

ON ALMOST ORTHOGONALITY IN SIMPLE THEORIES

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ABSTRACT. 1. We show that if p is a real type which is internal in a set Σ of partial types in a simple theory, then there is a type p' interbounded with p , which is finitely generated over Σ , and possesses a fundamental system of solutions relative to Σ .
2. If p is a possibly hyperimaginary Lascar strong type, almost Σ -internal, but almost orthogonal to Σ^ω , then there is a canonical non-trivial almost hyperdefinable polygroup which multi-acts on p while fixing Σ generically. In case p is Σ -internal and T is stable, this is the binding group of p over Σ .

INTRODUCTION

In this paper we shall study the interaction of a type p (over some set A in a simple theory) with a family Σ of partial types over A . Recall that p is

- (1) *(almost) Σ -internal* if for every realization a of p there are $B \downarrow_A a$ and realizations \bar{c} of types in Σ over B , such that $a \in \text{dcl}(B\bar{c})$ (resp. $a \in \text{bdd}(B\bar{c})$).
- (2) *(almost) generated over Σ* if there is $B \supseteq A$ such that for any realization a of a p there are realizations \bar{c} of types in Σ over B with $a \in \text{dcl}(B\bar{c})$ (resp. $a \in \text{bdd}(B\bar{c})$).

In a stable theory internality and finite generation are the same, and are an important tool in the analysis of a structure (for instance in Hrushovski's proof that unidimensional stable theories are superstable). Pillay has given examples of simple theories (even of SU -rank 1) where they differ [SW02, Examples 2 and 3]. The way out seems to be almost internality and almost generation, as they agree in any simple theory. However, *definable* as opposed to *algebraic* closure played an important rôle in the definition of the *binding group* of p over Σ , namely the group $\text{Aut}(p/A \cup \Sigma)$ of all permutations of the realizations of p induced by automorphisms fixing A and all realizations of Σ . If p is Σ -internal, this group and its action on p are definable in the stable case; moreover the action is transitive if p is a strong type almost orthogonal to Σ over A . For more details, the reader may consult [Bue96, Section 4.4], [Pil96, Section 7.4], and [Poi87, Section 2.e].

Our Theorem 1.2 improves Theorem 6 of [SW02]. Recall that for $p \in S(A)$ and Σ a family of partial types over A , a *(weak) fundamental system of solutions* for p over Σ is a tuple \bar{a} of realizations of p such that every realization of p is definable (bounded) over $A\bar{a}$ together with some realizations of Σ ; note that the existence of a (weak) fundamental system of solutions implies (almost) generation. Shami and Wagner show

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how to obtain, from an almost Σ -internal real type p in a simple theory, an imaginary type p' which has a fundamental system of solutions over Σ , such that a realization of p' is interalgebraic with a finite tuple of realizations of p . Theorem 1.2 modifies that argument and obtains such a type p' whose realizations are interalgebraic with a single realization of p . It follows that Pillay's examples are, in a way, the only ones possible: Up to a finite cover almost internality and generation (even the existence of a fundamental system) are the same in a simple theory.

But obtaining a fundamental system of solutions is not even the main problem. In a stable theory, for every two tuples a and a' , the following are equivalent:

- (1) a and a' are conjugate under an automorphism fixing $A \cup \Sigma$ pointwise.
- (2) $\text{tp}(a/A \cup \Sigma) = \text{tp}(a'/A \cup \Sigma)$.
- (3) $\text{Cb}(a/A \cup \Sigma) = \text{Cb}(a'/A \cup \Sigma)$.

(here we take the canonical base of the *type*, not of the *strong type*). However, in the simple case, there is only the implication from top to bottom. The classical definition of the the binding group in a stable theory as $\text{Aut}(p/A \cup \Sigma)$ corresponds to the first condition. This is shown to be unsuitable for the simple case in [SW02], as the group thus obtained can easily be trivialized by adding a generic bipartite graph between p and Σ , which will not affect independence. The proposed solution, the group $\text{Pél}(p, \Sigma)$ of permutations of the realizations of p elementary over $A \cup \Sigma$, corresponds to the second condition (this uses existence of a weak fundamental system of solutions), but it suffers from the same defect (although it is sometimes bigger). In the second part of the paper we shall give a construction corresponding to the third and weakest condition, and therefore to the largest group (or in fact, polygroup). In fact we shall construct a generic poly-chunk multi-acting on p in the sense of [Ben02]. This allows us to invoke the machinery of [BTW, Ben03, TW01] to obtain a coreless almost A -hyperdefinable polygroup, or, over some additional parameters, an almost hyperdefinable group acting transitively on an almost hyperdefinable set X , whose generic elements are interbounded (over independent parameters) with realizations of p . If the theory was stable to start with, the group obtained will be the original binding group.

As usual, we shall fix a complete first-order theory T and work inside a monster model \mathfrak{C} ; we shall suppose throughout that T is simple. We shall follow the terminology and notation of [Wag00]; in particular the class of a tuple a modulo an equivalence relation E is denoted by a_E . All types and partial types are hyperimaginary, and tuples can be infinite (of small length), unless stated otherwise. We shall often — and have already done so in the introduction — confound a type p , or even a set of partial types Σ , with the set of its (their) realizations. We shall write $a \equiv_A a'$ for $\text{tp}(a/A) = \text{tp}(a'/A)$, and $a \equiv_A^{Ls} a'$ for $\text{lstp}(a/A) = \text{lstp}(a'/A)$. If $(x_i : i < \alpha)$ is a sequence, we put $x_{<j} = (x_i : i < j)$ for any $j \leq \alpha$.

1. FROM INTERNALITY TO A FUNDAMENTAL SYSTEM

Definition 1.1. Let Φ and Σ be two families of partial types over A . The group $\text{Pél}(\Phi, \Sigma)$ of elementary permutations of Φ over Σ is the group of all permutations τ of Φ such that for all formulas $\varphi(\bar{x}, \bar{y})$ and all $\bar{b} \in \Phi$ and $\bar{a} \in A \cup \Sigma$

$$\models \varphi(\bar{b}, \bar{a}) \quad \Leftrightarrow \quad \models \varphi(\tau(\bar{b}), \bar{a}).$$

Theorem 1.2. *Let T be simple, A boundedly closed, and suppose $p = \text{lstp}(a/A)$ is real, and almost internal in a family Σ of partial types over A . Then there is an imaginary a' definable over Aa , such that a is algebraic over Aa' , and $\text{tp}(a'/A)$ is finitely generated over Σ . Moreover, $\text{tp}(a'/A)$ has a fundamental system of solutions relative to Σ .*

PROOF. By [SW02, Fact 1] (or in fact [Wag00, Proposition 3.4.9]) there is a finite tuple \bar{a} of realizations of p which is a weak fundamental system of solutions for p over Σ , with uniform algebraicity. Put $\bar{p} = \text{tp}(\bar{a}/A)$, and $G = \text{Pél}(p, \Sigma)$; recall from [SW02] that G is normal in $\text{Pél}(p \cup \Sigma, \emptyset)$, and in particular invariant under conjugation by $\text{Aut}(p/A)$.

For any $\bar{b} \models \bar{p}$ let $C(\bar{b})$ be those $g \in G$ which fix \bar{b} pointwise. Recall that two subgroups H and K of a group G are *commensurable* if their intersection has finite index in either group; a family \mathfrak{H} of subgroups of G is *uniformly commensurable* if any two $H, K \in \mathfrak{H}$ are commensurable and the index $|H : H \cap K|$ is bounded independently of the choice of H and K . Now since any realisation of p is uniformly algebraic over $A\bar{a}$ and a finite tuple of realizations of Σ , the family $\{C(\bar{b}) : \bar{b} \models \bar{p}\}$ is uniformly commensurable. We shall need the following fact:

Fact 1.3. [Sch80, BL89, Wag00, Theorem 4.2.4] *Let G be a group and \mathfrak{H} a family of uniformly commensurable subgroups of G . Then there is $N \leq G$ which is commensurable with any $H \in \mathfrak{H}$ (necessarily uniformly), and invariant under all automorphisms of G which stabilise \mathfrak{H} setwise. Moreover, N is a finite extension of a finite intersection of groups in \mathfrak{H} .*

Let $N \leq G$ be this group associated to the family $\{C(\bar{b}) : \bar{b} \models \bar{p}\}$; note that the family, and hence N , is $\text{Aut}(p/A)$ -invariant. Let $a' = \{na : n \in N\}$, an imaginary element. If $\bar{b} \models \bar{p}$ and contains a , then $|N : N \cap C(\bar{b})|$ is finite, as is $|N : N \cap C(a)|$. It follows that a' is a finite set, $\text{Aut}(\mathfrak{C}/A, a)$ -invariant, and thus Aa -definable. Conversely, clearly $a \in \text{acl}(a')$.

As $a' \in \text{dcl}(Aa)$, G acts on $p' = \text{tp}(a'/A)$ as a group of elementary permutations over Σ , and we get a homomorphism $\text{Pél}(p, \Sigma) \rightarrow \text{Pél}(p', \Sigma)$. As any two tuples of realizations of p' which have the same type over $A \cup \Sigma$ have pre-images with the same type over $A \cup \Sigma$ (since there are only finitely many choices for the pre-image, the fact that Σ is large does not matter), this homomorphism is surjective; clearly its kernel contains N .

Prolonging \bar{a} if necessary, we may in fact assume that $C(\bar{a}) \leq N$. Let $\bar{a}' = (a' : a \in \bar{a})$. As $C(a') = NC(a)$, we get $C(\bar{a}') = N$, so the kernel of the homomorphism is precisely N , and \bar{a}' is a fundamental system of solutions for p' over Σ . \square

2. ALMOST ORTHOGONALITY

Throughout this section we fix a set Σ of partial types over \emptyset (or, equivalently, an \emptyset -invariant big subclass Σ of the monster model \mathfrak{C}). Recall that it may consist of real, imaginary or hyperimaginary elements.

2.1. Canonical bases. In [HKP00], the existence of canonical bases has been shown for Lascar strong types in a simple theory. In order not to deviate from standard notation, we shall follow the convention that $\text{Cb}(a/b)$ means $\text{Cb}(\text{lstp}(a/b))$, that is $\text{Cb}(\text{tp}(a/\text{bdd}(b)))$, and shall proceed to define a canonical base for a *type* (as opposed

to a *Laszar strong type*). Let us first notice that $\text{Cb}(a/\Sigma)$ is meaningful when a is a tuple of a fixed (small) length, even though Σ is not.

Lemma 2.1. *Let a be a hyperimaginary.*

- (1) *There exists a tuple $b \subseteq \Sigma$ such that $a \downarrow_b \Sigma$.*
- (2) *If $b, b' \subseteq \Sigma$ are such that $a \downarrow_b \Sigma$ and $a \downarrow_{b'} \Sigma$, then $\text{lstp}(a/b)$ and $\text{lstp}(a/b')$ are parallel (in fact they have a common non-forking extension).*
- (3) *If $b \subseteq \Sigma$ is such that $a \downarrow_b \Sigma$, then $\text{Cb}(a/b) \in \text{dcl}(a)$.*

PROOF.

- (1) By the local character of forking.
- (2) As $a \downarrow_{b'} b$ and $a \downarrow_b b'$.
- (3) By the previous item, an automorphism fixing a fixes $\text{Cb}(a/b)$. \square

We may therefore define $\text{Cb}(a/\Sigma) = \text{Cb}(a/b)$ for some (any) $b \subseteq \Sigma$ such that $a \downarrow_b \Sigma$. However, $\text{Cb}(a/\Sigma) = \text{Cb}(\text{tp}(a/\text{bdd}(\Sigma)))$, so this will only get us as far as defining the simple analogue of the binding group over $\text{bdd}(\Sigma)$.

In a stable theory, $\text{Cb}(a/\Sigma)$ would be the set of canonical parameters for the definition of $\text{tp}(a/\text{bdd}(\Sigma))$, and the set of orbits of such parameters over Σ would suffice to define $\text{tp}(a/\Sigma)$. There is a simple analogue:

Lemma 2.2. *Let a be a hyperimaginary, and $c = \text{Cb}(a/\Sigma)$. Then there exists $\bar{c} \in \text{dcl}(a) \cap \text{dcl}(\Sigma)$ such that for every automorphism σ*

$$\sigma(\bar{c}) = \bar{c} \iff \sigma(c) \equiv_{\Sigma} c$$

One then has $\bar{c} \in \text{dcl}(c)$, $\text{bdd}(c) = \text{bdd}(\bar{c})$, and $a \downarrow_{\bar{c}} \Sigma$.

PROOF. We know that $\text{Cb}(a/\Sigma) \in \text{bdd}(\Sigma)$, so there is $b \subseteq \Sigma$ such that $\text{tp}(c/b)$ has a unique extension to Σ . By [BPW01], the (bounded) set of conjugates of c over b (or equivalently Σ) forms a hyperimaginary $\bar{c} \in \text{dcl}(b)$; clearly $c \in \text{bdd}(\bar{c})$. An automorphism σ fixing a fixes also c ; if $c' \equiv_{\Sigma} c$, then $\sigma(\Sigma) = \Sigma$ implies $\sigma(c') \equiv_{\Sigma} \sigma(c) = c$, whence $\sigma(\bar{c}) = \bar{c}$ and $\bar{c} \in \text{dcl}(a)$. Note that this also shows $\bar{c} \in \text{dcl}(c)$. Hence $\text{bdd}(c) = \text{bdd}(\bar{c})$, and $a \downarrow_c \Sigma$ implies $a \downarrow_{\bar{c}} \Sigma$. \square

Notation 2.3. If a and \bar{c} are as above, we write $\bar{c} = \text{Cb}_{\Sigma}(a)$. As $\bar{c} \in \text{dcl}(a)$, this is in fact a hyperdefinable function on $\text{tp}(a)$.

When we write $\text{Cb}_{\Sigma}(a) = \text{Cb}_{\Sigma}(a')$, we understand implicitly that $a \equiv a'$.

If T is stable, then $\text{tp}(a/\text{Cb}_{\Sigma}(a))$ has a unique non-forking extension to Σ , and $\text{Cb}_{\Sigma}(a)$ is the canonical base for $\text{tp}(a/\Sigma)$ (the *type*, not the *strong type*). In the simple case, we have an independence theorem relative to Σ :

Lemma 2.4. *Let a and d be any hyperimaginaries, and $b, b' \subseteq \text{dcl}(\Sigma)$. Assume that $a \downarrow_{\text{Cb}_{\Sigma}(a)} d$, $b' \downarrow_{\text{Cb}_{\Sigma}(a)} d$ and $b \equiv_{\text{Cb}_{\Sigma}(a)} b'$. Then there exists $b'' \downarrow_{\text{Cb}_{\Sigma}(a)} ad$ realizing $\text{tp}(b/a) \cup \text{lstp}(b'/\text{Cb}_{\Sigma}(a)d)$.*

Note that $b \downarrow_{\text{Cb}_{\Sigma}(a)} a$ is automatically true.

PROOF. Put $c = \text{Cb}(a/\Sigma)$ and $\bar{c} = \text{Cb}_{\Sigma}(a)$. Then

$$\bar{c} \in \text{dcl}(c) \cap \text{dcl}(\Sigma) \subseteq \text{dcl}(a),$$

and $\text{bdd}(c) = \text{bdd}(\bar{c})$. Let a' be such that $ab \equiv_{\bar{c}} a'b'$, and put $c' = \text{Cb}(a'/\Sigma)$. Then $c \equiv_{\bar{c}} c'$, whereby $c \equiv_{\Sigma} c'$, and in particular $c \equiv_{b'\bar{c}} c'$. Let a'' be such that $a''b'c \equiv_{\bar{c}} a'b'c'$. Then $a'' \equiv_c a$, whence $a'' \equiv_c^{\text{Ls}} a$. Let \hat{b} be such that $a\hat{b} \equiv_c^{\text{Ls}} a''b'$. Then $\hat{b} \equiv_{\bar{c}}^{\text{Ls}} b'$ and $\hat{b} \equiv_a b$, and we may apply the independence theorem to find $b'' \perp_{\bar{c}} ad$ such that $b'' \equiv_a^{\text{Ls}} \hat{b} \equiv_a b$ and $b'' \equiv_{\bar{c}d}^{\text{Ls}} b'$. \square

Lemma 2.5. *For any x, a we have $x \in \text{bdd}(a\Sigma)$ iff $x \in \text{bdd}(a, \text{Cb}_{\Sigma}(xa))$.*

PROOF. As $xa \perp_{\text{Cb}_{\Sigma}(xa)} \Sigma$, we get $x \perp_{a, \text{Cb}_{\Sigma}(xa)} \Sigma$; the equivalence follows. \square

Note that in a stable theory, Lemma 2.5 is true even with bounded replaced by definable closure (see the claim in the proof of Lemma 2.12).

2.2. Getting a poly-chunk. We shall assume familiarity with the theory of generic actions, as developed in [Ben02]. We recall a few of the notions that are used below:

- A partial type $\pi(x)$ over A has *definable independence* if for any partial type $\pi'(y)$ over A the set $\pi(x) \wedge \pi'(y) \wedge x \perp_A y$ is type-definable. (Every complete type has definable independence.)
- A partial type $\pi(x, y, z)$ is an *invertible generic action* if
 - (1) $\text{Func}(\pi) = \pi \upharpoonright_z$ and $\text{Arg}(\pi) = \pi \upharpoonright_x$ have definable independence,
 - (2) π implies that x, y, z are pairwise independent,
 - (3) If $f \models \text{Func}(\pi)$, for any x there are at most boundedly many y , and for any y there are at most boundedly many x , such that $\models \pi(x, y, f)$. We note $f(a) = \{b : \models \pi(a, b, f)\}$ and $f^{-1}(b) = \{a : \models \pi(a, b, f)\}$.

If π is an invertible complete reduced generic action, so is $\pi^{-1}(x, y, z) = \pi(y, x, z)$.

- If $f \models \text{Func}(\pi)$ then the set of its possible *germs* is

$$\hat{f} = \{\text{Cb}(ab/f) : b \in f(a)\}$$

If $\pi(x, y, f)$ is a Lascar strong type (we say that f is complete), then its unique germ is denoted by \hat{f} .

- The *reduction* of $\pi(x, y, z)$ is

$$\bar{\pi}(x, y, \bar{z}) := \exists z [\pi(x, y, z) \wedge \bar{z} = \text{Cb}(xy/z)]$$

(so in particular, $\bar{z} \in \hat{z}$). This is a generic action, whose functions are precisely the germs of functions of π , whence the notation $\text{Germ}(\pi) = \text{Func}(\bar{\pi})$.

- Two generic actions $\pi(x, y, z)$ and $\pi'(x', y, z)$ are *equivalent*, denoted $\pi \approx \pi'$, if they have the same reduction.
- If $\pi(x, y, z)$ and $\pi'(y, w, z')$ are generic actions, then so is their composition:

$$\pi' \circ \pi(x, w, z'z) := z \perp z' \wedge x \perp z z' \wedge \exists y [\pi(x, y, z) \wedge \pi'(y, w, z')].$$

The composition is *generic* if for every independent $f \in \text{Germ}(\pi)$ and $g \in \text{Germ}(\pi')$ every $h \in \widehat{g \circ f}$ is independent from each of f and g .

Let us return to the problem of constructing an analogue of the binding group. Let p be a Lascar strong type over \emptyset (in other words, $\text{Cb}(p) \in \text{dcl}(\emptyset)$), and suppose that p is almost Σ -internal, but almost orthogonal to $\Sigma(\mathfrak{C})$ (i.e. $x \perp \Sigma(\mathfrak{C})$ for every $x \models p$). Put

$$\mathfrak{R} = \{\text{tp}(x, a) : x \models p, x \perp a, x \in \text{bdd}(a\Sigma)\},$$

and for any $r(t, u) \in \mathfrak{R}$ define

$$\pi_r(t, t', uu') = r(t, u) \wedge r(t', u') \wedge \text{Cb}_\Sigma(tu) = \text{Cb}_\Sigma(t'u') \wedge u \downarrow_{\text{Cb}_\Sigma(tu)} u'.$$

Lemma 2.6. (1) π_r is a generic action for every $r \in \mathfrak{R}$, and $\pi_r^{-1} \approx \pi_r$.
 (2) If $\text{tp}(x, a') \in \mathfrak{R}$, let $(x_i : i \leq \alpha)$ be a Morley sequence in $\text{tp}(x/a')$ for some infinite ordinal α . Then $\text{tp}(x_\alpha, x_{<\alpha}) \in \mathfrak{R}$.

PROOF.

(1) Assume that $(x, y, a_0a_1) \models \pi_r$, and put

$$c = \text{Cb}_\Sigma(xa_0) = \text{Cb}_\Sigma(ya_1) \quad \text{and} \quad d = \text{Cb}_\Sigma(a_0) = \text{Cb}_\Sigma(a_1).$$

Then $xa_0 \downarrow_c ya_1$ and $a_1 \downarrow_d \Sigma$ imply $xa_0 \downarrow_d a_1$, whence $x \downarrow_{a_0} a_1$, and $x \downarrow a_0a_1$. Since $x \downarrow c$ this also yields $x \downarrow y$. We also have $y \in \text{bdd}(a_1c) \subseteq \text{bdd}(xa_0a_1)$. To see that $\pi_r \approx \pi_r^{-1}$, note that:

$$\pi_r^{-1}(t, t', uu') = \pi_r(t', t, uu') = \pi_r(t, t', u'u).$$

Thus every inverse function of π_r is interdefinable with a function of π_r which has the same graph, and their germs are therefore equal.

(2) Set $c_i = \text{Cb}_\Sigma(x_i a') \subseteq \text{dcl}(x_i a')$ for $i \leq \alpha$. Then $(x_i c_i : i \leq \alpha)$ is a Morley sequence over a' , and

$$\text{Cb}(x_\alpha c_\alpha / a') \in \text{dcl}(x_{<\alpha} c_{<\alpha}) \subseteq \text{bdd}(x_{<\alpha} \Sigma).$$

Then $x_\alpha c_\alpha \downarrow_{\text{Cb}(x_\alpha c_\alpha / a')} a'$ implies $x_\alpha \downarrow_{c_\alpha, \text{Cb}(x_\alpha c_\alpha / a')} a' c_\alpha$; as $x_\alpha \in \text{bdd}(a' c_\alpha)$ by Lemma 2.5, we get

$$x_\alpha \in \text{bdd}(c_\alpha, \text{Cb}(x_\alpha c_\alpha / a')) \subseteq \text{bdd}(x_{<\alpha} \Sigma).$$

□

The next lemma says that if $\models \pi_r(x, y, a_0a_1)$, then we can replace r with any other $r' \in \mathfrak{R}$, and moreover control $\text{lstp}(a_0/x)$.

Lemma 2.7. Let $r, r' \in \mathfrak{R}$, and suppose $\models \pi_r(x, y, a_0a_1)$. If a'' is such that $\models r'(x, a'')$, then there are a'_0, a'_1 such that:

- (1) $\models \pi_{r'}(x, y, a'_0 a'_1)$.
- (2) $\text{Cb}(xy/a_0a_1) = \text{Cb}(xy/a'_0 a'_1)$.
- (3) $a'_0 \equiv_x^{\text{Ls}} a''$.

Proof. Since we are only interested in $\text{lstp}(a''/x)$ we may assume that $a'' \downarrow_x a_0$, whereby $x \downarrow a_0 a''$. As moreover $x \in \text{bdd}(a_0 a'' \Sigma)$, we have

$$r'' := \text{tp}(x, a_0 a'') \in \mathfrak{R}.$$

Let $c = \text{Cb}_\Sigma(xa_0) = \text{Cb}_\Sigma(ya_1)$, and let $C = \text{Cb}_\Sigma(xa_0 a'')$. Note that $c \subseteq C \cap \text{dcl}(xa_0) \cap \text{dcl}(ya_1)$. Since $xa_0 \equiv_c ya_1$, we can find C' such that $xa_0 C \equiv_c ya_1 C'$. We have $xa_0 \downarrow_c C$, $ya_1 \downarrow_c C'$ and $xa_0 \downarrow_c ya_1$. By Lemma 2.4 there is C'' such that

$$C'' \downarrow_c xa_0 ya_1, \quad C'' \equiv_{ya_1} C' \quad \text{and} \quad C'' \equiv_{xa_0 c}^{\text{Ls}} C.$$

Hence

$$xa_0 \downarrow_{C''} ya_1 \quad \text{and} \quad ya_1 C'' \equiv_c ya_1 C'' \equiv_c xa_0 C,$$

So there exist a'_0 and a'_1 such that

$$\begin{aligned} a'_0 C'' &\equiv_{xa_0 c}^{Ls} a'' C \quad \text{with} \quad a'_0 \downarrow_{xa_0 C''} ya_1, \quad \text{and} \\ ya_1 a'_1 C'' &\equiv_c xa_0 a'' C \equiv_c xa_0 a'_0 C'' \quad \text{with} \quad a'_1 \downarrow_{ya_1 C''} xa_0 a'_0. \end{aligned}$$

Therefore $C'' = \text{Cb}_\Sigma(xa_0 a'_0) = \text{Cb}_\Sigma(ya_1 a'_1)$.

By standard independence calculus we obtain that $xa_0 a'_0 \downarrow_{C''} ya_1 a'_1$, and conclude that $\models \pi_{r''}(x, y, a_0 a'_0 a_1 a'_1)$.

Let $c' = \text{Cb}_\Sigma(xa'_0)$; since $\text{Cb}_\Sigma(xa_0 a'_0) = \text{Cb}_\Sigma(ya_1 a'_1)$, we get $c' = \text{Cb}_\Sigma(ya'_1)$ as well. Then $xa'_0 \downarrow_{c'} C''$ implies $xa'_0 \downarrow_{c'} ya'_1$, so $\models \pi_{r'}(x, y, a'_0 a'_1)$. As $\pi_{r''}$ is a generic action, $x \downarrow a_0 a_1 a'_0 a'_1$, whereby

$$xy \downarrow_{a'_0 a'_1} a_0 a_1 \quad \text{and} \quad xy \downarrow_{a_0 a_1} a'_0 a'_1.$$

This implies $\text{Cb}(xy/a_0 a_1) = \text{Cb}(xy/a'_0 a'_1)$.

Finally, $a'_0 \equiv_x^{Ls} a''$ holds, as even $a'_0 \equiv_{xa_0 c}^{Ls} a''$. \square

Corollary 2.8. *We have $\pi_r \approx \pi_{r'}$ for every $r, r' \in \mathfrak{R}$. In other words, the reduction $\bar{\pi}_r$ does not depend on r .*

Let $\pi = \bar{\pi}_r$ be this common reduction. Then $\pi = \pi^{-1}$ by Lemma 2.6.1. Now Lemma 2.7 can be restated as:

If $f \in \text{Func}(\pi)$ (i.e., $f \in \text{Germ}(\pi_r)$ for some, or equivalently for every, $r \in \mathfrak{R}$) and $y \in f(x)$, then for every $r \in \mathfrak{R}$ there are a_0, a_1 such that $\models \pi_r(x, y, a_0 a_1)$ and $f = \text{Cb}(xy/a_0 a_1)$. Moreover, $\text{lstp}(a_0/x)$ may be chosen to be any extension of $r(x, u)$ to a Lascar strong type over x .

Lemma 2.9. *Let $f, g \in \text{Func}(\pi)$ be independent, let $h \in \widehat{g \circ f}$, and let this be witnessed by xyz . Then there are $r \in \mathfrak{R}$ and $(a_i : i < 3)$ such that:*

- (1) $\models r(x, a_0) \wedge r(y, a_1) \wedge r(z, a_2)$.
- (2) $\text{Cb}_\Sigma(xa_0) = \text{Cb}_\Sigma(ya_1) = \text{Cb}_\Sigma(za_2)$.
- (3) $\{xa_0, ya_1, za_2\}$ are independent over this common canonical base.
- (4) $f = \text{Cb}(xy/a_0 a_1)$, $g = \text{Cb}(yz/a_1 a_2)$, and $h = \text{Cb}(xz/a_0 a_2)$.
- (5) fgh is independent of each of the a_i .
- (6) f, g, h are pairwise independent.

Thus in particular,

$$\models \pi_r(x, y, a_0 a_1) \wedge \pi_r(x, z, a_0 a_2) \wedge \pi_r(y, z, a_1 a_2),$$

so $h \in \text{Germ}(\pi_r) = \text{Func}(\pi)$.

Moreover, there is $\bar{r}(t, \bar{u}) \in \mathfrak{R}$ such that we can always take r to be the restriction of \bar{r} to (t, u) where $u \subseteq \bar{u}$; and there is a manner to choose u such that if $r(t, u)$ and $r'(t, u')$ were both constructed in this manner then $r \cap r' = \bar{r}|_{(t, u \cap u')}$ is also in \mathfrak{R} and can serve in place of both r and r' (and in fact this holds for any “small” number of r_i , not only for two).

Proof. Let $\bar{r}(t, \bar{u}) \in \mathfrak{R}$ be constructed from any type in \mathfrak{R} as in Lemma 2.6.2, where the sequence is of length $\alpha = |T|^+$. Then for every infinite subsequence $u' \subseteq \bar{u}$ we still have $\bar{r}|_{(t, u')} \in \mathfrak{R}$.

By assumption $y \in f(x)$ and $z \in g(y)$. By Lemma 2.7 (restated) $f, g \in \text{Germ}(\pi_{\bar{r}})$, and there are $(\bar{a}_i : i < 4)$ such that

- (1) $\models \pi_{\bar{r}}(x, y, \bar{a}_0 \bar{a}_1) \wedge \pi_{\bar{r}}(y, z, \bar{a}_2 \bar{a}_3)$.
- (2) $f = \text{Cb}(xy/\bar{a}_0 \bar{a}_1)$ and $g = \text{Cb}(yz/\bar{a}_2 \bar{a}_3)$.
- (3) $\bar{a}_1 \equiv_y^{\text{Ls}} \bar{a}_2$.

As f, g, h are at worst countable hyperimaginaries, there is $\alpha_0 < |T|^+$ such that $\text{tp}(fgh/\bar{a}_i)$ does not fork over the first α_0 elements of the sequence \bar{a}_i , for all $i < 4$. If a'_i is the sequence of the remaining elements (with indices $\geq \alpha_0$), then $fgh \perp a'_i$, as \bar{a}_i is independent. Let

$$r = \bar{r}(t, \bar{u})|_{(t, u_{\geq \alpha_0})} = \text{tp}(x, a'_0) = \text{tp}(y/a'_1) = \text{tp}(y/a'_2) = \text{tp}(z, a'_3).$$

Then $r \in \mathfrak{R}$; as $x \perp \bar{a}_0 \bar{a}_1$ implies $xy \perp_{a'_0 a'_1} \bar{a}_0 \bar{a}_1$, and similarly $yz \perp_{a'_2 a'_3} \bar{a}_2 \bar{a}_3$, we get

- (1) $\models \pi_{r'}(x, y, a'_0 a'_1) \wedge \pi_{r'}(y, z, a'_2 a'_3)$,
- (2) $f = \text{Cb}(xy/a'_0 a'_1)$ and $g = \text{Cb}(yz/a'_2 a'_3)$,
- (3) $a'_1 \equiv_y^{\text{Ls}} a'_2$.

Set $d'_0 = \text{Cb}_{\Sigma}(a'_0) = \text{Cb}_{\Sigma}(a'_1) \in \text{dcl}(a'_0) \cap \text{dcl}(a'_1) \cap \text{dcl}(\Sigma)$. Then

$$a'_0 \perp_{\text{Cb}_{\Sigma}(ya'_1)} ya'_1 \implies a'_0 \perp_{d'_0} ya'_1 \implies a'_0 \perp_{a'_1} y \implies f \perp_{a'_1} y \implies f \perp_{a'_1} ya'_1 \implies xf \perp_y a'_1.$$

Similarly, we obtain $zg \perp_y a'_2$. As $xf \perp_y zg$ (this is just because xyz witness $h \in \widehat{g \circ f}$) and $a'_1 \equiv_y^{\text{Ls}} a'_2$, the independence theorem yields a_1 with

$$a_1 \equiv_{xyf}^{\text{Ls}} a'_1 \quad , \quad a_1 \equiv_{yzg}^{\text{Ls}} a'_2 \quad , \quad \text{and} \quad a_1 \perp_y xfzg.$$

Then $xf \perp_{ya_1} zg$.

Choose $a_0 \perp_{a_1 xyf} zg$ such that $a_0 a_1 \equiv_{xyf} a'_0 a'_1$, and $a_2 \perp_{a_1 yzg} a_0 xf$ such that $a_1 a_2 \equiv_{yzg} a'_2 a'_3$. Then $xf a_0 \perp_{ya_1} z g a_2$; moreover $f = \text{Cb}(xy/a_0 a_1)$ and $g = \text{Cb}(yz/a_1 a_2)$. As $\text{Cb}_{\Sigma}(x a'_0) = \text{Cb}_{\Sigma}(y a'_1)$ and $\text{Cb}_{\Sigma}(y a'_2) = \text{Cb}_{\Sigma}(z a'_3)$, we get

$$\text{Cb}_{\Sigma}(x a_0) = \text{Cb}_{\Sigma}(y a_1) = \text{Cb}_{\Sigma}(z a_2) =: c;$$

as $x a'_1 \perp_{\text{Cb}_{\Sigma}(x a'_1)} y a'_2$ and $y a'_2 \perp_{\text{Cb}_{\Sigma}(y a'_2)} z a'_3$, we see that $\{x a_0, y a_1, z a_2\}$ are independent over c .

Put $d = \text{Cb}_{\Sigma}(a_i) \in \bigcap_{i < 3} \text{dcl}(a_i)$. Then $a_1 \perp_c a_2$ implies $a_1 a_2 \perp_d c$, and

$$x a_0 \perp_c a_1 a_2 \implies x a_0 c \perp_d a_1 a_2 \implies x \perp_{a_0} a_1 a_2 \implies x \perp_{a_0} a_0 a_1 a_2 \implies x \perp_f g a_0 a_2 ,$$

whence $\text{Cb}(xz/a_0 a_2) = \text{Cb}(xz/fg) = h$.

We saw earlier that $f \perp y a'_1$, whereby $f \perp y a_1$; by symmetry $f \perp x a_0$, $g \perp y a_1$ and $g \perp z a_2$. Since we also know that $f a_0 \perp_{y a_1} g a_2$, independence calculus yields that fg is independent of each of a_0 , a_1 and a_2 ; since $h \in \text{bdd}(fg)$, the same holds for fgh .

Finally, recall that $f \perp a'_0$, whence $f \perp a_0$. Hence

$$a_1 \underset{c}{\perp} a_0 a_2 \implies a_1 \underset{d}{\perp} a_0 a_2 \implies a_1 \underset{a_0}{\perp} a_2 \implies f \underset{a_0}{\perp} h \implies f \perp h ;$$

$g \perp h$ is proved similarly.

The moreover part is clear by the construction. \square

We obtain:

Corollary 2.10. *The composition π^2 is generic, and $\pi \approx \pi^2$ (or equivalently, $\pi_r^2 \approx \pi_{r'}$ for any $r, r' \in \mathfrak{R}$).*

Proof. We saw that whenever $f, g \in \text{Func}(\pi)$ and $h \in \widehat{g \circ f}$, then f, g, h are pairwise independent, which accounts for the genericity of the composition π^2 , and that $h \in \text{Func}(\pi)$, whereby $\pi^2 \approx \pi$. \square

Theorem 2.11. *π is a generic poly-chunk in the sense of [Ben02, Definition 3.7], and $\text{Arg}(\pi) = p$. If P is the set of its germs with product given by composition, then P is a polygroup chunk and $\text{SU}(P) \geq \text{SU}(p)$.*

Moreover, if p is in a real sort, then P is in a finitary sort (by “real” we mean “real or imaginary”, as for us the important distinction is rather from hyperimaginaries).

PROOF. By [Ben02, Proposition 2.4 and Theorem 3.9].

For the moreover part, if x is real, then there is a real a such that $r = \text{tp}(x, a) \in \mathfrak{R}$. Assume now that f is a germ of π_r , say $f = \text{Cb}(xy/aa')$. Then $\text{lstp}(xy/aa') = \text{tp}(xy/(\text{dcl}(xyaa') \cap \text{bdd}(aa')))$, whereby $f \in \text{dcl}(xyaa')$. As this can be done uniformly, we see that P can be defined in a hyperimaginary sort which is a quotient of $r \times r$. \square

We can now apply tools from [BTW, Ben03]: Let \bar{P} be the core-reduct of P . Then by [BTW], for every $f \in P$ there is a group G_f almost hyperdefinable over f , whose set of generic elements is the blow-up (along \bar{f}) of P' ; by [Ben03], there is a unique coreless polygroup \tilde{P} , almost hyperdefinable over \emptyset , such that P' is the set of generic elements of \tilde{P} .

Alternatively, we can apply [TW01, Theorem 1.9] and obtain an almost hyperdefinable group G acting transitively and faithfully on an almost hyperdefinable set X , such that a generic group element is interbounded over independent parameters with a realization of $\text{Germ}(\pi)$, and a generic element of X is interbounded over independent parameters with a realization of p .

2.3. The stable case. Assume now that T is stable, and p is a Σ -internal strong type over \emptyset . We shall show that the polygroup chunk from Theorem 2.11 is the set of generic elements of the usual binding group (in particular composition is unique, and P is a group chunk), and the generic poly-chunk π is the generic action of the binding group on p .

Lemma 2.12. *Let $x_{\leq \omega}$ be a Morley sequence in p . Then $x_\omega \in \text{dcl}(x_{< \omega}, \text{Cb}_\Sigma(x_{\leq \omega}))$.*

PROOF. Let $y \models p$, and choose $A \perp y$ and $a \in \Sigma$ with $y \in \text{dcl}(Aa)$. Let $(y_i a_i : i \leq \omega)$ be a Morley sequence in $\text{stp}(ya/A)$ with $y_\omega a_\omega = ya$. Then $(y_i : i \leq \omega)$ is a Morley sequence in p ; as p is stationary, we may assume $x_i = y_i$ for $i \leq \omega$.

Claim. In a stable theory, if $a \in \text{dcl}(bc)$ and $c' = \text{Cb}(ab/c)$ (the canonical base of the type, not of the strong type), then $a \in \text{dcl}(bc')$.

PROOF OF CLAIM. Suppose $a' \equiv_{bc'} a$ with $a' \perp_{bc'} ac$. Then $a'b \equiv_{c'} ab$ and $a'b \perp_{c'} c$, whence $a'b \equiv_c ab$; as $a \in \text{dcl}(bc)$, we get $a = a'$. \square

It follows that $x \in \text{dcl}(a, \text{Cb}(xa/A))$. But since $(x_i a_i : i \leq \omega)$ is a Morley sequence over A , we have

$$\text{Cb}(xa/A) \in \text{dcl}(x_i a_i : i < \omega) \subseteq \text{dcl}(x_{<\omega} \Sigma),$$

whence $x \in \text{dcl}(x_{<\omega} \Sigma)$, and $x \in \text{dcl}(x_{<\omega}, \text{Cb}_\Sigma(x_{<\omega}))$ by the claim again. \square

Fix $r(t, u) = \text{tp}(x_\omega, x_{<\omega})$ where $(x_i : i \leq \omega)$ is a Morley sequence in p . Note that r is precisely the type constructed in Lemma 2.6.2, and $r \in \mathfrak{R}$.

Recall that a *generic function* on p is a function f whose domain are the realizations of p independent of f (i.e. of the parameters needed to define f), and whose values are again independent of f . By definability of types in a stable theory, the relation “to agree on an independent realization of p ” is a definable equivalence relation for generic functions on p , which is a congruence for composition; the equivalence class of f is called the *germ* of f .

Lemma 2.13. *Germ(π) is a set of germs of generic functions on p , closed under inverse and generic composition.*

PROOF. If $\models \pi_r(t, t', uu')$, then $\text{Cb}_\Sigma(tu) \in \text{dcl}(tu)$, and $t' \in \text{dcl}(u', \text{Cb}_\Sigma(t'u'))$ by Lemma 2.12. Hence $t \in \text{dcl}(t, uu')$; as $t \perp uu'$ by Lemma 2.6.1, we see that $\text{Func}(\pi)$ is a set of generic functions on p (i.e. π is a well-defined generic action on p in the terminology of [Ben02]). We know that π is closed under inverse and generic composition; in a stable theory, the reduction of a well-defined generic action is well-defined (i.e. $\text{Germ}(\pi)$ is a set of germs of generic functions on p), and the composition of well-defined actions is again well-defined and corresponds to the composition of germs. \square

So the construction gives a generic group chunk which acts generically on p . We want to show that this is the generic part of the binding group of p over Σ . Let G be the binding group of p over Σ , i.e. the group of permutations of p induced by automorphisms fixing Σ pointwise.

Lemma 2.14. *Any $g \in G$ induces a germ \bar{g} of a generic function on p . If $(x_i : i < \omega)$ is a Morley sequence in p independent of g , then $g \in \text{dcl}(x_{<\omega}, g(x_{<\omega}))$.*

PROOF. If $x \models p$ with $x \perp g$, then $g(x) \models p$. But $\text{tp}(g(x)) = p = \text{tp}(x)$, $\text{tp}(x/g)$, and $\text{tp}(g(x)/g)$ have the same left stratified ranks with respect to that action; as the stratified ranks witness forking, $g(x) \perp g$, so g is a generic function on p and induces a germ.

Now suppose g and g' have the same type over $x_{<\omega} g(x_{<\omega})$. Any $y \models p$ independent of $x_{<\omega}$ is in $\text{dcl}(x_{<\omega} \Sigma)$; since g and g' fix Σ , we must have $g(y) = g'(y)$. Hence g and g' agree on all realizations of p independent of $x_{<\omega}$, and thus on any $y \models p$ (just consider a Morley sequence $y_{<\omega}$ independent of $x_{<\omega} g(x_{<\omega}) y$). \square

Define

$$\begin{aligned} \tilde{\pi}(t, t', x_{<\omega} x'_{<\omega}) &:= r(t, x_{<\omega}) \wedge r(t', x'_{<\omega}) \wedge \text{Cb}_\Sigma(tx_{<\omega}) = \text{Cb}_\Sigma(t'x'_{<\omega}) \\ &\wedge t \perp x_{<\omega} x'_{<\omega} \wedge t' \perp x_{<\omega} x'_{<\omega} \end{aligned},$$

so $\text{Germ}(\tilde{\pi})$ yields a set of germs of generic functions on p . As every $g \in G$ fixes Σ , we have $\text{Cb}_\Sigma(x_{<\omega}) = \text{Cb}_\Sigma(g(x_{<\omega}))$ for any Morley sequence $x_{<\omega}$ in p independent of g . Therefore $x_{<\omega}g(x_{<\omega}) \in \text{Func}(\tilde{\pi})$, and its (unique) germ $\overline{x_{<\omega}g(x_{<\omega})}$ does not depend on the choice of $x_{<\omega}$. It follows that

$$\tau : g \mapsto \overline{x_{<\omega}g(x_{<\omega})}$$

is an embedding preserving multiplication (composition) and inverse. Since the germ $\overline{x_{<\omega}g(x_{<\omega})}$ maps any $t \models p$ independent of $x_{<\omega}g(x_{<\omega})$ to $g(t)$, the embedding τ also preserves the action on p .

Lemma 2.15. τ is surjective, and thus a group isomorphism.

PROOF. Consider $tt'x_{<\omega}x'_{<\omega} \models \tilde{\pi}$. Then $tx_{<\omega}$ and $t'x'_{<\omega}$ have the same type over $\text{Cb}_\Sigma(tx_{<\omega})$, and hence over Σ . It follows that there is an automorphism σ fixing Σ pointwise, and mapping $tx_{<\omega}$ to $t'x'_{<\omega}$. If $g \in G$ is the element induced by σ , then $\tau(g) = \overline{x_{<\omega}x'_{<\omega}}$. \square

Lemma 2.16. $\text{Germ}(\pi)$ is a subset of $\text{Germ}(\tilde{\pi})$ containing all generic types.

PROOF. $\text{Germ}(\pi)$ is clearly a subset of $\text{Germ}(\tilde{\pi})$ closed under inverse and independent multiplication; we have to show that $\text{Germ}(\pi)^2 = \text{Germ}(\tilde{\pi})$. But given two Morley sequences $x_{<\omega}$ and $x'_{<\omega}$ with $\text{Cb}_\Sigma(x_{<\omega}) = \text{Cb}_\Sigma(x'_{<\omega})$, consider a third Morley sequence $x''_{<\omega} \downarrow_{\text{Cb}_\Sigma(x_{<\omega})} x_{<\omega}x'_{<\omega}$ with $\text{Cb}_\Sigma(x''_{<\omega}) = \text{Cb}_\Sigma(x_{<\omega})$. Then $\overline{x_{<\omega}x''_{<\omega}}$ and $\overline{x''_{<\omega}x'_{<\omega}}$ are both in $\text{Germ}(\pi)$, and their composition is $\overline{x_{<\omega}x'_{<\omega}}$. \square

Finally, we note that a generic group chunk is coreless; in this case the group construction of [BTW] is hyperdefinable (and in fact definable, since equality of germs is definable by stability). The group obtained there is definably isomorphic to the binding group (or to $\text{Germ}(\tilde{\pi})$), since a generic group chunk determines its group up to isomorphism.

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