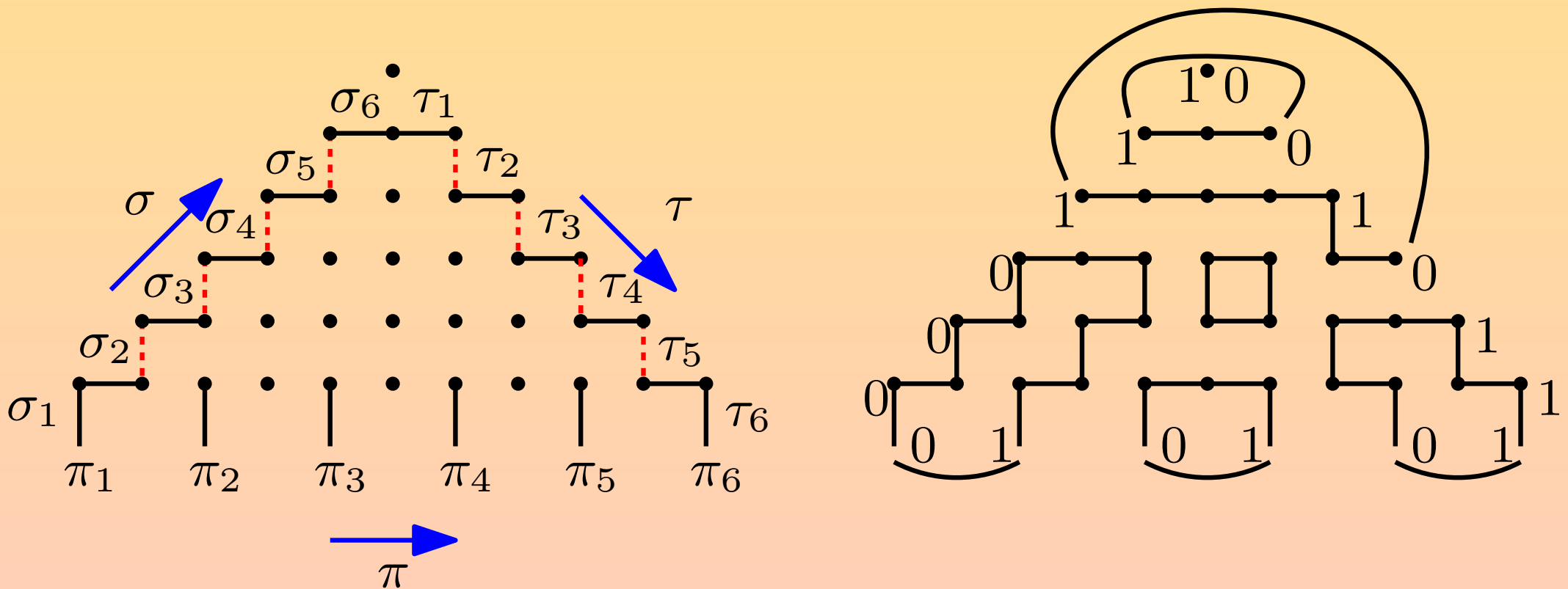


- verify  $c_{\alpha\pi} = c_{\alpha^*\pi^*}$  and  $c_{0\alpha 1, 0\pi 1} = c_{\alpha\pi}$ .

Challenge : conjecture a direct combinatorial description of these coefficients.

# The triangle $\mathcal{T}_n$

We now study the FPL configurations in the triangle, in short **TFPL configurations**. Given the boundary data  $\sigma, \pi, \tau$ , we want to compute  $t_{\sigma, \tau}^{\pi}$  which is the number of TFPL configurations with these boundaries.



# The numbers $t_{\sigma,\tau}^\pi$

## Proposition

$$t_{\sigma,\tau}^\pi = t_{\tau^*,\sigma^*}^{\pi^*}.$$

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## Theorem [CKLN '04]

$$t_{\sigma,\tau}^\pi = 0 \text{ unless } \sigma \leq \pi.$$

Sketch of the proof : the idea is to show that there exist integers  $N_i(f) \geq 0$  (with  $N_0(f) = 0$ ) attached to a TFPL configuration  $f$ , such that if  $f$  has boundary data  $\sigma, \pi, \tau$ , then  $\sigma_i - \pi_i = N_i(f) - N_{i-1}(f)$  for all  $i$ .



# Common prefixes and suffixes

Now we study the case where  $\sigma$  and  $\pi$  have common prefixes and suffixes.

## Proposition

Let  $\pi, \sigma, \tau \in \mathcal{D}_n$ , and suppose that there exist words  $u, \sigma', \pi', v$  such that  $\sigma = u\sigma'v$  and  $\pi = u\pi'v$ . Define  $a, b$  by  $n - a = |u|_0 + |v|_0$  and  $n - b = |u|_1 + |v|_1$ .

Then  $t_{\sigma, \tau}^{\pi} = 0$  unless  $\tau = 0^{n-a} \tau' 1^{n-b}$  for a certain  $\tau'$ .

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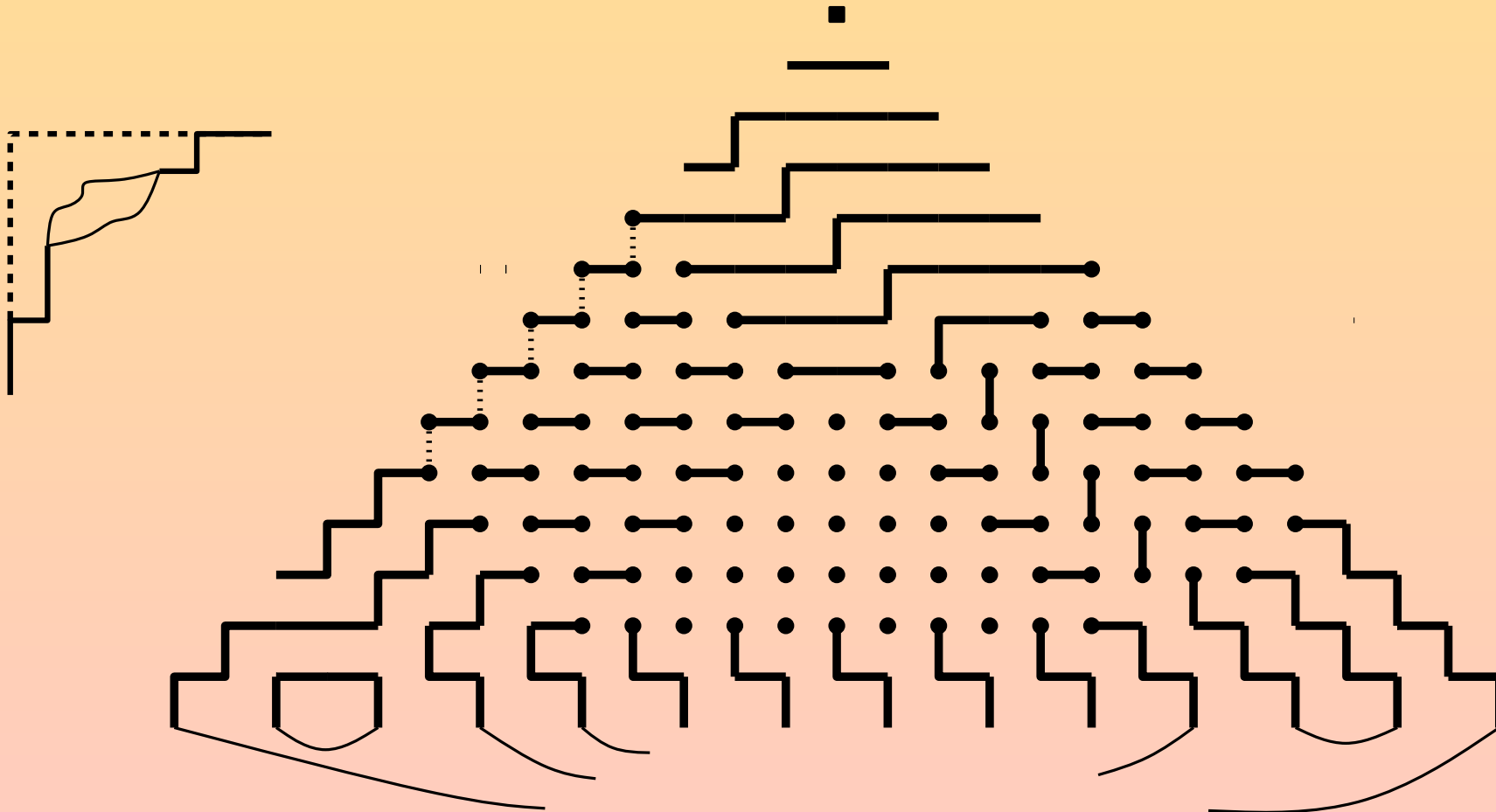
## Proposition

If in addition  $\pi' = 1^b 0^a$ , then  $t_{\sigma, \tau}^{\pi}$  is given by the determinant of a matrix of size  $a$  (or  $b$ ) with entries given by certain binomial coefficients.

# Common prefixes and suffixes

Proof : lots of fixed edges.

Case where  $\sigma$  and  $\pi$  have common prefix and suffix.

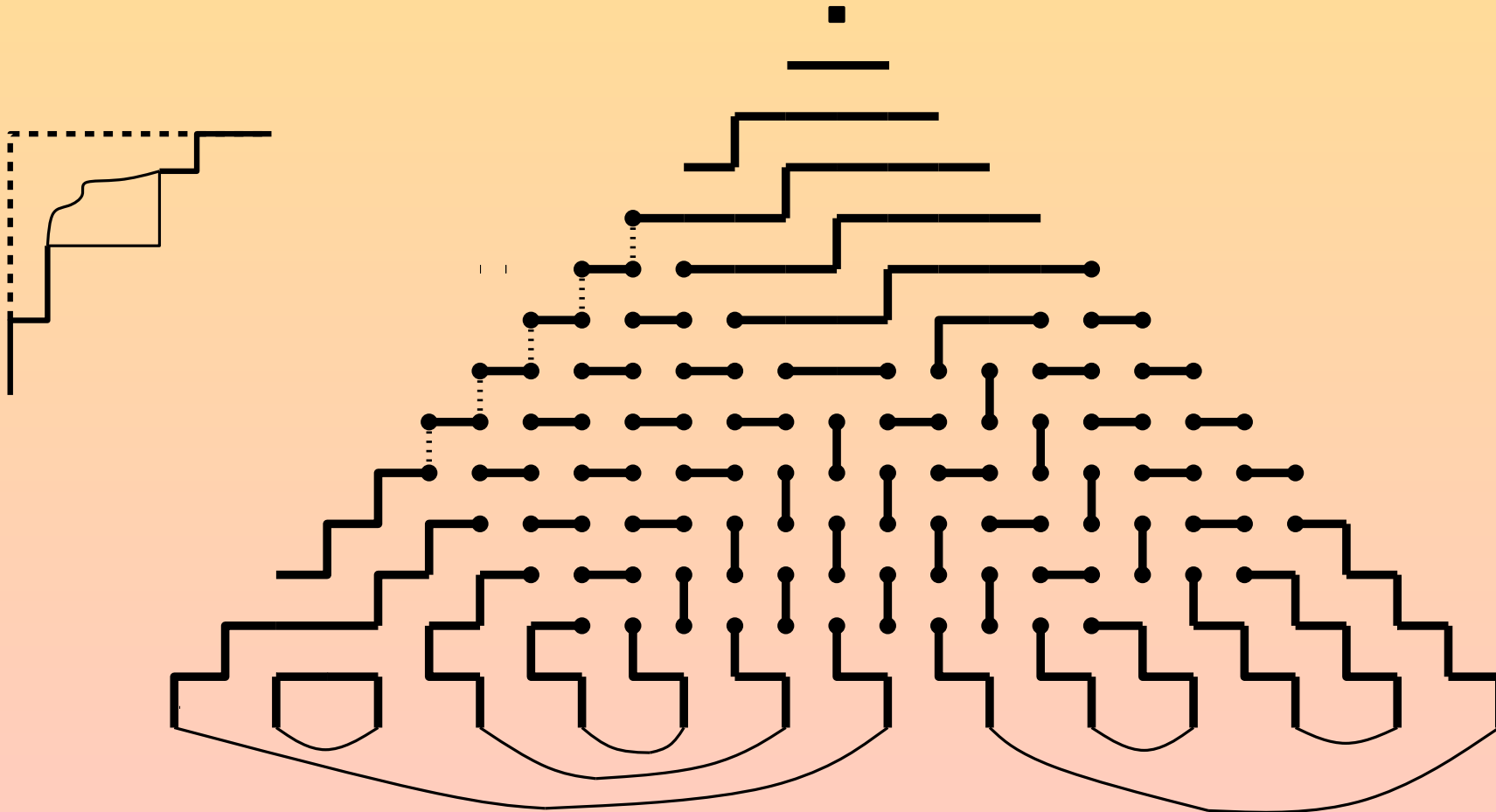


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And  $\pi'$  is equal to  $1^a 0^b$ .



# Triangles and Littlewood-Richardson coefficients

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**Proposition** For every  $\pi \in \mathcal{D}_n$

$$\frac{1}{h(\pi)} = \sum_{\substack{\sigma, \tau \in \mathcal{D}_n \\ \ell(\sigma) + \ell(\tau) = \ell(\pi)}} t_{\sigma, \tau}^{\pi} \cdot \frac{1}{2^{\ell(\sigma)} h(\sigma)} \cdot \frac{1}{2^{\ell(\tau)} h(\tau)}$$

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*Sketch of the proof* : remember that  $A_{\pi}(m)$  is a polynomial of degree  $\ell(\pi)$  and leading coefficient  $1/h(\pi)$ . It can be written as

$$\sum t_{\sigma, \tau}^{\pi} \cdot SSYT(\sigma, n + k) \cdot SSYT(\tau^*, m + 1 - k - 2n).$$

Choose  $k = m/2$ , and compare the coefficients in degrees  $\ell(\pi)$  and higher to get the formula.

# Littlewood Richardson coefficients

Let  $\lambda, \mu, \nu$  be partitions, and  $\Lambda(x)$  be the ring of symmetric functions of the variables  $x_1, x_2, \dots$ . The **Schur functions**  $s_\lambda(x)$  can be defined as

$$s_\lambda(x) = \sum_T x_i^{T_i},$$

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Schur functions form a basis of  $\Lambda(x)$ . We can expand  $s_\mu(x)s_\nu(x)$  on this basis, where the coefficients  $c_{\mu,\nu}^\lambda$  are often called the **Littlewood-Richardson (LR) coefficients**.

$$s_\mu(x)s_\nu(x) = \sum_\lambda c_{\mu,\nu}^\lambda s_\lambda(x)$$

Since all terms in  $s_\lambda$  have degree  $\ell(\lambda)$ , we get

$$c_{\mu,\nu}^\lambda = 0 \text{ unless } \ell(\lambda) = \ell(\mu) + \ell(\nu).$$

# Littlewood Richardson coefficients

These coefficients appear in other places in the theory of symmetric functions ; we have for instance :

$$s_{\lambda/\mu}(x) = \sum_{\nu} c_{\mu,\nu}^{\lambda} s_{\nu}(x)$$

We have also, if  $s_{\lambda}(x, y)$  is the symmetric function  $s_{\lambda}$  in the variables  $x_1, x_2, \dots, y_1, y_2, \dots$

$$s_{\lambda}(x, y) = \sum_{\mu, \nu} c_{\mu,\nu}^{\lambda} s_{\mu}(x) s_{\nu}(y)$$

If we evaluate this at  $x_i = y_i = 1$  for  $i = 1, \dots, m/2$ ,  $x_i = y_i = 0$  for  $i > m/2$ , we obtain polynomials in  $m$  which give the following identity in top degree  $\ell(\lambda)$  :

$$\frac{1}{h(\lambda)} = \sum_{\mu, \nu} c_{\mu,\nu}^{\lambda} \cdot \frac{1}{2^{\ell(\mu)} h(\mu)} \cdot \frac{1}{2^{\ell(\nu)} h(\nu)}$$

# Littlewood Richardson coefficients

As a consequence, there exist  $a_{\sigma\tau} > 0$  such that, for any  $\pi \in \mathcal{D}_n$ ,

$$\sum_{\sigma, \tau} a_{\sigma\tau} c_{\sigma, \tau}^{\pi} = \sum_{\sigma, \tau} a_{\sigma\tau} t_{\sigma, \tau}^{\pi} \quad (E)$$

in which  $\sigma, \tau$  go through all words such that  $\ell(\sigma) + \ell(\tau) = \ell(\pi)$

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For all words  $\pi, \sigma, \tau \in \mathcal{D}_n$  verifying  $\ell(\sigma) + \ell(\tau) = \ell(\pi)$ , we have

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Thanks to equation (E), we need only prove that  $c_{\sigma, \tau}^{\pi} \leq t_{\sigma, \tau}^{\pi}$  for all valid  $\sigma, \tau, \pi$ .

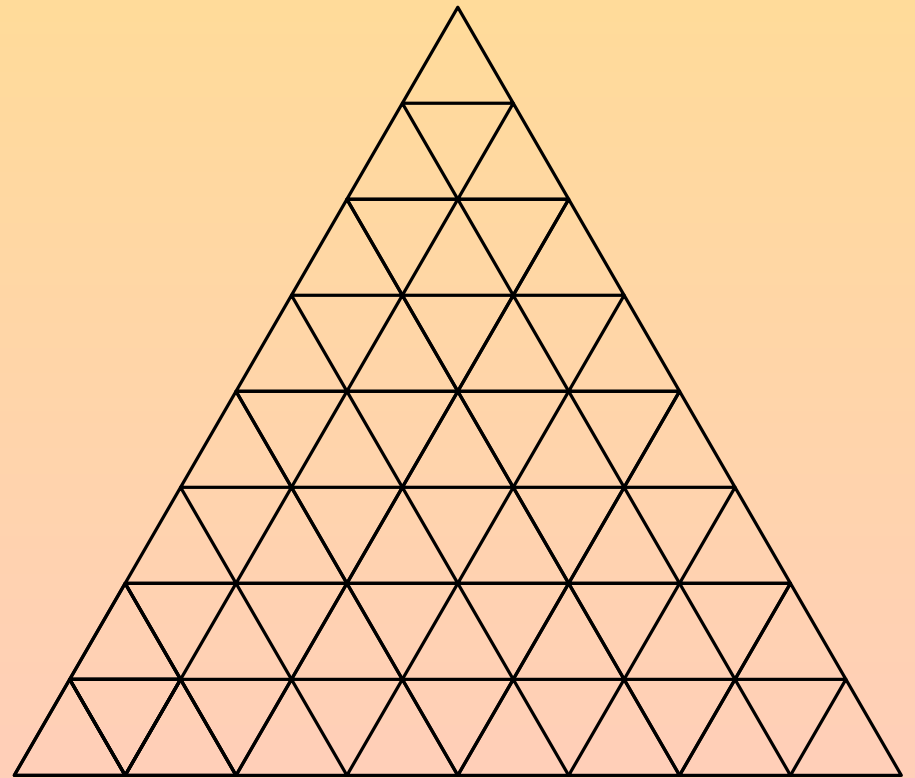
# Computing LR coefficients

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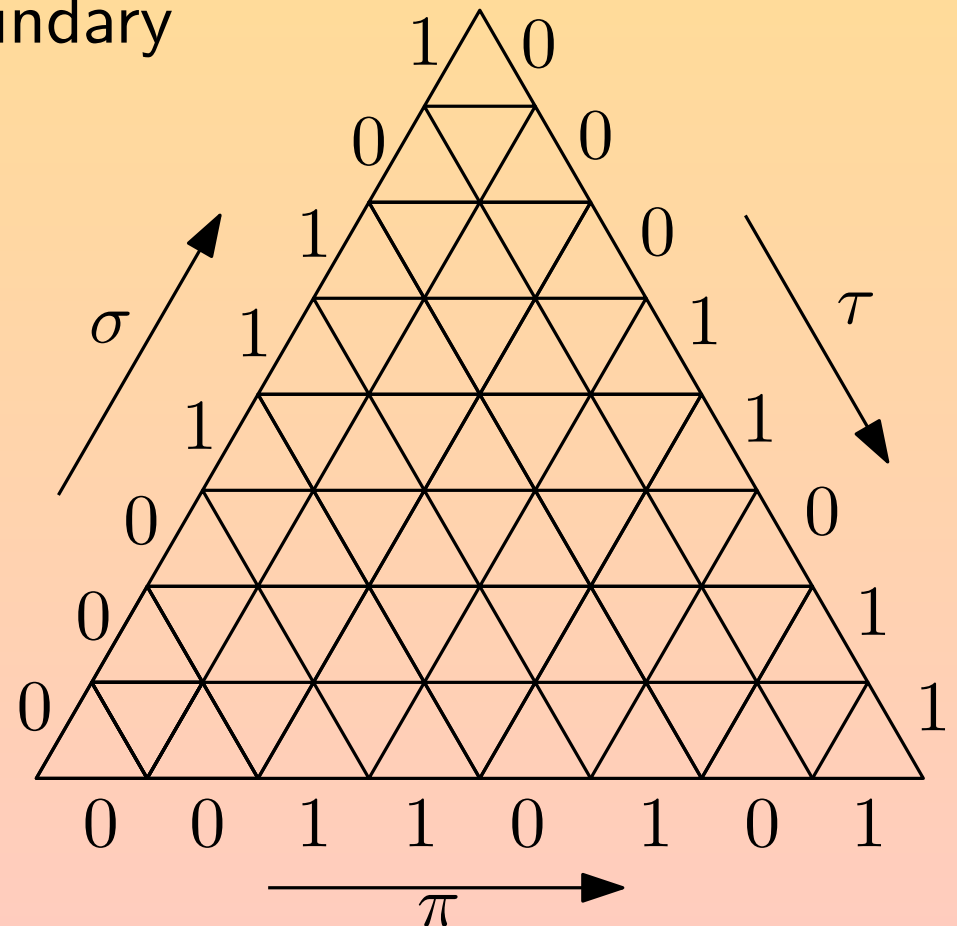
Consider a triangle of size  $2n$  on the triangular lattice.

Fix  $\sigma, \pi, \tau \in \mathcal{D}_n$ , and label the boundary edges of the triangle.

$$\pi = 00110101$$

$$\sigma = 00011011$$

$$\tau = 00011011$$

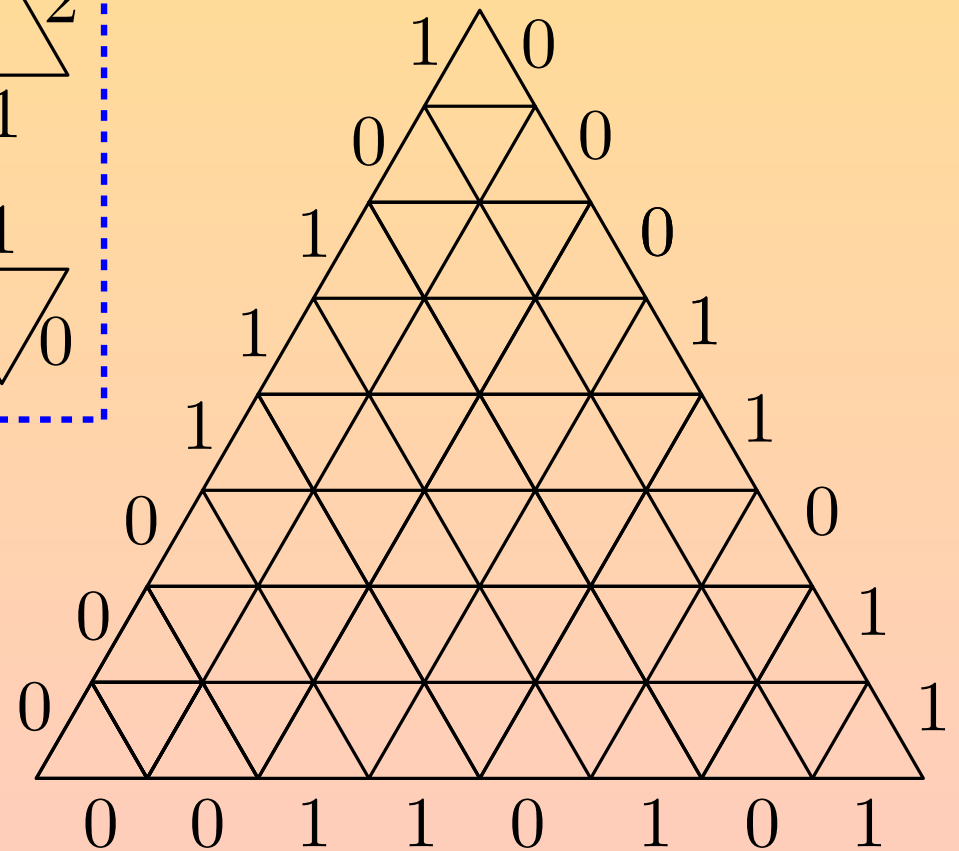
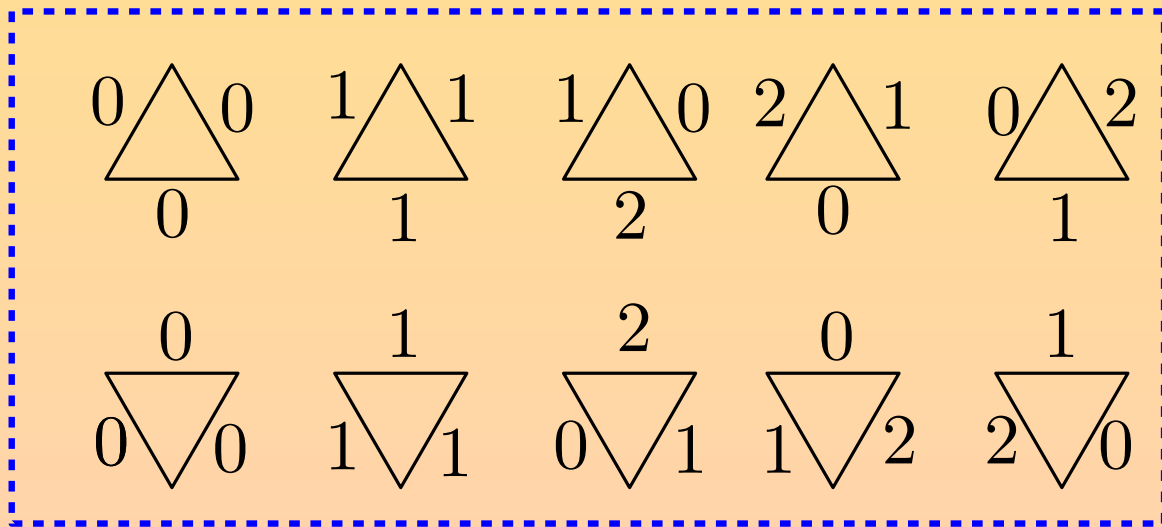




# Definition

A **Knutson-Tao puzzle** with boundary data  $\sigma, \pi, \tau$  is a labeling of each edge of the triangle by 0, 1 or 2, such that :

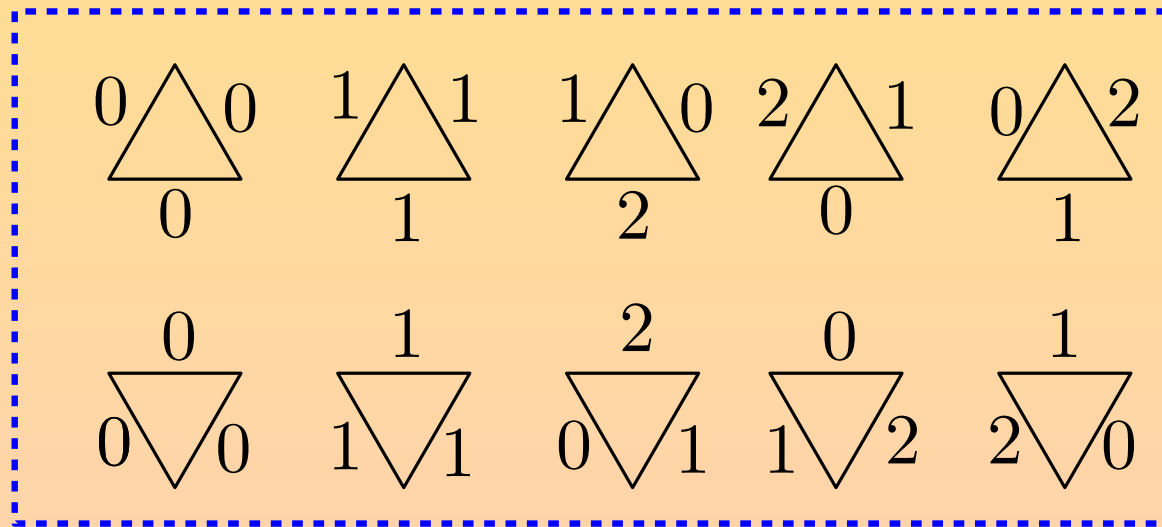
- the labels on the boundary is given by  $\sigma, \pi, \tau$  ;
- on each triangle, the induced labeling must be among :



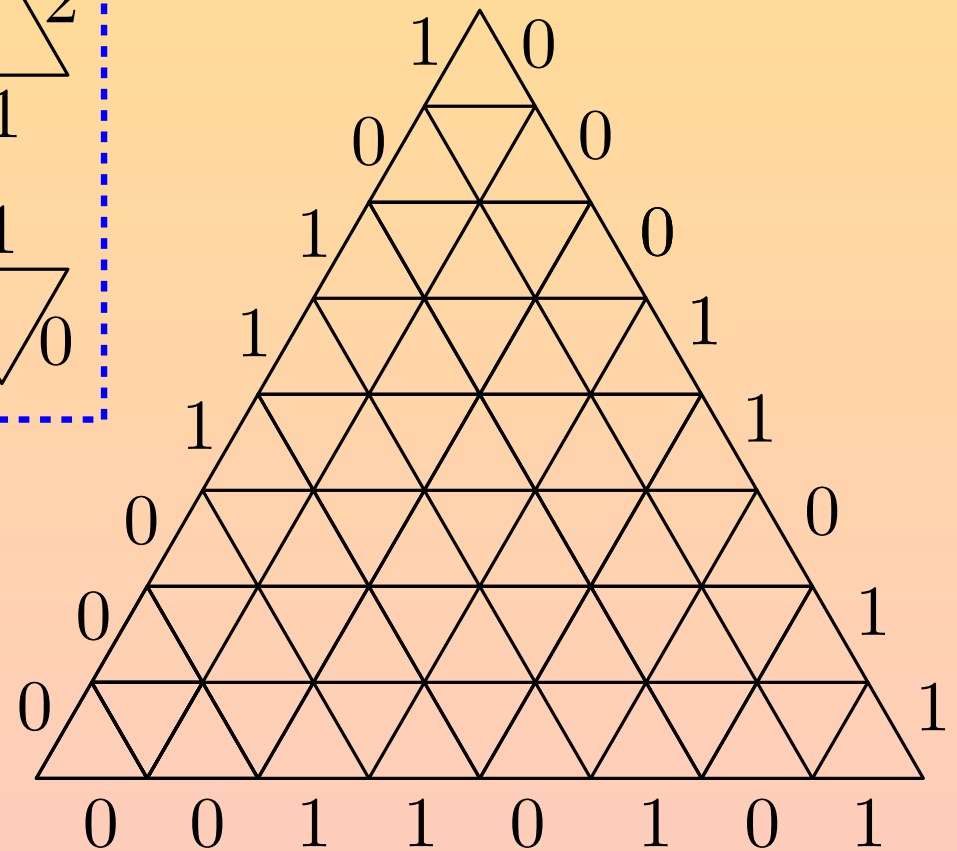
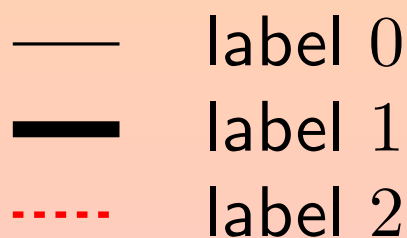
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We will picture the labeling of edges as follows :

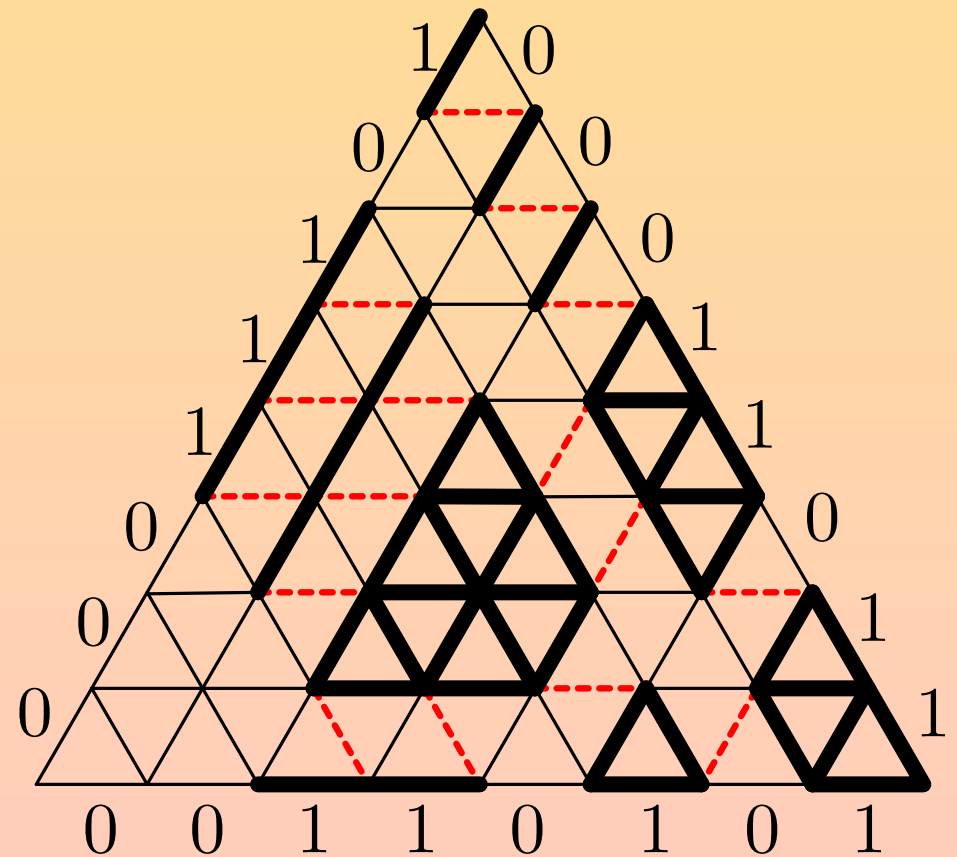
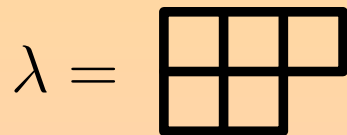


# Theorem [Knutson, Tao '03][K., T. and Woodward '03]

Let  $\sigma, \tau, \pi \in \mathcal{D}_n$ . Then the number of KT-puzzles with boundary data  $\sigma, \pi, \tau$  is equal to the LR coefficient  $c_{\sigma, \tau}^{\pi}$ .

For example, it is easy to see that there is only one puzzle with the boundary data of the example.

so  $c_{\mu, \nu}^{\lambda} = 1$  where



# From KT puzzles to TFPL configurations.

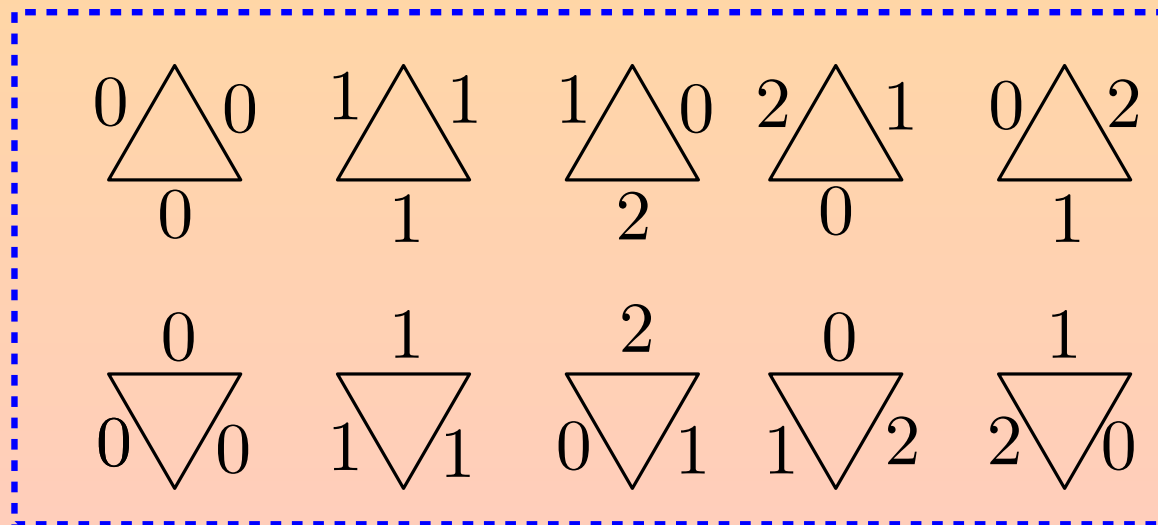
We fix  $\sigma, \pi, \tau \in \mathcal{D}_n$ , such that  $\ell(\sigma) + \ell(\tau) = \ell(\pi)$ . We will define a map  $\Phi$ .

KT puzzles with boundary data  $\sigma, \pi, \tau$



TFPL configurations with boundaries  $\sigma, \tau, \pi$

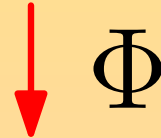
The map is **local** : it changes every small labeled triangle of the puzzle to a piece of a path of a TFPL configuration.



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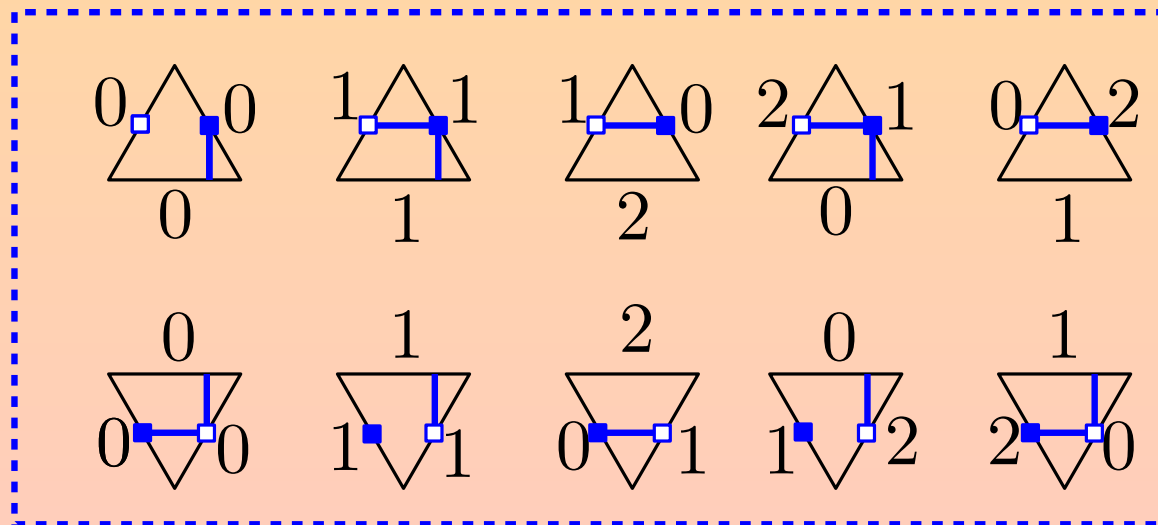
KT puzzles with boundary data  $\sigma, \pi, \tau$



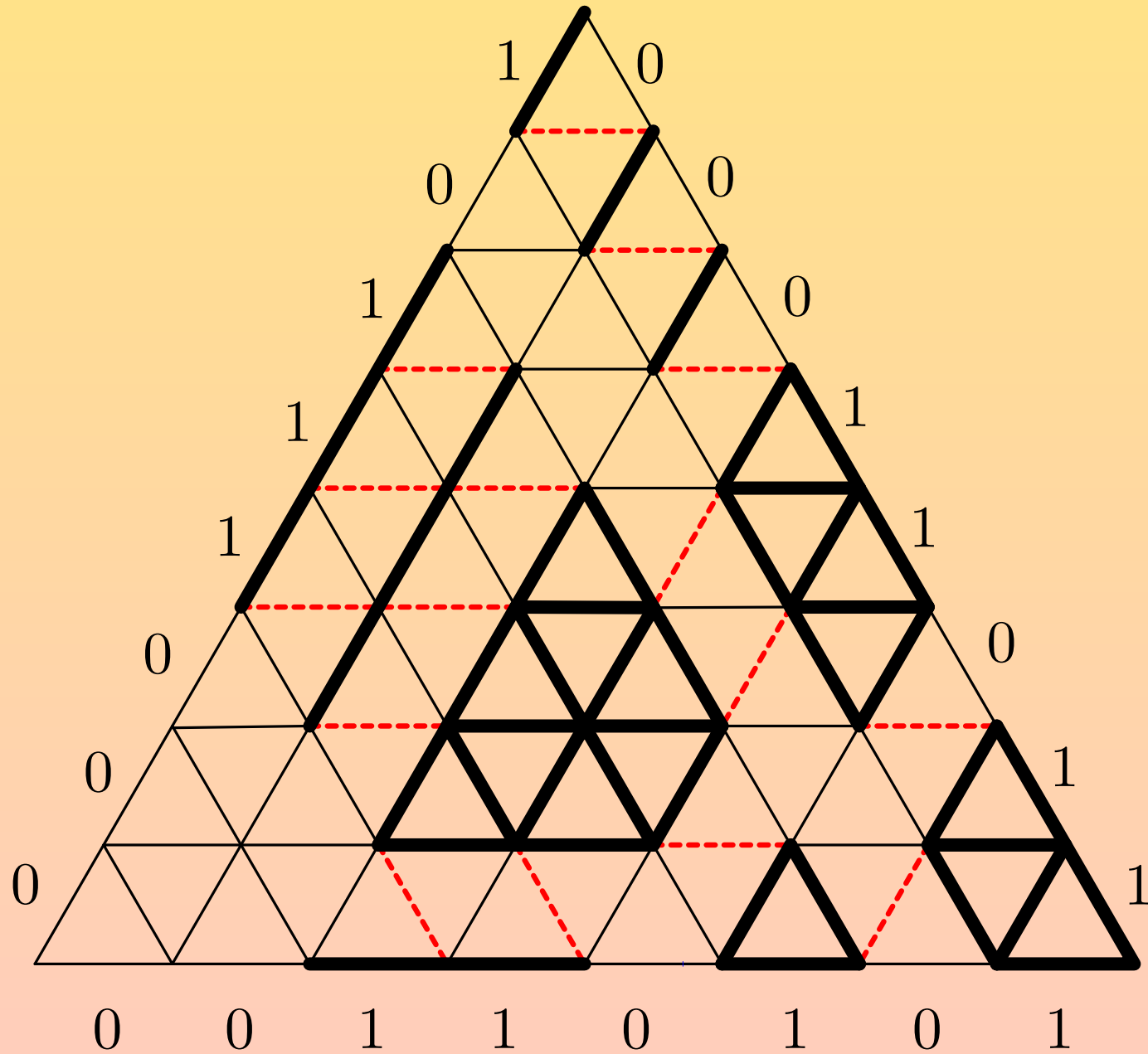
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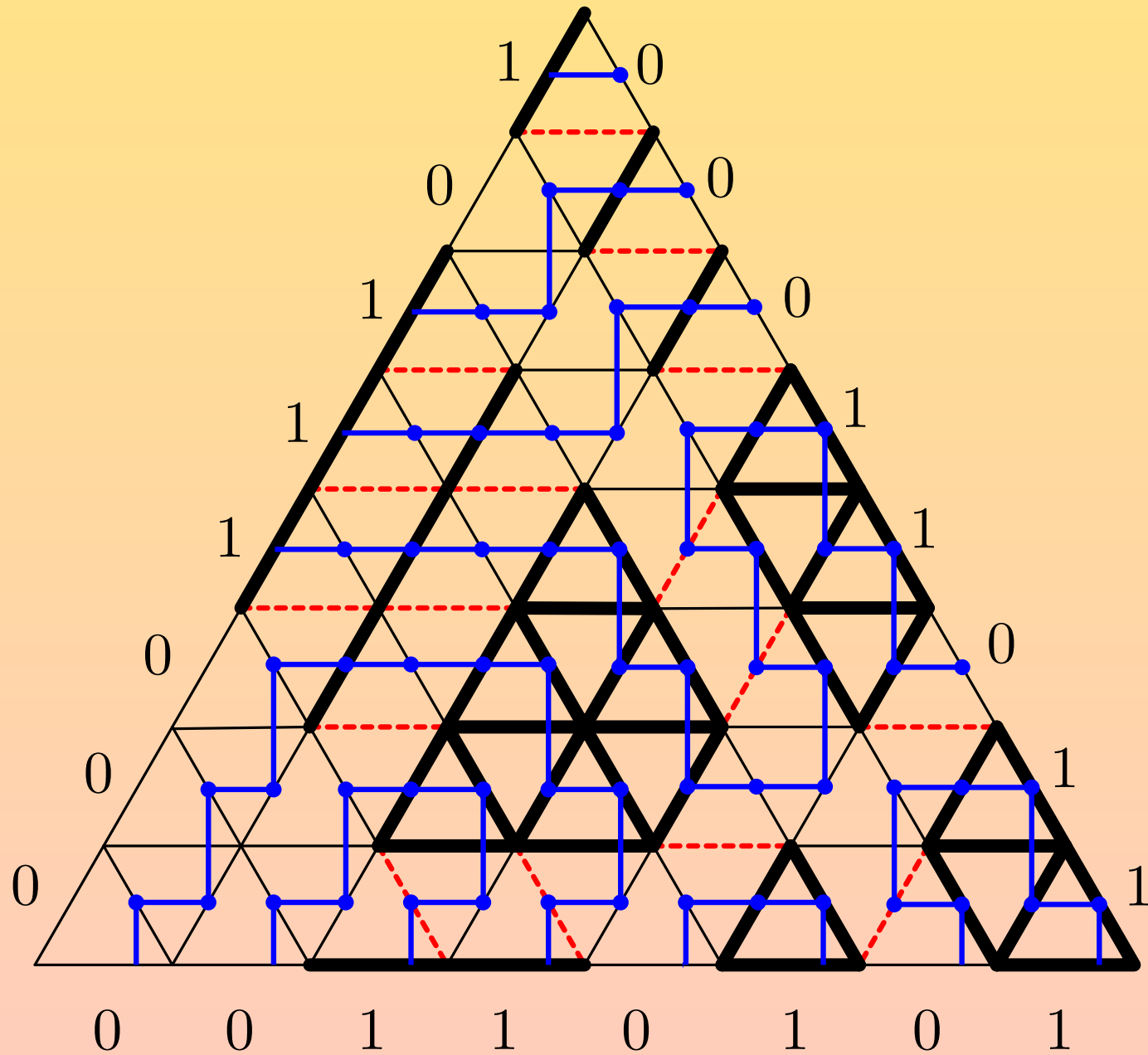
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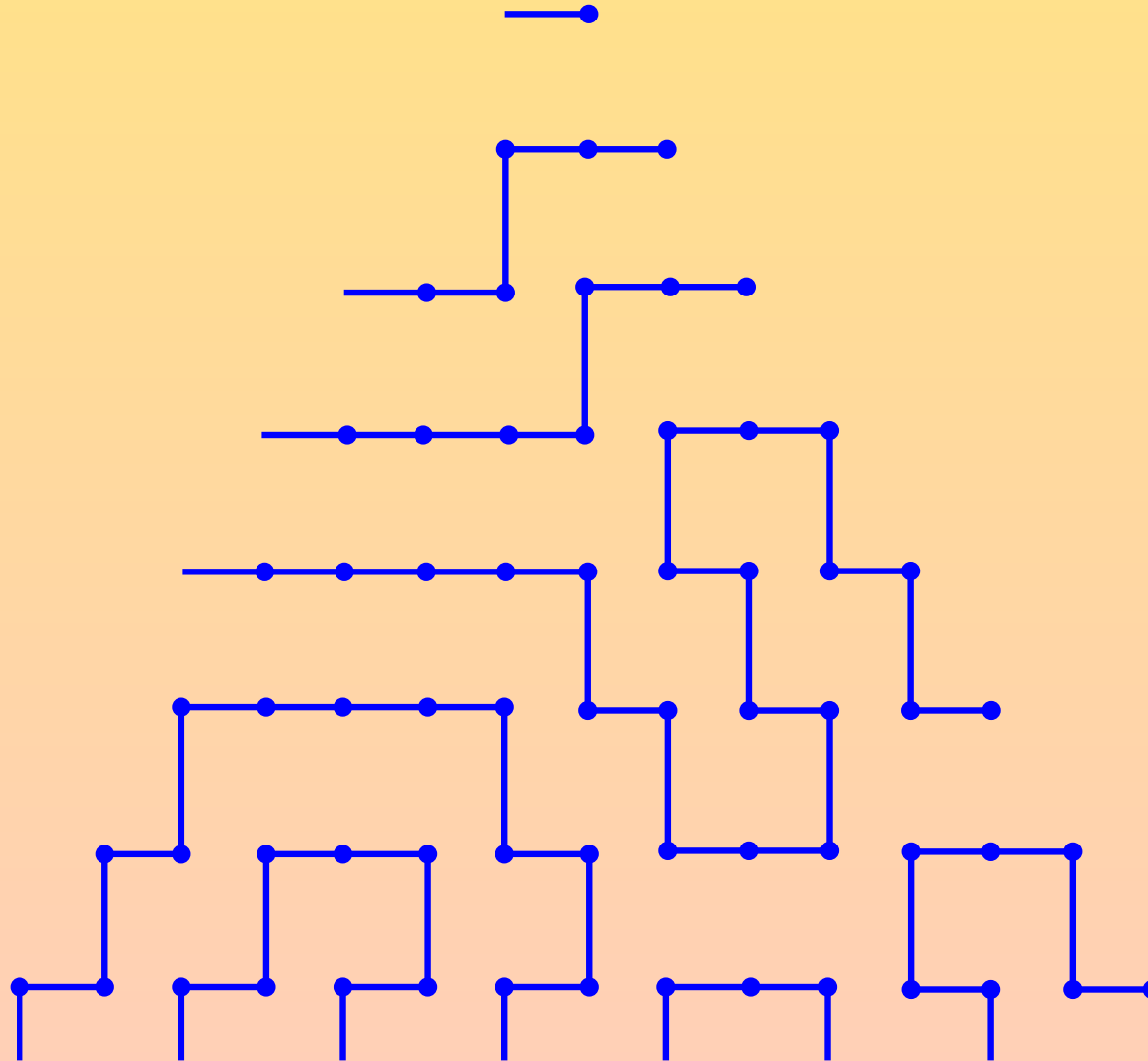
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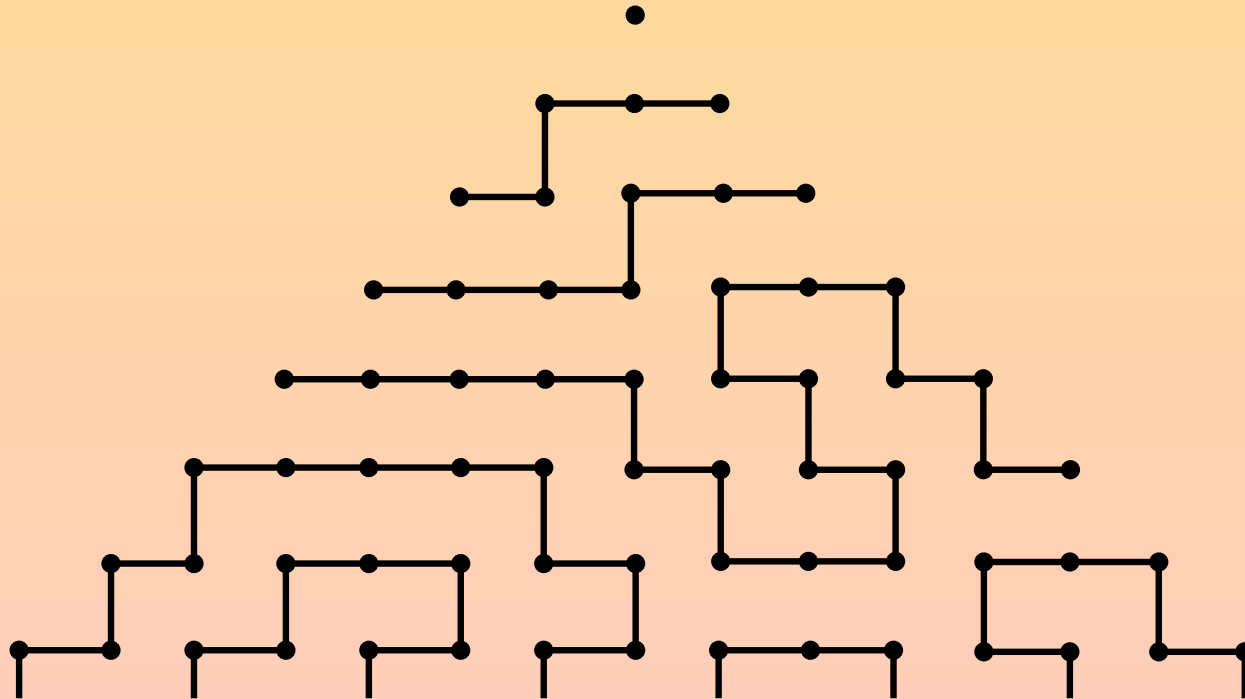
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The result is a TFPL configuration, with boundary data  $\sigma, \pi, \tau$ .

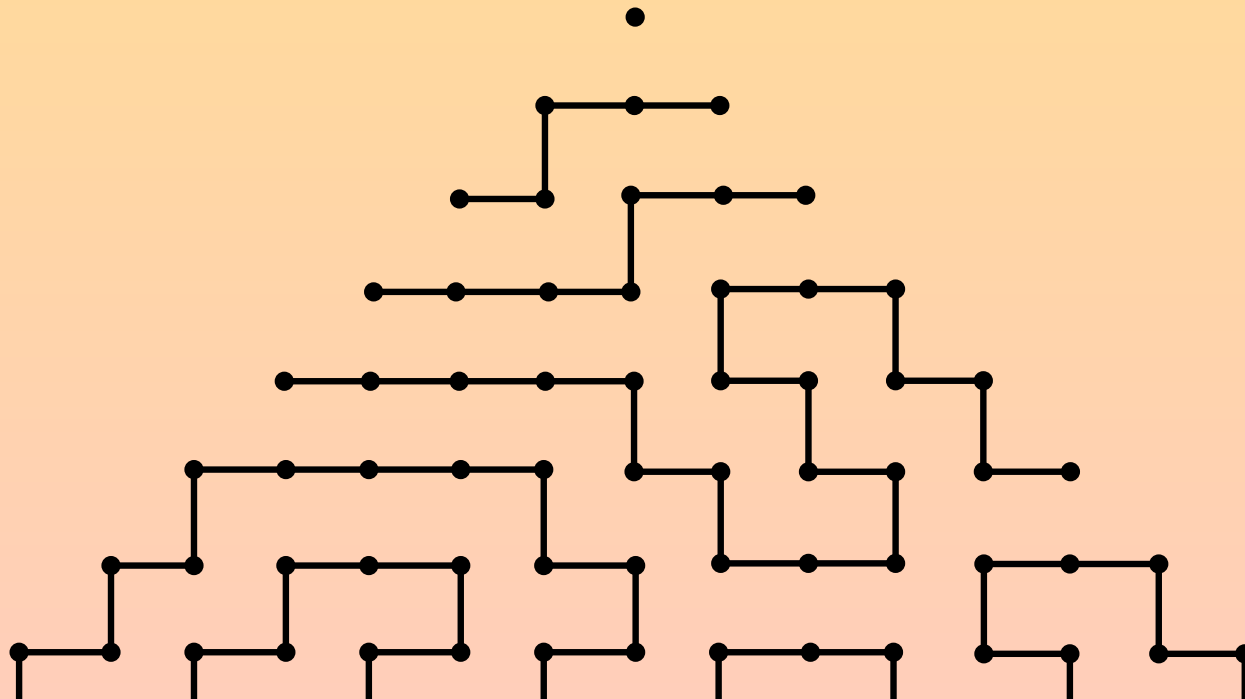


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The result is a TFPL configuration, with boundary data  $\sigma, \pi, \tau$ .

To finish the proof, one checks that this map  $\Phi$  is :

- **well defined**, i.e.  $\Phi(\text{puzzle})$  is fully packed, and verifies the boundary data  $\sigma, \pi, \tau$
- **injective** ;



# Conclusion

This diagram shows the possible indices for the numbers  $t_{\sigma, \tau}^{\pi}$  when  $\pi$  is fixed; in blue are coefficients we managed to compute, and in red are those involved in the definition of the  $c_{\alpha\pi}$ .

