# **ON INTEGERS OF THE FORM** $p + 2^k$

# LAURENT HABSIEGER AND XAVIER-FRANÇOIS ROBLOT

ABSTRACT. We investigate the density of integers that may be written as  $p + 2^k$ , where p is a prime and k a nonnegative integer.

### 1. INTRODUCTION

Troughout this paper, the symbol p will denote a prime and k will be a nonnegative integer. Romanov [5] proved that the integers of the form  $p + 2^k$  have positive density. He also raised the following question : does there exists an arithmetic progression consisting only of odd numbers, no term of which is of the form  $p + 2^k$ ? Erdős [1] found such an arithmetic progression by considering integers which are congruent to 172677 modulo  $5592405 = (2^{24} - 1)/3$ . Thus the density of numbers of the form  $p + 2^k$ is less than 1/2, the trivial bound obtained from the odd integers. For convenience we introduce

$$\underline{d} = \liminf_{x \to \infty} \frac{\#\{p + 2^k \le x\}}{x/2} \quad \text{and} \quad \overline{d} = \limsup_{x \to \infty} \frac{\#\{p + 2^k \le x\}}{x/2}.$$

The aim of this paper is to give an explicit version of the estimates  $0 < \underline{d} \leq \overline{d} < 1$ .

# **Theorem 1.** We have

$$0.1866 < \underline{d} \le \overline{d} < 0.9819$$
 .

This range is pretty large and Bombieri conjectured the more precise value 0.868 (see [4]).

In section 2, we obtain the lower bound  $0.1866 < \underline{d}$ , by slightly refining a straightforward application of a recent result of Pintz and Ruzsa [3], in their study of Linnik's approximation of Goldbach problem (see also [2,3]). In section 3, we get the upper bound, using computations on residue classes.

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### 2. The lower bound

Let N be a large integer and put  $L = \lfloor \log N / \log 2 \rfloor$ . Define the functions

$$r(n) = \#\{(p,k) : n = p + 2^k, p \le N, 1 \le k \le L\}$$

and

$$s(N) = \#\{(p_1, p_2, k_1, k_2) : p_1 - p_2 = 2^{k_2} - 2^{k_1}, p_j \le N, 1 \le k_j \le L, j = 1, 2\}$$

so that

$$s(N) = \sum_{n=1}^{N} r^2(n)$$
.

Pintz and Ruzsa [3] proved the following lemma.

**Lemma 1.** For N large enough, we have

$$s(N) \le \frac{2}{\log^2 2} CN \,,$$

where C < 5.3636.

Let d(N) denote the number of positive integers  $n \leq N$  which may be written in the form  $n = p + 2^k$ . The Cauchy-Schwarz inequality implies easily that

$$(\pi(N)L)^2 \le d(N)s(N)\,,$$

where  $\pi(N)$  denotes the number of primes  $p \leq N$ . We deduce from Lemma 1 and from the prime number theorem that  $2Cd(N) \geq (1 + o(1))N$ , and the lower bound  $\underline{d} \geq 1/C > 0.1864$  follows from the definitions.

To get the bound from the theorem, we need further notations. Put

$$\epsilon_N = \frac{\sum_{1 \le n \le N, r(n) > 0} r(n)}{\sum_{1 \le n \le N, r(n) > 0} 1} \quad \text{and} \quad \epsilon = \frac{2}{\underline{d} \log 2}.$$

By the definitions, there exists a subsequence of  $(\epsilon_N)_{N \in \mathbb{N}}$  which converges to  $\epsilon$ . Let us now refine the Cauchy-Schwarz inequality by studying

$$\Delta_N = \sum_{1 \le n \le N, r(n) > 0} (r(n) - \epsilon_N)^2,$$

so that

$$\Delta_N = \sum_{1 \le n \le N} r^2(n) - \frac{\left(\sum_{1 \le n \le N} r(n)\right)^2}{\sum_{1 \le n \le N, r(n) > 0} 1} = s(N) - \frac{(\pi(N)L)^2}{d(N)}$$
$$\le \left(5.3636 - \frac{1}{\underline{d}} + o(1)\right) \frac{2N}{\log^2 2},$$

for infinitely many N. Without loss of generality we may assume that  $\epsilon \in ]15, 15.5[$ : otherwise we would get either  $\underline{d} \geq 0.19$  which would be better, or  $\underline{d} \leq 0.1862$  which is false. For infinitely many N we thus have

$$\Delta_N \ge \sum_{1 \le n \le N, \ r(n) > 0} (15 - \epsilon_N)^2 \ge \left( \sum_{1 \le n \le N, \ r(n) > 0} (15 - \epsilon)^2 + o(1) \right) N$$
$$= \left( \frac{d}{2} \left( 15 - \frac{2}{d \log 2} \right)^2 + o(1) \right) N.$$

We deduce from these estimates the inequality

$$\frac{\underline{d}}{2} \left( 15 - \frac{2}{\underline{d} \log 2} \right)^2 \le \frac{2}{\log^2 2} \left( 5.3636 - \frac{1}{\underline{d}} \right) \,,$$

which may be written as

$$56.25 \log^2 2\underline{d}^2 - (15 \log 2 + 5.3636)\underline{d} + 1 \le 0.$$

The lower bound  $\underline{d} \ge 0.1866$  then follows.

# 3. The upper bound

#### A. Basic ideas.

Let us introduce further notations. Let M be a positive odd integer and let  $\omega$  denote the order of 2 in  $(\mathbb{Z}/M\mathbb{Z})^*$ . For  $\overline{m}$  a residue class modulo M, put

$$f_M(\overline{m}) = \{ \overline{k} \in \mathbb{Z}/\omega\mathbb{Z} : \overline{m} - 2^{\overline{k}} \in (\mathbb{Z}/M\mathbb{Z})^* \}$$

and

$$\delta_M(\nu) = |\{\overline{m} \in \mathbb{Z}/M\mathbb{Z} : |f_M(\overline{m})| = \nu\}|.$$

The basic tool to get an upper bound for  $\overline{d}$  is the following lemma.

**Lemma 2.** With the previous notations, we have

$$\overline{d} \le \sum_{\nu=0}^{\omega} \delta_M(\nu) \min\left(\frac{1}{M}, \frac{2\nu}{\omega\varphi(M)\log 2}\right) \,,$$

where  $\varphi$  denotes Euler's function.

*Proof.* Let  $\overline{m}$  be a congruence class modulo M, with  $|f_M(\overline{m})| = \nu$ . Let us study the proportion of odd integers congruent to  $\overline{m}$  that may be written in the form  $p + 2^k$ . This proportion is clearly at most 1/M, and we only need to prove the alternative upper bound.

Since all the primes but a finite number are invertible modulo M, there exist  $\nu$  congruence equations  $\overline{m} = \overline{p_i} + 2^{\overline{k_i}}$ ,  $i \in \{1, \ldots, \nu\}$ , such that all but finitely many representations  $p + 2^k$  come from one of these congruence equations. The number of primes up to N which are congruent to  $p_i$  modulo M is asymptotic to  $N/(\varphi(M) \log N)$ , while the number of powers of 2 which are congruent to  $2^{k_i}$  modulo M is asymptotic to

 $\log N/(\omega \log 2)$ . Thus the number of integers congruent to  $\overline{m}$  that may be written in the form  $p + 2^k$  is at most  $(\nu/(\varphi(M)\omega \log 2) + o(1))N$ . This implies that the proportion of odd integers enjoying these properties is at most  $2\nu/(\varphi(M)\omega \log 2)$  and the lemma follows.

This lemma provides a non trivial upper bound for  $\overline{d}$  as soon as there exist residue classes  $\overline{m}$  modulo M such that

$$f_M(\overline{m}) < \frac{\omega\varphi(M)\log 2}{2M},$$
 (1)

a condition that occur for a small number of classes. The main problem is to compute the distribution of the  $f_M(\overline{m})$ 's in an efficient way. The direct computation of all the  $f_M(\overline{m})$ 's is quickly limited by memory problems. However one can obtain significant results this way.

Take  $M = 23205 = (2^{24} - 1)/723$ , so that  $\omega = 24$  and  $\varphi(M) = 9216$ . The condition (1) is equivalent to  $f_M(\overline{m}) \leq 3$ . We find

$$(\delta_M(0), \delta_M(1), \delta_M(2), \delta_M(3)) = (0, 48, 720, 320),$$

and we get this way  $\overline{d} \leq 0.985049$ .

### B. Refined algorithms and results.

It appears that the function  $f_M$  takes very few possible values, when compared to the subset set of  $\mathbb{Z}/\omega\mathbb{Z}$ . So let us introduce

$$g_M(I) = \{\overline{m} \in \mathbb{Z}/M\mathbb{Z} : f_M(\overline{m}) = I\}$$
 and  $G_M(I) = |g_M(I)|,$ 

for  $I \subset \mathbb{Z}/\omega\mathbb{Z}$ . Note that

$$\delta_M(\nu) = \sum_{|I|=\nu} G_M(I) \,.$$

So it is sufficient to know the distribution of the  $G_M(I)$ 's to compute an upper bound for  $\overline{d}$ .

The main advantage of the function  $g_M$  is that it is easily computable by induction on the number of prime factors of M. The initial case is given by  $g_0(\{0\}) = \{0\}$ .

Let  $M_1$ ,  $M_2$  be two positive odd squarefree integers, with  $M_2 = pM_1$ for some prime p not dividing  $M_1$ . Let  $\omega_1$ ,  $\omega_2$  and  $\omega_p$  denote the order of 2 in  $(\mathbb{Z}/M_1\mathbb{Z})^*$ ,  $(\mathbb{Z}/M_2\mathbb{Z})^*$  and  $(\mathbb{Z}/p\mathbb{Z})^*$ , respectively. The image of  $f_p$  is easy to compute. There is the subset  $I_{p,0} = \{\bar{2}^{\bar{k}} \in (\mathbb{Z}/p\mathbb{Z})^* : \bar{k} \in \mathbb{Z}/\omega_p\mathbb{Z}\}$ with  $G_p(I_{p,0}) = p - \omega_p$ , for each  $\bar{j} \in \mathbb{Z}/\omega_p\mathbb{Z}$  the subset  $I_{p,\bar{j}} = \{\bar{2}^{\bar{k}} \in$  $(\mathbb{Z}/p\mathbb{Z})^* : \bar{k} \in \mathbb{Z}/\omega_p\mathbb{Z}, \ \bar{k} \neq \bar{j}\}$  with  $G_p(I_{p,\bar{j}}) = 1$ . Now, let  $I_2$  and  $I_p$  be in the image of  $f_{M_2}$  and  $f_p$  respectively. Denote by  $\tilde{I}_2$  and  $\tilde{I}_p$  the subsets of  $\mathbb{Z}/M_1\mathbb{Z}$  which are inverse images of  $I_2$  and  $I_p$  by the map on subsets induced by the natural surjections  $\mathbb{Z}/M_1\mathbb{Z} \to \mathbb{Z}/M_2\mathbb{Z}$  and  $\mathbb{Z}/M_1\mathbb{Z} \to \mathbb{Z}/p\mathbb{Z}$ respectively. Then it is easy to see that  $\tilde{I}_2 \cap \tilde{I}_p$  is in the image of  $f_{M_1}$  with  $G_{M_1}(\tilde{I}_2 \cap \tilde{I}_p) = G_{M_2}(I_2)G_p(I_p)$ , and that all subsets in the image of  $f_{M_1}$  are obtained in this way. This construction allows to build recursively the image of  $f_M$ . It also enables us to know how many classes have the same image. Therefore, one can compute  $G_M(I)$  without knowing  $g_M(I)$ .

Let us give an example. For  $M = 5592405 = 3.5.7.13.17.241 = (2^{24} - 1)/3$ , we have  $\omega = 24$ . There are 16401 subsets in the image of  $f_M$ , which is much fewer than  $2^{24}$ . Each of these subsets is obtained in r ways, with  $1 \le r \le 250068$ . Only subsets of cardinality at most 3 lead to an improved upper bound. The empty set appears 48 times. Each of the singletons from  $\mathbb{Z}/24\mathbb{Z}$  appears 540 times. For 2-subsets, the situation is slightly more complicated to describe. The subsets of the form  $\{a, a \pm 8\}$  appear 3625 times (there are 24 of them) while those of the form  $\{a, a + 12\}$  appear 7170 times (there are 12 of them). There are 224 interesting 3-subsets, appearing 3, 6, 225 or 9520 times.

This method requires much less memory than the algorithm from the previous subsection. It is still possible to save a bit more memory. Indeed the representation problem (by an invertible plus a power of 2) is invariant when multiplied by a power of 2. So we can use a representative of a collection of subsets, each of them being obtained by translation from the representative, instead of subsets of  $\mathbb{Z}/\omega\mathbb{Z}$ .

The best result found so far is given by

$$M = 3.5.7.11.13.17.19.31.41.73.241.257$$
.

It leads to the improvement

# $\overline{d} < 0.9818818607968211912960156368$ ,

and the upper bound from Theorem 1 follows. This computation took 35 minutes on an Intel Xeon 2.4GHz with a memory stack of 2.1Go. Indeed, the real limitation is the memory. Note that during the computations, subsets for which  $G_M(I)$  was quite large and thus unlikely to contribute in the density were dropped (still there were a total of 4469837 different subsets at the end). Hence the density obtained may be a little bigger than the actual density for this value of M.

#### Addendum

The referee informed the authors that, during the refereeing process, Janos Pintz would have improved on the lower bound. In a paper to appear in Acta Math. Hung., he would show  $\underline{d} \geq 0.18734$  by a more elaborate method.

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