

Coherence modulo relations and double groupoids

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Journées LHC

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I. Introduction and motivations

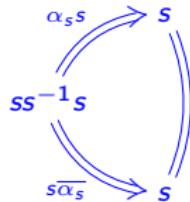
Motivations: algebraic context

- ▶ This work is part of **algebraic rewriting**, consisting in applying rewriting methods to study intrinsic properties of algebraic structures presented by generators and relations.
- ▶ For instance, computation of **syzygies** (relations among relations): for the group $\mathbb{Z}^3 = \langle x, y, z \mid [x, y] = 1, [y, z] = 1, [z, x] = 1 \rangle$, the Jacobi identity

$$[x^y, [y, z]][y^z, [z, x]][z^x, [x, y]] = 1$$

is such a syzygy, with $[x, y] = xyx^{-1}y^{-1}$ and $x^y = y^{-1}xy$.

- ▶ For monoids or categories, Squier's theorem gives a generating family for syzygies from a finite convergent presentation, **Guiraud-Malbos '09**, **Gaussent-Guiraud-Malbos '14**, **Hage-Malbos '16**.
- ▶ If a group $G = \langle X \mid R \rangle$ is presented as a monoid $M = \langle X \coprod \overline{X} \mid R \cup \{xx^{-1} \xrightarrow{\alpha_x} 1, x^{-1}x \xrightarrow{\overline{\alpha_x}} 1\}$, the confluence diagram



is an artefact induced by the algebraic structure and should not be considered as a syzygy.

Motivation: objectives

- ▶ **Objective:** Study diagrammatic algebras arising in **representation theory** using **algebraic rewriting**.
 - ▶ Khovanov-Lauda-Rouquier (KLR) algebras for categorification of quantum groups;
 - ▶ Temperley-Lieb algebras in statistical mechanics;
 - ▶ Brauer algebras and Birman-Wenzl algebras in knot theory.
- ▶ **Main questions:**
 - ▶ **Coherence** theorems; ✓
 - ▶ **Categorification** constructive results;
 - ▶ Computation of **linear bases** for these algebras by rewriting.
- ▶ Structural rules of these algebras make the study of local confluence complicated.

Example: Isotopy relations

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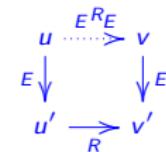
- ▶ We use **rewriting modulo**.
 - ▶ Algebraic axioms are not rewriting rules, but taken into account when rewriting.

Three paradigms of rewriting modulo

- ▶ Rewriting system R :
 - ▶ Usual rewriting theory;
 - ▶ Squier's theorem expressed in n -categories.
- ▶ In **rewriting modulo**, we consider a rewriting system R and a set of equations E .
- ▶ 3 paradigms of rewriting modulo:
 - ▶ Rewriting with R modulo E , Huet '80

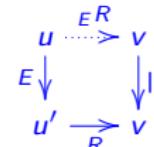


- ▶ $E R_E$: Rewriting with R on E -equivalence classes



- ▶ Rewriting with any system S such that $R \subseteq S \subseteq E R_E$, Jouannaud - Kirchner '84.

- ▶ Many results in rewriting modulo are expressed for $E R$.



II. Double groupoids

Double groupoids

- We introduce a cubical notion of coherence, related to n -categories enriched in **double groupoids**.
- A **double category** is an internal category $(\mathbf{C}_1, \mathbf{C}_0, \partial_-^{\mathbf{C}}, \partial_+^{\mathbf{C}}, \circ_{\mathbf{C}}, i_{\mathbf{C}})$ in \mathbf{Cat} .

$$\begin{array}{ccc} (\mathbf{C}_0)_0 & \xrightarrow{(\mathbf{C}_1)_0} & (\mathbf{C}_0)_0 \\ (\mathbf{C}_0)_1 \downarrow & \parallel \downarrow & \downarrow (\mathbf{C}_0)_1 \\ (\mathbf{C}_0)_0 & \xrightarrow{(\mathbf{C}_1)_0} & (\mathbf{C}_0)_0 \end{array}$$

- It gives four related categories

$$\begin{array}{ll} \mathbf{C}^{vo} := (\mathbf{C}^v, \mathbf{C}^o, \partial_{-,0}^v, \partial_{+,0}^v, \circ^v, i_0^v), & \mathbf{C}^{ho} := (\mathbf{C}^h, \mathbf{C}^o, \partial_{-,0}^h, \partial_{+,0}^h, \circ^h, i_0^h), \\ \mathbf{C}^{sv} := (\mathbf{C}^s, \mathbf{C}^v, \partial_{-,1}^v, \partial_{+,1}^v, \diamond^v, i_1^v), & \mathbf{C}^{sh} := (\mathbf{C}^s, \mathbf{C}^h, \partial_{-,1}^h, \partial_{+,1}^h, \diamond^h, i_1^h), \end{array}$$

where \mathbf{C}^{sh} is the category \mathbf{C}_1 and \mathbf{C}^{vo} is the category \mathbf{C}_0 .

- Elements of \mathbf{C}^o are called **point cells**, the elements of \mathbf{C}^h and \mathbf{C}^v are called **horizontal cells** and **vertical cells** respectively and pictured by

$$\begin{array}{ccc} & & x_1 \\ & \xrightarrow{f} & \\ x_1 & & x_2 \\ & \downarrow e & \\ & & x_2 \end{array}$$

Double groupoids

- Source and target maps make elements of \mathbf{C}_s be square cells

$$\begin{array}{ccc} \cdot & \xrightarrow{\partial_{-,\mathbf{1}}^h(A)} & \cdot \\ \partial_{-,\mathbf{1}}^v(A) \downarrow & \Downarrow A & \downarrow \partial_{+,\mathbf{1}}^v(A) \\ \cdot & \xrightarrow{\partial_{+,\mathbf{1}}^h(A)} & \cdot \end{array}$$

, with identities

$$\begin{array}{ccc} x_1 & \xrightarrow{f} & x_2 \\ i_0^v(x_1) \downarrow & \Downarrow i_1^h(f) & \downarrow i_0^v(x_2) \\ x_1 & \xrightarrow{f} & x_2 \end{array}$$

$$\begin{array}{ccc} x & \xrightarrow{i_0^h(x)} & x \\ e \downarrow & \Downarrow i_1^v(e) & \downarrow e \\ y & \xrightarrow{i_0^h(y)} & y \end{array}$$

- Compositions

$$\begin{array}{ccc} x_1 & \xrightarrow{f_1} & x_2 & \xrightarrow{f_2} & x_3 \\ e_1 \downarrow & \Downarrow A & \downarrow e_2 & \Downarrow B & \downarrow e_3 \\ y_1 & \xrightarrow{g_1} & y_2 & \xrightarrow{g_2} & y_3 \end{array} \rightsquigarrow \begin{array}{ccc} x_1 & \xrightarrow{f_1 \circ^h f_2} & x_3 \\ e_1 \downarrow & \Downarrow A \circ^h B & \downarrow e_3 \\ y_1 & \xrightarrow{g_1 \circ^h g_2} & y_3 \end{array}$$

$$\begin{array}{ccc} x_1 & \xrightarrow{f_1} & x_2 \\ e_1 \downarrow & \Downarrow A & \downarrow e_2 \\ y_1 & \xrightarrow{f_2} & y_2 \\ e'_1 \downarrow & \Downarrow A' & \downarrow e'_2 \\ z_1 & \xrightarrow{f_3} & z_2 \end{array} \rightsquigarrow \begin{array}{ccc} x_1 & \xrightarrow{f_1} & x_2 \\ e_1 \circ^v e'_1 \downarrow & \Downarrow A \circ^v A' & \downarrow e_2 \circ^v e'_2 \\ z_1 & \xrightarrow{f_3} & z_2 \end{array}$$

for all x_i, y_i, z_i in \mathbf{C}^o , f_i in \mathbf{C}^h , e_i, e'_i in \mathbf{C}^v and A, A', B in \mathbf{C}_s .

Double groupoids

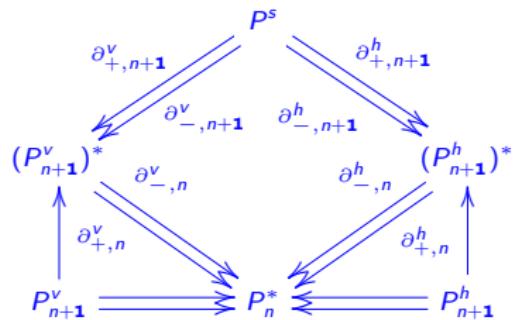
- ▶ These compositions satisfy the **middle four interchange law**:

$$\begin{array}{c} x_1 \xrightarrow{f_1} x_2 \\ e_1 \downarrow \Downarrow A \downarrow e_2 \\ y_1 \xrightarrow{g_1} y_2 \\ \diamond^h \end{array} \quad \begin{array}{c} x_2 \xrightarrow{f_2} x_3 \\ e_2 \downarrow \Downarrow B \downarrow e_3 \\ y_2 \xrightarrow{g_2} y_3 \\ \diamond^h \end{array} \quad = \quad \begin{array}{c} x_1 \xrightarrow{f_1} x_2 \\ e_1 \downarrow \Downarrow A \downarrow e_2 \\ y_1 \xrightarrow{g_1} y_2 \\ \diamond^v \end{array} \quad \begin{array}{c} x_2 \xrightarrow{f_2} x_3 \\ e_2 \downarrow \Downarrow B \downarrow e_3 \\ y_2 \xrightarrow{g_2} y_3 \\ \diamond^h \end{array}$$
$$\begin{array}{c} y_1 \xrightarrow{g_1} y_2 \\ e'_1 \downarrow \Downarrow A' \downarrow e'_2 \\ z_1 \xrightarrow{h_1} z_2 \end{array} \quad \begin{array}{c} y_2 \xrightarrow{g_2} y_3 \\ e'_2 \downarrow \Downarrow B' \downarrow e'_3 \\ z_2 \xrightarrow{h_2} z_3 \end{array} \quad \begin{array}{c} y_1 \xrightarrow{g_1} y_2 \\ e'_1 \downarrow \Downarrow A' \downarrow e'_2 \\ z_1 \xrightarrow{h_1} z_2 \end{array} \quad \begin{array}{c} y_2 \xrightarrow{g_2} y_3 \\ e'_2 \downarrow \Downarrow B' \downarrow e'_3 \\ z_2 \xrightarrow{h_2} z_3 \end{array}$$

- ▶ **Double groupoid** = double category $(\mathbf{C}_1, \mathbf{C}_0, \partial_-^{\mathbf{C}}, \partial_+^{\mathbf{C}}, \circ_{\mathbf{C}}, i_{\mathbf{C}})$ in which \mathbf{C}_1 and \mathbf{C}_0 are groupoids.
- ▶ **n -category enriched in double groupoids** = n -category \mathcal{C} such that any homset $\mathcal{C}_n(x, y)$ is a double groupoid.
 - ▶ The set $\mathcal{C}_n(x, y)^\diamond$ is the set of n -cells in $\mathcal{C}_n(x, y)$.
- ▶ Horizontal $(n+1)$ -category is the $(n+1)$ -category of **rewritings**; vertical $(n+1)$ -category is the $(n+1)$ -category of **modulo rules**.

Double $(n+2, n)$ -polygraphs

- ▶ A **double n -polygraph** is a data (P^v, P^h, P^s) made of:
 - ▶ two $(n+1)$ -polygraphs P^v and P^h such that $P_k^v = P_k^h$ for $k \leq n$,
 - ▶ a **2-square extension** P^s of the pair of $(n+1)$ -categories $((P^v)^*, (P^h)^*)$, that is a set equipped with four maps $\partial_{\pm, n}^\mu$, with $\mu \in \{v, h\}$, making Γ a **2-cubical set**:



- ▶ A **double $(n+2, n)$ -polygraph** is a double n -polygraph whose square extension P_{n+2}^s is defined on $((P^v)^\top, (P^h)^\top)$.
- ▶ A double n -polygraph (resp. double $(n+2, n)$ -polygraph) (P^v, P^h, P^s) generates a free n -category enriched in double categories (resp. in double groupoids), denoted by $(P^v, P^h, P^s) \mathbb{T}$.

Acyclicity

- A 2-square extension P^s of $((P^\nu)^\top, (P^h)^\top)$ is **acyclic** if for any square

$$S = \begin{array}{c} \cdot \xrightarrow{(P^h)^\top} \cdot \\ (P^\nu)^\top \downarrow \quad \downarrow A \quad \downarrow (P^\nu)^\top \\ \cdot \xrightarrow{(P^h)^\top} \cdot \end{array}$$

there exists a square $(n+1)$ -cell A in $(P^\nu, P^h, P^s)^\top\top$ such that $\partial(A) = S$.

- A **2-fold coherent presentation** of an n -category \mathbf{C} is a double $(n+2, n)$ -polygraph (P^ν, P^h, P^s) such that:
 - the $(n+1)$ -polygraph $P^\nu \coprod P^h$ presents \mathbf{C} ;
 - P^s is acyclic
- **Example:** Let E be a convergent $(n+1)$ -polygraph and \mathbf{C} the n -category presented by E .
 $\text{Cd}(E) :=$ square extension of $(E^\top, 1)$ containing squares

$$\begin{array}{c} \cdot \xrightarrow{=} \cdot \\ e_1 \star_{n-1} e'_1 \downarrow \quad \downarrow e_2 \star_{n-1} e'_2 \\ \cdot \xrightarrow{=} \cdot \end{array}$$

for a choice of confluence diagram of any critical branching (e_1, e_2) of E .

- From Squier's theorem, $(E, \emptyset, \text{Cd}(E))$ is a 2-fold coherent presentation of \mathbf{C} .

III. Polygraphs modulo

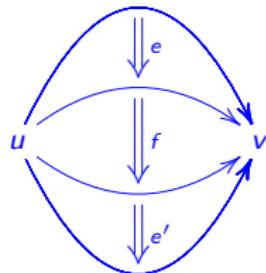
Polygraphs modulo

A *n*-polygraph modulo is a data (R, E, S) made of

- ▶ an *n*-polygraph R of **primary rules**,
- ▶ an *n*-polygraph E such that $E_k = R_k$ for $k \leq n-2$ and $E_{n-1} \subseteq R_{n-1}$, of **modulo rules**,
- ▶ S is a cellular extension of R_{n-1}^* such that $R \subseteq S \subseteq {}_E R_E$, where the cellular extension ${}_E R_E$ is defined by

$$\gamma^{{}_E R_E} : {}_E R_E \rightarrow \text{Sph}_{n-1}(R_{n-1}^*)$$

where ${}_E R_E$ is the set of triples (e, f, e') in $E^\top \times R^{*(1)} \times E^\top$ such that



and the map $\gamma^{{}_E R_E}$ is defined by $\gamma^{{}_E R_E}(e, f, e') = (\partial_{-,n-1}(e), \partial_{+,n-1}(e'))$.

Branchings and confluence modulo

- A **branching modulo E** of the n -polygraph modulo S is a triple (f, e, g) where f and g are n -cells of S^* with f non trivial and e is an n -cell of E^\top , such that:

$$\begin{array}{ccc} u & \xrightarrow{f} & u' \\ e \downarrow & & \\ v & \xrightarrow{g} & v' \end{array}$$

- It is **local** if f is an n -cell of $S^{*(1)}$, g is an n -cell of S^* and e an n -cell of E^\top such that $\ell(g) + \ell(e) = 1$.
- It is **confluent modulo E** if there exists n -cells f', g' in S^* and e' in E^\top :

$$\begin{array}{ccccc} u & \xrightarrow{f} & u' & \xrightarrow{f'} & w \\ e \downarrow & & \downarrow & & \downarrow e' \\ v & \xrightarrow{g} & v' & \xrightarrow{g'} & w' \end{array}$$

- S is said **confluent modulo E** (resp. **locally confluent modulo E**) if any branching (resp. local branching) of S modulo E is confluent modulo E .

IV. Coherence modulo

Coherent confluence modulo

- We consider Γ a 2-square extension of (E^\top, S^*) .
- A branching modulo E is Γ -confluent modulo E if there exist n -cells f', g' in S^* , e' in E^\top and an $(n+1)$ -cell A in $(E, S, E \rtimes \Gamma \cup \text{Peiff}(E, S))^{\overline{\Gamma}, v}$:

$$\begin{array}{ccccc} u & \xrightarrow{f} & u' & \xrightarrow{f'} & w \\ e \downarrow & & \downarrow A & & \downarrow e' \\ v & \xrightarrow{g} & v' & \xrightarrow{g'} & w' \end{array}$$

- $(E, S, -)^{\overline{\Gamma}, v}$ is the free n -category enriched in double categories generated by $(E, S, -)$, in which all vertical cells are invertible.
- $\text{Peiff}(E, S)$ is the 2-square extension containing the following squares for all $e, e' \in E^\top$ and $f \in S^*$.

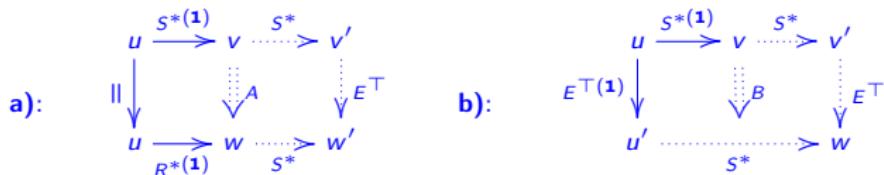
$$\begin{array}{ccc} u \star_i v \xrightarrow{f \star_i v} u' \star_i v & & w \star_i u \xrightarrow{w \star_i f} w \star_i u' \\ u \star_i e \downarrow & & e' \star_i u \downarrow \\ u \star_i v' \xrightarrow{f \star_i v'} u' \star_i v' & & w' \star_i u \xrightarrow{w' \star_i f} w' \star_i u' \end{array}$$

- $E \rtimes \Gamma$ is to avoid "redundant" elements in Γ for different squares corresponding to the same branching of S modulo E :

$$\begin{array}{ccc} u & \xrightarrow{f} & v & \xrightarrow{f'} & v' \\ e \downarrow & & & \downarrow e' & \\ u & \xrightarrow{g = e_1 g_1 e_2} & w & \xrightarrow{\quad} & w' \\ & & g' & & \end{array} \quad \text{and} \quad \begin{array}{ccc} u & \xrightarrow{f} & v & \xrightarrow{f'} & v' \\ e \star_{n-1} e_1 \downarrow & & & & \downarrow e' \\ u_1 & \xrightarrow{g_1 e_2} & w & \xrightarrow{\quad} & w' \\ & & g' & & \end{array}$$

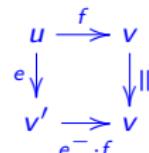
Coherent Newman and critical pair lemmas

- S is Γ -confluent modulo E (resp. locally Γ -confluent modulo E) if any of its branching modulo E (resp. local branching modulo E) is Γ -confluent modulo E .
- **Theorem:** The following assertions are equivalent:
 - S is Γ -confluent modulo E ;
 - S is locally Γ -confluent modulo E ;
 - S satisfies properties **a)** and **b)**:



for any local branching of S modulo E .

- S satisfies properties **a)** and **b)** for any critical branching of S modulo E .
- For $S = E R$, property **b)** is trivially satisfied.



- ▶ A set X of $(n-1)$ -cells in R_{n-1}^* is E -normalizing with respect to S if for any u in X ,

$$\text{NF}(S, u) \cap \text{Irr}(E) \neq \emptyset.$$

- ▶ **Theorem:** Let (R, E, S) be n -polygraph modulo, and Γ be a square extension of the pair of $(n+1, n)$ -categories (E^\top, S^\top) such that

- ▶ E is convergent,
- ▶ S is Γ -confluent modulo E ,
- ▶ $\text{Irr}(E)$ is E -normalizing with respect to S ,
- ▶ ${}_E R_E$ is terminating,

then $\Gamma \cup \text{Cd}(E)$ is acyclic.

- ▶ A normalization strategy for an n -polygraph P is a map σ that sends every $(n-1)$ -cell u to an n -cell $\sigma_u : u \rightarrow \hat{u}$.
- ▶ σ and ρ normalization strategies for S and E weakly commute if:

$$\begin{array}{ccc} u & \xrightarrow{\sigma_u} & \hat{u} \\ \rho_u \downarrow & & \downarrow \rho_{\hat{u}} \\ \tilde{u} & \xrightarrow{\quad \eta_u \quad} & \tilde{\hat{u}} \end{array}$$

Coherent extensions

- ▶ A **coherent completion modulo E** of S is a square extension denoted by $\mathcal{C}(S)$ of the pair of $(n+1, n)$ -categories (E^\top, S^\top) containing square cells $A_{f,g}$ and $B_{f,e}$:

$$\begin{array}{ccc} u \xrightarrow{f} u' \xrightarrow{f'} w & & u \xrightarrow{f} u' \xrightarrow{f'} w \\ \Downarrow \downarrow A_{f,g} \quad \downarrow e' & & \downarrow \downarrow B_{f,e} \quad \downarrow e' \\ u \xrightarrow{g} v \xrightarrow{g'} w' & & v \xrightarrow{g'} w' \end{array}$$

for any critical branchings (f, g) and (f, e) of S modulo E .

- ▶ **Corollary:** Let (R, E, S) be an n -polygraph modulo such that

- ▶ E is convergent,
- ▶ S is confluent modulo E ,
- ▶ $\text{Irr}(E)$ is E -normalizing with respect to S ,
- ▶ ${}_E R_E$ is terminating,

For any coherent completion Γ of S modulo E , $\Gamma \cup \text{Cd}(E)$ is acyclic.

- ▶ **Corollary:** Let R be an n -polygraph.

Conclusion

Example: The 2-category \mathcal{KLR}

- Let \mathcal{KLR} be the 2-linear category defined by:

- \mathcal{KLR}_0 is a set X corresponding to the weight lattice of a Kac-Moody algebra;
- $\mathcal{KLR}_1 = \{\varepsilon = (\varepsilon_1, \dots, \varepsilon_{\ell(\varepsilon)}) \text{ with } \varepsilon_i \in \{-, +\}\}$.
- \mathcal{KLR}_2 admits for generating 2-cells:



- Subject to the following relations:

- "Nil-Hecke relations" for both orientations of strands:

$$\begin{array}{c} \text{Diagram: two strands crossing, both labeled } \alpha, \text{ with a dot at the top-left, labeled } \lambda. \end{array} - \begin{array}{c} \text{Diagram: two strands crossing, both labeled } \alpha, \text{ with a dot at the top-right, labeled } \lambda. \end{array} = \begin{array}{c} \text{Diagram: two strands crossing, left labeled } \alpha, \text{ right labeled } \alpha, \text{ top-left dot, labeled } \lambda. \end{array} - \begin{array}{c} \text{Diagram: two strands crossing, left labeled } \alpha, \text{ right labeled } \alpha, \text{ bottom-right dot, labeled } \lambda. \end{array} = \begin{array}{c} \text{Diagram: two vertical strands, left labeled } \alpha, \text{ right labeled } \alpha, \text{ top-right dot, labeled } \lambda. \end{array}$$

$$\begin{array}{c} \text{Diagram: two strands crossing, both labeled } \alpha, \text{ with a dot at the top-left, labeled } \lambda. \end{array} = 0 \quad \begin{array}{c} \text{Diagram: three strands crossing, left labeled } \alpha, \text{ middle labeled } \alpha, \text{ right labeled } \alpha, \text{ top-left dot, labeled } \lambda. \end{array} = \begin{array}{c} \text{Diagram: three strands crossing, left labeled } \alpha, \text{ middle labeled } \alpha, \text{ right labeled } \alpha, \text{ top-right dot, labeled } \lambda. \end{array}$$

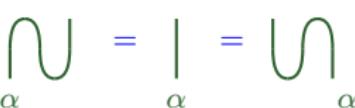
- Bubble relations:

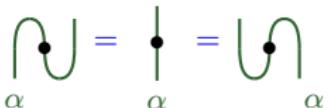
$$n \cdot \begin{array}{c} \text{Diagram: a circle with a dot at the top, labeled } \lambda. \end{array} \Rightarrow \begin{cases} 1_{\mathbf{1}_\lambda} & \text{if } n = h - 1 \\ 0 & \text{if } n < h - 1 \end{cases} \quad ; \quad \begin{array}{c} \text{Diagram: a circle with a dot at the bottom, labeled } \lambda. \end{array} \cdot n \Rightarrow \begin{cases} 1_{\mathbf{1}_\lambda} & \text{if } n = -h - 1 \\ 0 & \text{if } n < -h - 1 \end{cases}$$

Example: The 2-category \mathcal{KLR}

$$h-1+\alpha \cdot \text{Diagram} \Rightarrow - \sum_{l=1}^{\alpha} h-1+\alpha-l \cdot \text{Diagram} \cdot \text{Diagram}^{-h-1+l} \text{ for any } \lambda \in X \text{ and } \alpha > 0$$

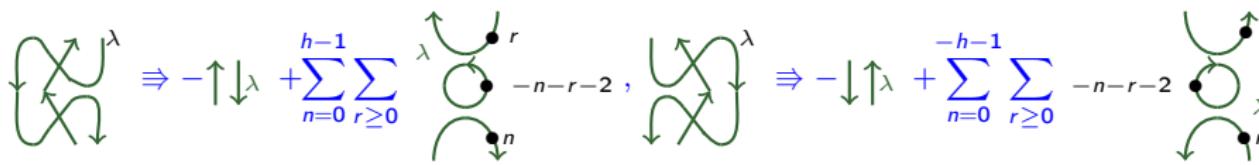
► Isotopy relations:



$$\text{Diagram}_\alpha = \text{Diagram}_\alpha = \text{Diagram}_\alpha$$


$$\text{Diagram}_\alpha = \text{Diagram}_\alpha = \text{Diagram}_\alpha$$

► "Quantum" relations



$$\text{Diagram}^\lambda \Rightarrow -\uparrow\downarrow_\lambda + \sum_{n=0}^{h-1} \sum_{r \geq 0}^{\lambda} \text{Diagram}^{\lambda-r-n-2}, \quad \text{Diagram}^\lambda \Rightarrow -\downarrow\uparrow_\lambda + \sum_{n=0}^{-h-1} \sum_{r \geq 0}^{-\lambda} \text{Diagram}^{-\lambda-r-n-2}$$



$$\text{Diagram}^\lambda \Rightarrow \sum_{n=0}^h \text{Diagram}^{-n-1}; \quad \text{Diagram}^\lambda \Rightarrow - \sum_{n=0}^{-h} \text{Diagram}^{-n-1};$$



$$\text{Diagram}^\lambda \Rightarrow - \sum_{n=0}^{-h} \text{Diagram}^{-n-1}; \quad \text{Diagram}^\lambda \Rightarrow \sum_{n=0}^h \text{Diagram}^n;$$